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A COMPILER COMPILER

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0. Introduction.

The purpose of this report is to give an informal description of a compiler compiler for languages defined by means of a Compiler Description Language based on affix grammars, together with a description of that compiler compiler itself in that Compiler Description Language (CDL). Affix grammars have been described formally in [2].

In the first chapters of this report, an explanation is given of the use of this compiler compiler. This explanation is neither very formal nor very complete, but should enable the reader to understand the description given in the last chapter of the report.

The compiler compiler is a syntax-directed compiler of the top-to-bottom variety equipped with a macro mechanism, which accepts input written in a two-level extension of CF grammar, giving as output a program in ALGOL 60. It is intended as a tool in the study and production of compilers for high level languages, such as ALGOL 68.

0.1 Notation and terminology.

A Context-Free grammar (CF grammar) consists of two distinct finite collections of symbols (, termed the collections of "nonterminal" and "terminal" symbols respectively), and furthermore of one specific non-terminal, its "starting symbol", and a finite collection of "rules".

Each rule associates with some nonterminal its "alternatives", a number of possibly empty sequences of nonterminals and terminals.

In writing down a grammar, we will write down one rule for each non-terminal, with that nonterminal as its "left-hand-side", followed by a colon (:), followed by the "right-hand-side" of the rule, followed by a point (.). The right-hand-side consists of the alternatives of the left-hand-side, each separated from the next by a semicolon (;). The symbols within one alternative are separated by commas (,). As symbols we will use only "tags", i.e., sequences of letters and digits beginning with a letter.

The notation just described, the "Van Wijngaarden notation" [3], is a variation on the more common Backus-Naur form [4], with as advantages ease of writing and reading, and the fact that the end of a rule is clearly indicated.

By a "direct production" of a nonterminal we mean any of its alternatives. By a "production" of a nonterminal x we mean a sequence of symbols, which is either a direct production of x , or is the result of replacing, in some production of x , a nonterminal y by a direct production of y .

By a "terminal production" of a nonterminal x we mean a production of x that contains no nonterminals.

By a "sentence" of a CF grammar we mean a terminal production of its starting symbol.

By the "language" of a CF grammar G we mean the collection of its sentences. This language is denoted by $L(G)$.

0.2 CF grammars of type LL(1).

By the productions of a sequence of symbols y we mean the productions of a new nonterminal x with y as its only direct production, with the exclusion of y itself.

For an alternative x , we will indicate by $\text{first}(x)$ the set of terminal symbols which are the first symbol of some production of x , and by $\text{follow}(x)$ the set of terminal symbols which immediately follow x in any production of the starting symbol.

We then say a CF grammar is "of type LL(1)" [6] if it satisfies the following 3 conditions:

- 1) For any rule $A: a_1; a_2; \dots; a_n$, the sets $\text{first}(a_1)$, $\text{first}(a_2)$, ..., $\text{first}(a_n)$ are mutually disjoint.
- 2) At most one of $a_1, \dots, a_n$ can produce an empty string ϵ .
- 3) If a_p produces ϵ , then $\text{first}(a_q)$ has no elements in common with $\text{follow}(A)$.

It is decidable whether a given CF grammar is of type LL(1). Without loss of generality we may assume that only the last alternative can produce ϵ (condition 2).

1. Syntax-directed compiling.

1.1 What is a compiler compiler?

A "program" is a text in some programming language more or less directly executable on a computer, which, when executed with some input, gives some output.

A "compiler" for some high level language is a program which, when executed with a program in that language as input, gives as output a program in some other language (preferably more directly executable).

A "compiler compiler" is a compiler which, when executed with a formal description of a compiler for some language as input, gives as output a compiler for that language.

A "family of compiler compilers" is a collection of compiler compilers which, when executed with identical input, give as output equivalent programs, possibly in different languages.

Ideally, a formal description of some programming language should be capable of yielding also a formal description of a compiler for it; certainly a formal description of a compiler may be used to define a language. The formalism for definition of a language consists of two coupled mechanisms, its syntax and semantics. Strictly speaking, the syntax should indicate whether a given sequence of symbols is a program in the language, and if so, derive a parsing tree for it, while the semantics should attach a meaning to the parsing tree of a correct program.

In the definition of most languages, for instance that of ALGOL 60, this separation is not so clean cut: one can argue that part of the syntax is given under the heading of semantics. In ALGOL 68, the "context-conditions" are an example of syntactical restrictions on programs appearing as semantics, so that the syntax by itself is not sufficient to define what is a proper program.

Formalization of the syntax of programming languages is now well enough understood to obtain part of a compiler by mechanical means, using "syntax-directed" techniques. On the other hand, there exists no formalization of the semantics which could usefully give the "semantics-directed" parts of a compiler. For this, manual techniques are used, for instance by translating to some set of macros which then reflect the semantics.

Our goal is to extend the applicability of syntax-directed techniques so far that all but the very kernel of the semantics of a compiler can be defined by syntax, so that only that kernel remains to be defined by other means. To this end, a notational system, "Compiler Description Language" (CDL) will be introduced, strongly based on an extension of CF grammars ("affix grammars"), a kind of grammars which, on the one hand, are powerful enough to define major programming languages, and, on the other hand, lend themselves to syntax-directed compiling techniques.

In the light of the foregoing discussion the Compiler Description Language can best be seen as a programming language, which allows one, at a high level of abstraction, to write compilers in a machine independent way.

Although it is our intention to describe a family of compiler compilers, the discussion will center around one particular member of that family, whose input is in Compiler Description Language, and whose output is in ALGOL 60. ALGOL 60 as an output language provides a clear enough framework for the notions involved while being broadly available. ALGOL 68 would certainly provide a better framework, but its availability is still very limited.

Other members of the same family of compiler compilers are being constructed, giving output in various assembly languages and in ALGOL 68.

1.2 Parsing according to a CF grammar.

The basis of a syntax-directed compiler is a parser, so we must first go into matters of parsing.

By a "parse" of a sentence according to a grammar we mean a sequence of substitutions of a direct production for some nonterminal, which leads from the starting symbol of the grammar to that sentence. The problem of parsing a given sequence of symbols then consists of

- 1) deciding whether that sequence of symbols is a sentence of the grammar;
- 2) if so, then constructing a parse for it.

Using a CF grammar as a generating device is simple enough, but it takes more trouble to find a practical algorithm for parsing, even though

the parsing problem can easily be shown to be decidable. Still, at the moment the parsing problem for CF grammars is wholly solved, and literature abounds with parsing methods.

These methods can be classified along various criteria:

- a) bottom up or top down
- b) general algorithm & table for the particular grammar or particular parser for that grammar only
- c) parser for general CF grammars or parser for a reduced class obtained by imposing restrictions on the grammars.

The compiler compiler to be described generates a top down particular parser, written in ALGOL 60, for a CF grammar which satisfies the LL(1) conditions.

From a grammar, the particular parser is obtained by a transsscription process where each rule is turned into a <declaration> for a parsing procedure corresponding to the nonterminal in its left-hand-side; these procedures are then embedded in an environment, which contains at least some as yet unspecified means of input and output. Several such transcriptions are possible, and we will investigate two of them.

1.2.1 The first parsing method.

The transcription process is best introduced by an example.

Assume the following part of a grammar:

```
number: digit, numbertail.                                G1
numbertail: digit, numbertail; endmarker.
```

Its transcription would run:

```
Boolean procedure number;                                P1
begin if  $\neg$  digit then goto 1;
      if  $\neg$  numbertail then goto 0;
      number:= true; goto end;
  1:0: number:= false; end;
end;
```

```

Boolean procedure numbertail;
begin if  $\neg$  digit then goto 1;
    if  $\neg$  numbertail then goto 0;
    numbertail := true; goto end;
  1: if  $\neg$  endmarker then goto 2;
    numbertail := true; goto end;
  2: 0: numbertail := false; end:
end;

```

We will assume *digit* and *endmark* to be procedures declared in the environment, such that they return true if the current symbol in the input-stream is a digit or an endmark respectively, simultaneously advancing the input by one symbol, and return false otherwise.

When we now call *number* it returns either true or false, depending on whether the input stream, starting with the current symbol, contained a sequence of digits followed by an endmarker.

Now consider another CF grammar for number 1:

number 1: digit, number 1; digit, endmarker.

G2

with transcription:

```

Boolean procedure number 1;
begin if  $\neg$  digit then goto 1;
    if  $\neg$  number 1 then goto 0;
    number 1 := true; goto end;
  1: if  $\neg$  digit then goto 2;
    if  $\neg$  endmarker then goto 0;
    number 1 := true; goto end;
  2: 0: number 1 := false; end:
end;

```

P2.a

The first transcription method fails to parse any number 1. Consider, e.g., 9 #, where 9 is a digit and # the endmarker. It finds the string begins with a digit, so advances the input, and then tries to parse the remainder as a number 1, which fails; so number 1 returns false.

Clearly, there are CF grammars for which method 1 does not work satisfactorily.

1.2.2 A second parsing method.

The trouble with grammar G2 had to do with the fact that, upon meeting a digit, the input could not safely be advanced. Assume the environment to contain an input pointer *pin*, indicating at every moment the current symbol of an input array. Advancing the input is done by incrementing *pin*, so backtracking the input can be done by restoring *pin* to its original value.

The second transcription method transforms G2 into

```

Boolean procedure number 1;                                P2.b
begin integer pold; pold := pin;
    if  $\neg$  digit then goto 1;
    if  $\neg$  number 1 then goto 1;
    number 1 := true; goto end;
1: pin := pold;
    if  $\neg$  digit then goto 2;
    if  $\neg$  endmarker then goto 2;
    number 1 := true; goto end;
2: pin := pold; number 1 := false; end:
end;

```

This second version will recognize the string 9 # as follows: it finds 9 is a digit and increments *pin*; it finds # is not a number 1, so control goes to the label 1; *pin* is restored to its old value; a digit is found, incrementing *pin*; an endmarker is found, again incrementing *pin* and the procedure returns true.

Clearly this second transcription method works for a wider class of grammars than the first did.

1.2.3 Comparison of the two methods

In this section we will compare the merits of the two methods and the conditions under which they can be used.

First some terminology: We will term a procedure x , yielding a Boolean value, and obtained from a nonterminal y by one of the transcription processes, a "parsing procedure" for y .

We say a parsing procedure x "recognizes" the nonterminal x if it satisfies the following two requirements:

- 1) x always terminates;
- 2) if the inputstream, starting with the current symbol, contains a terminal production of x , then x yields true, simultaneously advancing the input to the first symbol after that terminal production; otherwise it yields false.

We will say x "exactly recognizes" x when x recognizes x and advances the inputstream only when returning true.

In this terminology *digit* is a parsing procedure exactly recognizing a digit. By induction on the length of the inputstream, one can deduce that *number* and *numbertail* recognize number and numbertail, though not exactly, as shown by the inputstring *1a* where *a* is not an endmarker: number 1 will return false, but has then already advanced the input.

According to the first method, number 1 (P2.a) fails to recognize anything, but the second version (P2.b) exactly recognizes number 1.

For convenience we will term the first transcription method "non-restoring" and the second "restoring".

We can make a number of observations concerning the two transcription methods:

- 1) For a CF grammar its nonrestoring parser recognizes its starting symbol if and only if that grammar satisfies the LL(1) conditions. The first condition assures that, upon recognizing a terminal symbol, the input can safely be advanced without a chance of taking a wrong alternative. The second and third condition assure that no problems arise because of empty terminal productions.
- 2) For a CF grammar, its restoring parser exactly recognizes its starting symbol if that grammar satisfies the LL(1) conditions. These conditions assure that an alternative taken and found present is indeed the only correct one. They are sufficient, but not necessary: the class of grammars recognized exactly contains the LL(1) grammars.

The difference between the restoring and nonrestoring parser for a LL(1) grammar comes out only for input which does not form a sentence: while the restoring parser rejects it (recognizing exactly), the non-restoring parser may reject it but still have advanced the input, or even accept it, as is clear from the example:

```
statement: ass stat; if stat.                                G3.a
ass stat:  identifier, becomes symbol, expression.
if stat:   if symbol, bool expr, then symbol, statement.
```

With statement as starting symbol, and appropriate exact recognizers for identifier, becomes symbol, etc., the nonrestoring parser for G3 would accept not only

```
      v := 0
and    if v > 2 then v := v + 1
but also v if v > 2 then v := v + 1
```

without giving any error message. A nonrestoring parser may accept erroneous sentences without warning. This deficiency must be mended by explicitly taking possible errors into account while constructing the syntax:

```
ass stat:      identifier, rest ass stat.                    G3.b
rest ass stat: becomes symbol, expression; errormessage.
```

Errormessage should print out some appropriate warning.

Advantages of the nonrestoring parser over the restoring parser are:

- 1) Time-efficiency: the parser has less work to do than the restoring parser, which may have to perform lots of backtracking.
- 2) Memory-efficiency: because backtracking can never occur, there is no need for any array to hold all of the input, and no need for temporary storage to retain old values of *pin*.

The disadvantage of the nonrestoring parser lies in the restrictions it imposes on the grammar.

In a practical situation, the advantages of the nonrestoring parser may outweigh its disadvantage: when constructing a syntax, one has good

grounds to make it LL(1).

The compiler compiler allows one to choose between the two parsing methods, translating some rules in a restoring, others in a nonrestoring fashion. This facility is essential for making compilers as efficient as possible, but makes explaining the workings of the compiler compiler somewhat complicated. For this reason we will show mainly nonrestoring parsers in the rest of this chapter, leaving the construction of the restoring version to the reader.

1.3 Extensions to CF syntax.

In this section we will give an account of the extensions that have to be made to CF syntax in order to turn it into an appropriate input-language for the compiler compiler.

1.3.1 Embedding actions into the syntax.

In order to turn a parser into a compiler, one has to provide it with means to perform, as a side-effect of parsing, some semantic actions. To accomplish this, these actions are embedded in the syntax as follows: Consider the parsing of a number with as translation its value:

number 2: digit, action 1, numbertail 2.

G4

numbertail 2: digit, action 2, numbertail 2; endmarker.

where *action 1* assigns the value of the last digit read to some global variable *s* while *action 2* multiplies *s* by 10 and then adds the value of the last digit read to *s*.

We divide the nonterminals into "predicates" and "actions", where predicates are transcribed as described in 1.2, while actions are transcribed as insertion of a procedure call at the corresponding place of the parser:

P4

```

Boolean procedure number 2;
begin if  $\neg$  digit then goto 1;
    action 1;
    if  $\neg$  numbertail 2 then goto 0;
    number 2 := true; goto end;
1:0: number 2 := false; end:
end;

```

```

Boolean procedure numbertail 2;
begin if  $\neg$  digit then goto 1;
    action 2;
    if  $\neg$  numbertail 2 then goto 0;
    numbertail 2 := true; goto end;
1: if  $\neg$  endmarker then goto 2;
    numbertail 2 := true; goto end;
2:0: numbertail 2 := false; end:
end;

```

We will allow "primitive" actions and predicates to be defined not by a rule of the grammar but by some other means, e.g., macros. If we take as such primitive actions:

```

alpha = s := 0
beta  = s := 10 * s
and gamma = s := s + last digit read

```

then we can define the actions action 1 and action 2 in terms of those primitive actions:

```

action 1: alpha, gamma.
action 2: beta, gamma.

```

G5

with transcription:

```

procedure action 1; begin s := 0; s := s + last digit read end;
procedure action 2; begin s := 10 * s; s := s + last digit read end;

```

An example of a primitive predicate could be digit.

The transcription of a rule for an action is the <declaration> of a procedure (not a Boolean procedure) which may involve again actions and predicates, primitive or not primitive. More involved examples will follow.

1.3.2 Affixes.

The primitive actions alpha, beta and gamma mentioned in 1.3.1 are neither really primitive nor very practical: they involve a global variable *s* which may not be used for other purposes by other actions, e.g., some action taken by digit.

We want to provide parsing procedures with parameters, as a mechanism for communicating information to and from other parsing procedures.

To this end, a nonterminal may be accompanied by a number of "affixes" in the form of symbols written after it in the syntax and each preceded by a plus (+), e.g.:

```

number 3 + value:                                     G6
    digit, action 3, numbertail 3 + value.
numbertail 3 + value:
    digit, action 4, numbertail 3 + value;
    endmarker.
```

Where action 3 and action 4 are primitive actions:

```

action 3 = value := last digit read
and action 4 = value := 10 * value + last digit read
```

with transcription:

```

Boolean procedure number 3 (value); integer value;           P6
begin if  $\neg$ digit then goto 1;
    value := last digit read;
    if  $\neg$ numbertail 3 (value) then goto 0;
    number 3 := true; goto end;
1:0: number 3 := false; end;
end;
```

```

Boolean procedure numbertail 3 (value); integer value;
begin if  $\neg$  digit then goto 1;
    value := 10 * value + last digit read;
    if  $\neg$  numbertail 3 (value) then goto 0;
    numbertail 3 := true; goto end;
1: if  $\neg$  endmarker then goto 2;
    numbertail 3 := true; goto end;
2:0: numbertail 3 := false; end:
end;

```

Note that *value* is an output parameter. All parameters are called by name. Of course, also primitive actions and predicates can have parameters.

Apart from parameters, we will also need local variables, which we will indicate in the left-hand-side of the rule as symbols each preceded by a minus (-), e.g.:

```

number 4 + value - d:                                     G7
    digit 4 + d, action 5 + value + d, numbertail 4 + value.
numbertail 4 + value - d:
    digit 4 + d, action 6 + value + d, numbertail 4 + value;
    endmarker.
digit 4 + d: digit, action 5 + d + last digit read.

```

where

```

    action 5 + x + y = x := y
and  action 6 + x + y = x := 10 * x + y

```

with transcription

```

Boolean procedure number 4 (value); integer value;       P7
begin integer d;
    if  $\neg$  digit 4 (d) then goto 1;
    value := d;
    if  $\neg$  numbertail 4 (value) then goto 0;
    number 4 := true; goto end;
1:0: number 4 := false; end:
end;

```

```

Boolean procedure numbertail 4 (value); integer value;
begin integer d;
    if  $\neg$  digit 4 (d) then goto 1;
    value := 10 * value + d;
    if  $\neg$  numbertail 4 (value) then goto 0;
    numbertail 4 := true; goto end;
1: if  $\neg$  endmarker then goto 2;
    numbertail 4 := true; goto end;
2:0: numbertail 4 := false; end:
end;

```

```

Boolean procedure digit 4 (d); integer d;
begin if  $\neg$  digit then goto 1;
    d := last digit read;
    digit 4 := true; goto end;
1: digit 4 := false; end:
end;

```

Those symbols which occur in the grammar to indicate a parameter or local variable are affixes, and must be distinct from all other symbols.

Affixes occurring in the left-hand-side of a rule we term "bound" affixes of that rule if they are preceded by a plus and "free" affixes if preceded by a minus. Affixes occurring in a rule which are neither bound nor free are "terminal".

For sake of simplicity we assume the bound, free and terminal affixes to be disjoint sets. If a bound affix is used as an input parameter to a parsing procedure we will term it "inherited"; otherwise, it is "derived".

A nonterminal must, when occurring in a right-hand-side, always be accompanied by as many affixes as there are bound affixes in the left-hand-side of its defining rule.

Of course, care has to be taken that an inherited affix should get a value before being used. In an article of D.E. Knuth [5] it is shown how to investigate by graph-theoretical means whether a system of inherited and derived affixes is consistent (in his terminology, whether a system of inherited and synthesized attributes is well defined).

A nonterminal together with its affixes is termed an "affix expression", of which the nonterminal is the "handle". The handle with its appended affixes, in a definition of an affix grammar as a generating device, takes the same place as a nonterminal in the definition of a CF grammar. Affixes derive their name from the fact that they are considered to be attached to the handle of an affix expression.

1.3.3 Grouping; jumps and labels.

Consider again G2:

number 1: digit, number 1; digit, endmarker.

When recognizing a number 1, a restoring parser for this grammar has to recognize the the last digit twice. The grammar can be rewritten to G1 in order to obviate this backtracking, but then we must introduce an extra nonterminal numbertail. When we have recognized the first digit of a number 1, we expect either an endmarker or another number 1; this could be denoted as

number 5: digit, G8
(number 5; endmarker).

with as transcription:

Boolean procedure number 5; P8.a
begin if $\neg$ digit then goto 1;
 begin if $\neg$ number 5 then goto 2;
 number 5 := true; goto end;
 2: if $\neg$ endmarker then goto 3;
 number 5 := true; goto end;
 3: goto 0;
 end;
1: 0: number 5 := false; end;
end;

The restoring version of this parser runs:

```

Boolean procedure number 5;                                P8.b
begin integer pold; pold:= pin;
  if  $\neg$  digit then goto 1;
  begin integer pold; pold:= pin;
    if  $\neg$  number 5 then goto 2;
    number 5:= true; goto end;
  2: pin:= pold;
    if  $\neg$  endmarker then goto 3;
    number 5:= true; goto end;
  3:
  end;
  1: pin:= pold; number 5:= false; end;
end;

```

As a restoring parser, P8.b recognizes the same language as P2.b but without recognizing the last digit twice. As a nonrestoring parser, P8.a succeeds where P2.a failed; consequently G8 can be considered as of type LL(1). The introduction of grouping brackets allows one to construct grammars which can be parsed more efficiently, and to enlarge the applicability of nonrestoring parsers.

Still, the parser for number 5 is inefficient because it is an unnecessarily recursive procedure; the recursion in number 5 is used only to effectuate repetition: a number consists of a first digit, optionally followed by some number of digits, followed by an endmarker. We will need a device for performing repetition. Therefore we will allow both labels and jumps within one rule of a grammar, indicating a label as a symbol acting as label-identifier followed by a colon, and a jump by a colon followed by such a symbol, e.g. (with label *rep*):

```

number 6:                                                    G9
  digit,
  rep: (digit, :rep; endmarker).

```

with transcription:

```

Boolean procedure number 6;
begin if  $\neg$  digit then goto 1;
  rep: begin if  $\neg$  digit then goto 2;
    goto rep;
    2: if  $\neg$  endmarker then goto 3;
      number 6 := true; goto end;
    3: goto 0;
  end;
  1:0: number 6 := false; end;
end;

```

P9

A jump is an action and may only occur at the end of some alternative, to a label within that same rule.

Under these conditions, the introduction of jumps and labels can easily be understood as shorthand for a corresponding recursive definition, so that the language accepted by a parser is unchanged, and none of the problems associated with the unrestricted introduction of jumps and labels is encountered. As an example, the recursive counterpart of number 6 is simply:

```

number 7: digit, rep.
rep: digit, rep; endmarker.

```

1.3.4 Some more syntactic sugar.

In order to turn affix grammars into a Compiler Description Language, some more features are needed.

Apart from parameters and local variables, parsing procedures may require global variables. The variable *pin* used by restoring parsers is an example. For the sake of simplicity we will use integer variables only, as a very general and useful datatype. Thus, for a pointer *pin*, a <declaration> will appear in the parser, as

```

integer pin;

```

Furthermore, a facility is needed to declare arrays, e.g. for holding input, and, as we will see later, to be used as stacks. Lastly, a facility

is needed for declaring a Boolean variable, termed a "flag", which may then occur instead of a predicate. In 2.4 the precise notation for such declarations is given. In 2.3 it is described how one can indicate to the compiler compiler that a specific symbol is an action or predicate, defined globally, and thus to be used without further definition.

A simple mechanism for specifying macros is also described in 2.4.

1.4 What is the use of a compiler compiler?

The conventions introduced in the previous sections, and described in chapter 2, allow one to program the input to the compiler compiler in a reasonable level detail giving a reasonable efficient compiler, where especially one can leave many administrative details to the system. Of course, ALGOL 60 is not the most practical language to write compilers in. But also in the hardest machine code, subroutines (with local variables and parameters kept on a stack), global variables, labels and jumps etc. have their counterparts, so a compiler compiler can be written towards every decent machine code.

The primitive actions and predicates are not defined in the Compiler Description Language itself, but are borrowed from some other language in the form of macros. These primitive actions and predicates therefore have to be chosen with particular care, so that they can be implemented in an equivalent way on any reasonable computer.

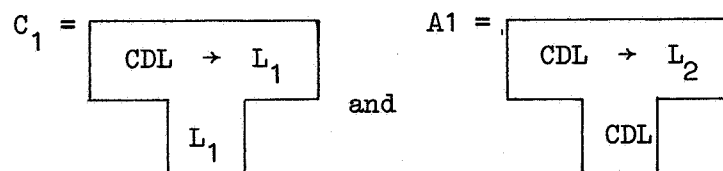
The Compiler Description Language allows one to define, in a machine independent fashion, how to construct recursively more involved actions and predicates out of the primitive ones, and out of other actions and predicates.

A compiler description in CDL describes not one single compiler, but a collection of these, one for every language to which a compiler compiler exists.

In particular, because the compiler compiler itself is defined in CDL (chapter 4) it is possible to obtain a compiler compiler to another language, and written in that other language, by a process of bootstrapping [7]:

Preparation Assume we have a compiler compiler C_1 , written in L_1 , which turns a compiler description (= text in CDL) into a compiler written in L_1 again. We then construct a compiler description A1 of a compiler compiler which turns a compiler description into a compiler written in L_2 .

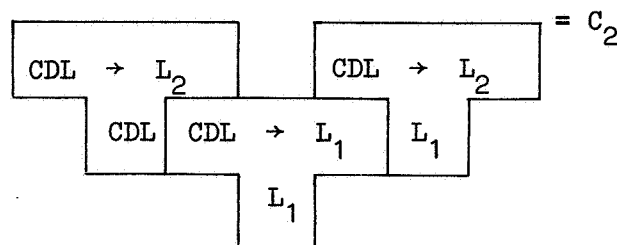
In pictures:



Step 1

We execute C_1 with A1 as input, giving as output C_2 , which is a compiler compiler written in L_1 , turning a compiler description into a compiler written in L_2 .

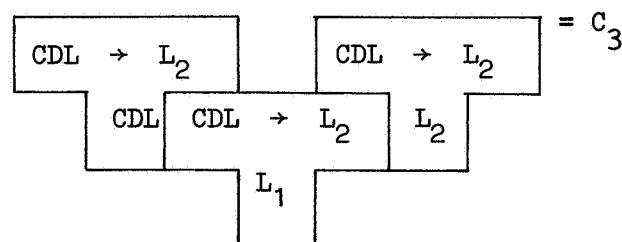
In a picture:



Step 2

We execute C_2 with A1 as input, giving as output C_3 , which is a compiler compiler written in L_2 , turning a compiler description into a compiler written in L_2 .

In a picture:



Using CDL for writing some compiler is preferable both to using machine-code and to using a general purpose high-level language as, e.g., ALGOL 60, for the following reasons:

- Compiler Description Language forms a very high framework for writing compilers, where one is little bothered with the administrative details of the parsing process.

- A Compiler Description is highly machine-independent and transferable, since a compiler can be obtained in any object language for which a compiler compiler exists. Because of their simplicity, compiler compilers for CDL can easily be made compatible.
- There is the possibility of first testing the compiler description using a high-level object code with good diagnostics, which is guaranteed never to blow up the system, and later to use a version in machine-code for production.
- By forcing oneself to use only specific programming tools, there is less possibility for tricky programming, using local features of some system which make transfer impossible.
- A Compiler Description is well enough readable to be of help in documentation.

Of course a carefully handwritten program can be more efficient both in time and in space than the mechanically obtained result of a compiler compiler. But one must not exaggerate the importance of this point: by a suitable choice of macros and careful construction of the crucial part of the syntax, one can get as near to machine-code as one should wish.

On a more general level, the Compiler Description Language can be seen as a mechanism for the syntactical composition of semantics, which should be suitable for the definition of programming languages, not by an in principle generative grammar plus verbally formulated semantics, but as a system describing the meaning of a program in terms of very simple and basic semantics actions, i.e., a machine-independent compiler.

2. About this compiler compiler.

In this chapter, a description is given of the Compiler Description Language, in the form of a context-free grammar exhibiting its syntax, accompanied by semantic remarks. For a precise definition, the description of the compiler compiler itself is to be found in chapter 4.

2.1 Compiler description.

compiler description:

- specification, compiler description;
- declaration, compiler description;
- command, compiler description;
- comment, compiler description;
- rule, compiler description;
- starting symbol.

A compiler description consists of various building stones; some of them provide only information to the compiler compiler (as specifications and commands), others give rise to a translation. The compiler compiler treats those building stones in one scan as it goes along, at the same time printing out a copy of the text with possibly some diagnostic messages, producing a translation and collecting information, some of which is displayed afterwards as a diagnostic aid (see section 2.9).

Since all work is done in one forward scan, as a general rule defining occurrences of symbols must precede applied occurrences: if one wants to use, e.g., some symbol as an action, one has to specify it as such before its defining rule. Very limited use is made of block structure: there are only two levels; symbols are either global, or local to some rule; in the latter case they may be redefined in a later rule.

2.2 Symbols.

The input is seen as a sequence of symbols, between which layout characters are ignored. A symbol is of one of four kinds:

- 1) tag, consisting of a letter, possibly followed by a number of letters or digits, between which spaces are ignored.
- 2) constant, a sequence of digits, between which spaces are ignored.
- 3) spec, one of the following reserved special symbols:

plus	+
minus	-
times	*
semicolon	;
comma	,
sub	[
sup	]
open	(
close	)
colon	:
point	.
equals	=

- 4) bold, a sequence of characters other than accents, enclosed between accents. The following bolds are reserved bold symbols:

external symbol	'external'
action symbol	'action'
predicate symbol	'predicate'
pointer symbol	'pointer'
flag symbol	'flag'
macro symbol	'macro'
list symbol	'list'
restore symbol	'restore'
unrestore symbol	'unrestore'
short symbol	'short'
long symbol	'long'
trace symbol	'trace'
untrace symbol	'untrace'
first parameter symbol	'1'
second parameter symbol	'2'

third parameter symbol '3'
 fourth parameter symbol '4'
 fifth parameter symbol '5'
 result symbol 'result'

Tags are used as nonterminals, terminals, affixes, labels, in macros, etc. Constants may occur, e.g. as inherited affixes and in macros. Non-reserved bold symbols may occur in macros.

2.3 Specifications.

specification:

external specification;
 internal specification;
 macro specification.

external specification:

external symbol, type, tag list, point.

type:

action symbol; predicate symbol;
 pointer symbol; flag symbol.

tag list:

tag, (comma, tag list;).

internal specification:

type of internal, tag list, point.

type of internal:

action symbol; predicate symbol.

macro specification:

macro symbol, type, rest macro spec, point.

rest macro spec:

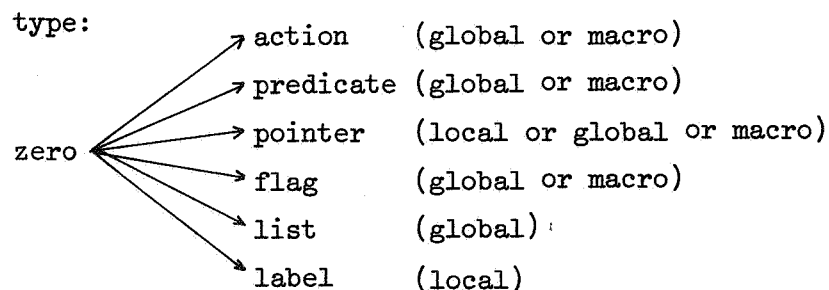
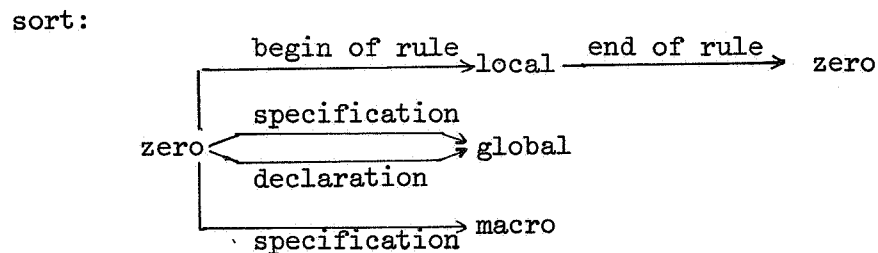
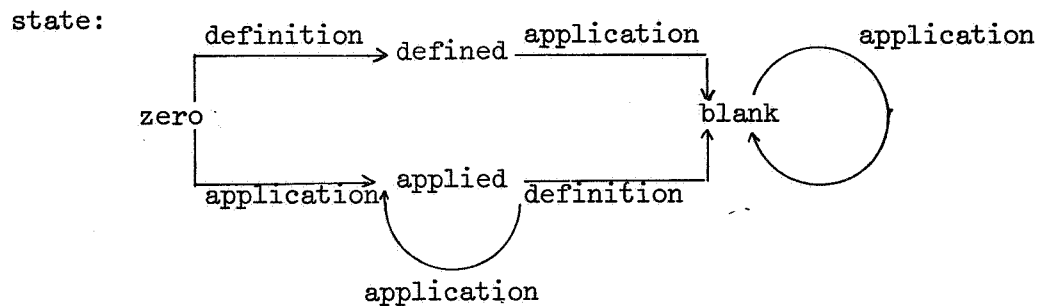
macro def,
 (comma, rest macro spec;).

macro def:

tag, equals, macrotext.

During execution of a compiler compiler with a compiler description as input, every tag occurring in the compiler description has three attributes, its "state", "sort" and "type", which may change according to fixed rules from occurrence to occurrence. Before its first occurrence, the attributes are all assumed to be zero. At every occurrence, the context may prescribe a state, sort and type, and an attempt is made to change the attributes of the tag accordingly. If such a change is not allowed, a diagnostic message is printed. Thus, it is not allowed to redefine a predicate, or use it as an affix, etc. Affixes local to a rule get, at the end of that rule, attributes zero, so that they can be used afterwards with other attributes. The attributes also serve to influence the output of the compiler compiler, so that it translates an action in a different way than a predicate.

Allowed changes of the attributes are



At the end of each rule, the state of each local should be blank (i.e. both defined and applied; if not then an error message is given); all locals again obtain state, sort and type zero.

At the end of the compiler description, the state of each global and macro should be blank, with the exception that an applied but not defined global pointer is treated as a terminal symbol, i.e. a declaration is produced for it, and also a read statement. (See 2.8.)

Specifications serve to prescribe the attributed of a tag, as follows:

	<u>state</u>	<u>sort</u>	<u>type</u>
'external' 'action'	defined	global	action
'external' 'predicate'	defined	global	predicate
'external' 'pointer'	defined	global	pointer
'external' 'flag'	defined	global	flag
'action'	applied	global	action
'predicate'	applied	global	predicate
'macro' 'action'	defined	macro	action
'macro' 'predicate'	defined	macro	predicate
'macro' 'flag'	defined	macro	flag
'macro' 'pointer'	defined	macro	pointer

An external specification indicates that the tags mentioned are supposed to have been declared externally by other means; such a tag may not be redefined and must be applied later on in the compiler description. In this way, e.g., system procedures can be used.

An internal specification indicates that the tags mentioned are to be defined later in in a specific way; it serves as an application: the tags mentioned are henceforth applied globals. Example: 'action' x.

If now a rule follows with x in the left-hand-side, then it will be translated as a <declaration> for an action-procedure. Had the rule not been preceded by this specification, then x would have been treated as a global predicate: an internal specification 'predicate' is always assumed by default.

A macro specification not only establishes the tags mentioned as

defined macros of some type, but also stores their macrotext. A macrotext is any sequence of symbols not containing a comma or a point.

Examples of parameterless macros:

```
'macro' 'pointer' max list = 5000.
```

Whenever now in the text max list appears, it will be expanded to 5000.

If a macrotext must contain a point or comma, then this symbol can be turned into a different (bold) symbol, by enclosing it between accents:

```
'macro' 'pointer' pi = 3'. '141592654.
```

On output, the accents will disappear. If one wants to output an accent, then the quote may serve as accent-image, being output as an accent. (There is, unfortunately, no quote image, so a macrotext cannot contain a quote on output). Example:

```
'macro' 'flag' true = '"true"'. 
```

which, expanded on output, gives *'true'*.

Macros can also be parameterized, with up to 5 parameters. (Another upper limit might have been chosen.) In the macrotext, the reserved bold symbols '1', '2', '3', '4' and '5' stand for first, second, up to fifth parameter respectively. For a parameter may be substituted any affix, i.e. a local, global or macro, pointer, flag or list. In 2.7 is described how and when a macro is expanded. The macro system described here is admittedly simpleminded and limited; another system might be used instead.

2.4 Declarations.

declaration:

pointer declaration;

flag declaration;

list declaration.

pointer declaration:

pointer symbol, tag list, point.

flag declaration:

flag symbol, tag list, point.

list declaration:

list symbol, rest list declaration.

rest list declaration:

tag, sub, expr, comma, expr, bus,
(comma, rest list declaration; point).

expr: tag, rest expr; constant, rest expr.

rest expr: plus, expr; minus, expr; .

A declaration has an effect on the attributes of the tags mentioned, making then defined globals.

Furthermore, a pointer declaration is translated as a <declaration> for a number of integer variables, a flag declaration as one for a number of Boolean variables, and a list declaration as one for a number of integer arrays.

Macros occurring in an expr are expanded. Example:

'macro' 'pointer' min text = 1, max text = 10000.

'list' text [min text : max text].

'pointer' stackpointer.

with translation:

integer array text [1 : 10000];

integer stackpointer;

After these declarations, text is a defined global list, min text is a ... macro pointer (both defined and applied), and stackpointer a defined global pointer.

2.5 Commands.

command: restore symbol; unrestore symbol;

short symbol; long symbol;

trace symbol; untrace symbol.

The command 'restore' causes the compiler compiler to translate rules in the restoring mode, 'unrestore' in the nonrestoring mode. The default

mode is nonrestoring.

The command 'short' causes the compiler compiler to produce compact output by suppressing most layout; 'long' causes more legible output to be produced. The default command is 'long'.

The command 'trace' causes, at the end of each rule, and at the end of the compiler description, tracing output to appear indicating attributes of all locals and globals respectively; the normal situation is 'untrace'.

2.6 Comments.

comment: sub, rest comment.

rest comment: bus; nonbus, rest comment.

(Here, nonbus stands for any symbol except].)

A comment is translated as <comment>, e.g. [input] is translated as comment input;

2.7 Rules.

rule: left hand side, middle, right hand side, point.

left hand side:

handle, optional bound affixes.

handle: tag.

optional bound affixes:

plus, bound affix, optional bound affixes;

times, bound affix, optional bound affixes; .

bound affix: tag.

middle:

optional free affixes, colon.

optional free affixes:

minus, free affix, optional free affixes; .

free affix: tag.

right hand side:

alternative,

(semicolon, right-hand-side;); .

```

alternative:
    affix expression,
        (comma, alternative;);
    group;
    jump; .
affix expression:
    label, colon, affix expression;
    handle, optional affixes.
label: tag.
optional affixes:
    plus, affix, optional affixes; .
affix: tag; constant.
group: open, right hand side, close.
jump: colon, label.

```

A rule is translated as a <declaration> for a parsing procedure corresponding to its handle, where the left-hand-side furnishes the <procedure heading>.

Bound affixes that occur preceded by a plus appear in the <specification part> specified integer, whereas those preceded by a times remain unspecified (so as to allow, e.g., arrays to be passed as parameters).

The middle of a rule is translated as begin followed by <declarations> for the free affixes, if any, possibly followed by integer pold; pold:= pin;.

The translation of the right-hand-side consists of the translation of its alternatives, followed by

```

0: handle:= false; end:
end;

```

where *handle* is an <identifier> corresponding to the handle of the rule. The <label> 0 labels the only way in which the parsing procedure can return false.

The translation of an alternative consists of the translation of its constituents, if any, followed by *handle:= true; goto end;* followed by an unique <label>.

The translation of a jump is the corresponding <go to statement>. A jump furnishes an applied occurrence of a local label.

The translation of a group is

- 1) In the restoring mode:

begin integer pold; pold:= pin; followed by the translation of its constituent alternatives, followed by end;

- 2) In the nonrestoring mode:

begin followed by the translation of its constituent alternatives, followed by goto 0 ; end;

The translation of an affix expression depends on sort and type of its handle, as follows:

Let λ stand for:

- i) in the restoring mode: the <label> at the very end of the alternative of which the affix expression is a constituent (termed the current <label>);
- ii) in the nonrestoring mode: if this affix expression is the first one of the alternative with a predicate as handle, then the current <label>, and, otherwise, the <label> 0.

Let α stand for:

- i) if the handle is a flag, then the <identifier> corresponding to it;
- ii) if the handle is global, then a call of the procedure corresponding to it, with its affixes translated as <actual parameter>s;
- iii) if the handle is a macro, then its expansion, with its first, second, etc. affix substituted for the corresponding parameter in its macro text. If the handle is a flag or predicate, then this expansion is furthermore enclosed between brackets (and).

For an affix occurring as a parameter. the following translation is given:

- i) if the affix is a constant, then the corresponding <unsigned integer>;
- ii) if the affix is a tag but not a macro, then the corresponding <identifier>;
- iii) if the affix is a macro, then its expansion (without parameter substitution; a parameterized macro is not, for the moment, allowed as an affix).

The expansion of a macro with some number of parameters is its macro text after substitution of parameters and expansion of all parameterless macros occurring in it. (A parameterized macro occurring in the macro text is not, for the moment, expanded.)

The expansion of a parameterless macro is its macro text. (Clearly, recursive macros cannot occur.)

The translation of an affix expression is:

- 1) If the handle is a flag or predicate, then

if $\neg \alpha$. then goto λ ;

- 2) If the handle is an action, then

α ;

The compiler compiler may delete some elements from this most general translation:

- 1) it deletes the assignations to the procedure identifier if the handle of the rule is an action, correspondingly changing the Boolean procedure into a procedure;
- 2) it deletes the <label>s to which no jumps occur, together with those parts of the procedure that can only be reached via those <label>s.

As an example, consider a part of a compiler description:

```
'action' set to zero.
'macro' 'action' put = '1' ['2']:= '3', incr = '1':= '1' + 1,
      make = '1':= '2'.
'macro' 'predicate' equal = '1' = '2'.
set to zero * list + min + max - p:
      make + p + min,
      rep: put + list + p + 0,
            (equal + p + max;
             incr + p, :rep).
```

Translation:

```
procedure set to zero (list, min, max); integer max, min;
begin integer p;
      p:= min;
      rep: list [p]:= 0;
            begin if  $\neg$  (p = max) then goto 2;
                  goto end;
            2: p:= p + 1; goto rep;
            end;
end;
```

A tag with type zero, when applied as an affix, becomes by context an applied global pointer. If no declaration or external specification for some applied global pointer occurs in the compiler description, then it will be treated at the end as a terminal (2.8). This is the only provision for terminal symbols: in a compiler description they occur as affixes without defining occurrence.

2.8 Starting symbol.

```
starting symbol:
      result symbol, handle, point.
```

The starting symbol is translated as a call for the procedure corres-

ponding to its handle. In front of this call, for each applied but not defined global pointer a <declaration> and a read <statement> is given. Between these <declaration>s and the read <statement>s a preparation procedure *initialize for reading* is called once. If there are no terminals, then also this call is obviated.

Example: assume one terminal, *t*, and starting symbol sentence:

```
integer t;
initialize for reading;
read (t);
sentence
end
```

Consequences:

- 1) in executing the obtained ALGOL 60 program, a representation for each terminal is to be provided as input.
- 2) If there are terminals, the following actions have to be defined in the compiler description:
read, out and initialize for reading.
- 3) In restoring parsers, the pointer *pin* must be defined. (Other representations may be taken if the relevant input to the compiler compiler is changed.)

2.9 Diagnostics.

For debugging a compiler description, various diagnostic aids are provided.

In the first place, the whole text is printed out while it is read, with two empty lines between building stones such as rules, declarations, etc. Input is assumed to come from cards, so lines have a maximum width of 80 positions. Output is assumed to proceed on cards, with a maximum of 72 positions used. A \$ in column 72 of an output card is used to indicate continuation on the next card. Continuation cards are not counted. The number of the input line appears on the left of its printout, the number of the current output card on its right.

If, at the end of a rule, some local has not been both defined and applied, then a warning message is printed.

If, at the end of the compiler description, some global has not been both defined and applied, then also a warning message appears. Various other errors are signalled by messages which should speak for themselves.

The compiler compiler makes use of a number of tables of fixed size. Whenever the upper bound of one of those tables is exceeded, but also at the end of successful execution, a "post mortem" is printed, giving the lower and upper bound of each table, with the space occupied at that moment, and finally the number of cards read and punched.

These tables, with lower and upper bounds are:

ttag	100001	110000	(tags)
tbold	200001	200300	(bolds)
tspec	300001	300100	(specs)
tcons	400001	400300	(constants)
lloc	500001	500200	(locals)
lglob	600001	601000	(globals)
lmacr	700001	701100	(macrotext)
ltext	800001	804000	(line numbers of occurrences of globals).

The bounds are such that a compiler of the size and complexity of the compiler compiler itself needs only 50% of each of those tables.

Finally, at the end of the execution, a list is printed in alphabetical order of all the globals and macros occurring in the grammar, with their linenumbers of occurrence.

3. Examples.

In this chapter some examples are given and explained of texts in CDL. They are all assumed to occur in an environment containing:

```
'unrestore'
'macro' 'action'
get = '3':= '1' ['2'], put = '1' ['2']:= '3',
make = '1':= '2',
mark = '1':= -'1',
add = '3':= '1' + '2', subtr = '3':= '1' - '2',
addmult = '4':= '1' * '2' + '3',
divrem = '3':= '1' ÷ '2'; '4':= '1' - '2' * '3',
incr = '1':= '1' + 1, decr = '1':= '1' - 1.
'macro' 'predicate'
marked = '1' < 0,
less = '1' < '2', equal = '1' = '2', lseq = '1' ≤ '2'.
'global' 'action' resym, prsym, exit.
'global' 'pointer' newlinechar, spacechar.
```

This is a rather minimal environment with which a lot can be done. The macro actions *get* and *put* are for accessing array elements; *make* performs the assignation; *mark* inverts the sign of an integer variable, which we will use as marking bit; *add* and *subtr* allow arithmetic; *addmult* and *divrem* serve for packing and unpacking small integers into a word; and *incr* and *decr* serve to increment and decrement a variable by one.

The system procedure *resym* reads one character from input (without possibility of backtracking) and assigns to its only parameter an integer, viz. the internal value of that character, which does not exceed 255. We will assume that the digits have internal value 0 - 9 respectively, and the letters the internal values 10 - 35. Complementarily, *prsym* prints out the character corresponding to the value of its parameter. The global variables *newlinechar* and *spacechar* contain the internal values of the suggested characters. A call of *exit* terminates execution.

The following set of examples introduces one by one some building stones for a compiler for an ALGOL like language. Each example may make use of objects defined in a previous example.

3.1 Reading a number.

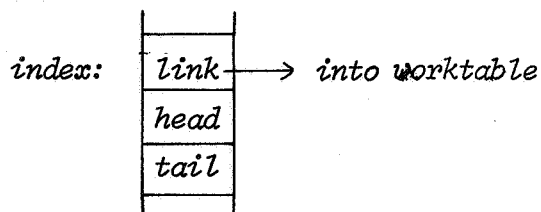
By now we have seen so many versions of number that it is time to see a good one:

```
'action' nextchar.
'pointer' char.
number + val - d:
    digit + d, make + val + d,
    rep: (digit + d, addmult + val + 10 + d + val, :rep;).
digit + d:
    lseq + char + 9, make + d + char, nextchar.
nextchar:
    resym + char,
    rep: (equal + char + spacechar, resym + char, :rep;).
```

Before the first number is read, *initialize for reading* has to take care of calling *nextchar* once.

3.2 Updating a chain.

In a worktable, we want to keep a number of chains of two elements of the form:



Each element points, via its link, to the next of the chain, the last element having link 0. Each chain is accessible by a variable which either points to its first element or contains 0 to indicate an empty chain. We want to be able to

- 1) search a chain for an element whose head has a specific value, and, if such an element exists, obtain its tail.
- 2) insert an element at the head of a chain.
- 3) see whether a specific element is already there, and in that case

obtain its index.

- 4) detach an element, if any, from the head of a chain, but leave it in the worktable (no garbage-collection).

An updating system with these properties is needed in a compiler, e.g., for collecting declarations: The most recently found declaration for a specific identifier is always the first to be looked at and the first to lose its significance, although the information obtained is needed again afterwards and may therefore not disappear from the worktable.

```
'macro' 'pointer' min work = 1, max work = 4095.
```

```
'list' work [min work : max work].
```

```
'pointer' pwork.
```

```
'action' insert 2, initialize chain adm, errormessage.
```

```
'macro' 'action'
```

```
get link = '2':= work ['1'], put link = work ['1']:= '2',
```

```
get head = '2':= work ['1' + 1], put head = work ['1' + 1]:= '2',
```

```
get tail = '2':= work ['1' + 2], put tail = work ['1' + 2]:= '2'.
```

```
search 2 + p + hd + tl - q - x:
```

```
make + q + p,
```

```
fnd: (lseq + min work + q, get head + q + x,
```

```
equal + x + hd, get tail + q + tl;
```

```
get link + q + q, :fnd)).
```

```
insert 2 + p + hd + tl - q:
```

```
make + q + p, make + p + pwork, add + pwork + 3 + pwork,
```

```
(less + max work + pwork, errormessage + work full;
```

```
put head + p + hd, put tail + p + tail, put link + p + q).
```

```
already there + p + hd + tl + q - x:
```

```
make + q + p,
```

```
fnd: (lseq + min work + q,
```

```
(get head + q + x, equal + x + hd, get tail + q + x,
```

```
equal + x + tl;
```

```
get link + q + q, :fnd)).
```

```

detach 2 + p + hd + tl:
    lseq + min work + p, get head + p + hd, get tail + p + tail,
        get link + p + p.
'pointer' access.
initialize chain adm:
    make + pwork + min work, make + access + 0.

```

After one call of *initialize chain adm*, *access* gives access to an empty list. If the worktable becomes too small, *errormessage* is called (see 3.4).

If the compiler is to be run on a computer where the wordlength is such that the head and tail may fit together in one word, then one may simply replace the macro text of put head, get head, put tail and get tail by a version which manipulates with halfwords.

An advantage of using macros for accessing fields of a datastructure is that one can change the datastructures without invalidating the whole grammar. One could, e.g., in testing the compiler perform no packing, and in a later production version introduce packing. This allows the development of quite machine-independent compilers without a resulting price in efficiency.

As an initialization, *initialize chain adm* has to be called.

Note that work full is a terminal symbol (being an affix without defining occurrence) and that a representation of that symbol will have to be at the beginning of the input to the compiler resulting from this compiler description.

3.3 Reading tags.

We want to read <identifier>s, i.e. tags, and obtain a unique key for each different tag. We want to be able to list the tags in alphabetical order. Accompanying each tag we must remember a pointer, originally zero.

We choose the following storage organization: the information is kept in cells linked together in a text table, one cell for each tag, a cell consisting of a fixed and a variable part.

The fixed part consists of 3 pointer:

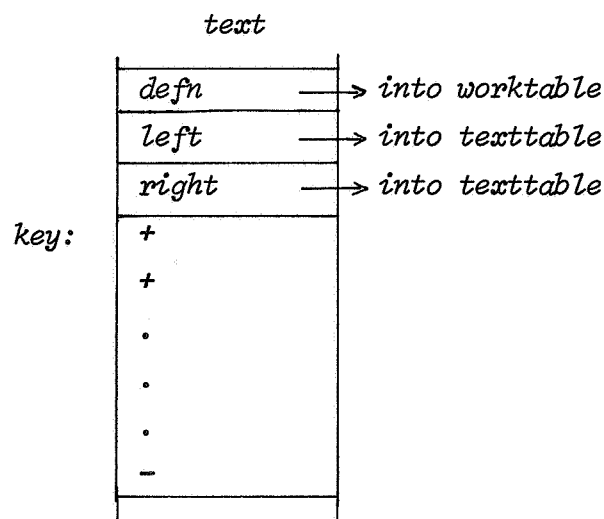
defn an access to a definition chain;
left a pointer to the chain of all tags alphabetically preceding the tag, zero if empty;
right idem for the alphabetically following tags.

The variable part contains the successive characters of the tag, packed three to a word, the last word being filled out with dummies (= 255) and provided with a - sign, the others with a + sign. (We will assume that the wordlength of the computer on which the compiler will run is sufficient, i.e. at least 25 bits.)

We will give out the index (in the text table) of the first word of the variable part as a key.

The chainstructure just described is a doubly linked namelist. It is more efficient to use than a linear namelist, and less efficient than a namelist using hash coding, over which again it has the advantage of keeping the tags in alphabetical order.

In a picture:



When reading a tag, we will first build up a cell for it in the text table, and then enter it, checking whether the same text already occurs in the table; if so, then the key of the old occurrence is delivered and the new cell removed; if not, then the new cell is fitted into the chain.

```

'macro' 'pointer' mintext = 1, max text = 4095.
'list text [min text : max text].
'pointer' ptext.
'macro' 'action'
put defn = text ['1' - 3]:= '2', get defn = '2':= text ['1' - 3],
put left = text ['1' - 2]:= '2', get left = '2':= text ['1' - 2],
put right= text ['1' - 1]:= '2', get right= '2':= text ['1' - 1].
'macro' 'predicate'
    is letter = '1' > 9 ^ '1' < 36,
    is letgit = '1' < 36.

[treatment of text]
'action' add to text, stack text, stack last, enter, reserve fixed part.
'pointer' word, c.
'macro' 'pointer' dummy = 255, bytesize = 256, maxbytes = 3.
'macro' 'predicate' left = abs('1') < abs('2').
add to text + x:
    equal + c + maxbytes, stack text + word, make + word + x,
                                     make + c + 1;
    addmult + word + bytesize + x + word, incr + c.
stack text + x:
    lseq + ptext + max text, put + text + ptext + x, incr + ptext;
    errormessage + text full.
stack last:
    rep: (less + c + maxbytes, addmult + word + bytesize + dummy + word,
          incr + c, :rep;
    mark + word, stack text + word).

```

Now follows the *pièce de résistance* of the treatment of text. The action enter checks whether the topmost element of the texttable, which is pointed to by x, already occurred somewhere else in the texttable; if such is the case, then x is made to point to that older element, and the new element is deleted. Also, this action takes care of the alphabetic ordering of the tags.

```

enter + x + pchain - y - x1 - y1 - wx - wy:
    equal + pchain + 0, make + pchain + x, reserve fixed part;
    make + y + pchain,
nxy: make + x1 + x, make + y1 + y,
nxw: get + text + x1 + wx, get + text + y1 + wy,
    (equal + wx + wy,
      (marked + wx,
        (equal + x + y, reserve fixed part;
          make + ptext + x, make + x + y);
        incr + x1, incr + y1, :nxw);
      left + wx + wy, get left + y + wy,
        (equal + wy + 0, put left + y + x, reserve fixed part;
          make + y + wy, :nxy);
      get right + y + wy,
        (equal + wy + 0, put right + y + x, reserve fixed part;
          make + y + wy, :nxy)).
reserve fixed part:
    stack text + 0, stack text + 0, stack text + 0.

[reading a tag]
'pointer' ptag.
read tag + x - t:
    letter + t, make + word + t, make + c + 1, make + x + ptext,
    nxt: (letgit + t, add to text + t, :nxt;
      stack last, enter + x + ptag).
letter + t:
    is letter + char, make + t + char, nextchar.
letgit + t:
    is letgit + char, make + t + char, nextchar.
initialize text adm:
    make + ptext + min text, reserve fixed part, make + ptag + 0.

```

Of course the left- and right-pointer of a tag can be packed together again in one word. As initialization, *initialize text adm* has to be called. The pointer ptag gives access to the chain of tags, text full is a terminal.

3.4 Printing, errormessages.

We are now in a position to define the printing of tags, and the treatment of errormessages.

```
'action' print, print 1, print 2, nlcr, errormessage.
print + x - p - el:
    make + p + x,
        rep: get + text + p + el,
            (marked + el, mark + el, print 1 + el;
            print 1 + el, incr + p, :rep).
print 1 + t - c:
    prsym + spacechar,
    make + c + max bytes, print 2 + c + t.
print 2 + c + t - el:
    equal + c + 1, prsym + el;
    divrem + t + bytesize + t + el, decr + c, print 2 + c + t,
                                                prsym + el.

nlcr: prsym + newlinechar.
errormessage + x: nlcr, print + x, exit.
```

The definition of print 2 has to be recursive in order to get the characters out in the right order.

The procedure *errormessage* prints a message and then terminates execution. A more subtle reaction might be envisaged.

3.5 An input administration.

Assuming that there also exist texts, similar to that in 3.3, for reading bold symbols, special symbols and constants, a complete input administration can be constructed:

```

'macro' 'pointer' min inpt = 1, max inpt = 2047.
'list' inpt [min inpt : max inpt].
'pointer' pin, xin, symb.
'action' next symbol, skip layout.
next symbol:
    nxt: skip layout,
        (read tag + symb;
        read bold + symb, add + symb + 20000 + symb;
        read spec + symb, add + symb + 30000 + symb;
        read const+ symb, add + symb + 40000 + symb;
        errormessage + incorrect char, nextchar, :nxt).
skip layout:
    skp: (equal + char + new line char, nextchar, :skp;).
rq + x:
    equal + symb + x, incr + pin,
        (equal + pin + xin, next symbol, put + inpt + pin + symb,
                                                incr + xin;
        get + inpt + pin + symb).
identifier + x:
    less + symb + 20000, make + x + symb, rq + x.
initialise for reading:
    nextchar, initialize chain adm, initialize text adm,
    make + xin + min inpt, make + pin + xin, nextsymbol, put + inpt +
                                                pin + symb.

```

The pointer *pin* serves as reading pointer, *xin* indicates up to what index the inputarray has been filled with symbol-keys. In *next symbol*, a multiple of 10000 is added to some keys as a type-indication, allowing discrimination between, e.g., tags and bolds. By means of *rq*, one can ask whether the current symbol (kept in *symb*), is equal to a given symbol, while *identifier* returns true only if the current symbol is a tag. A restoring parser can now work on the basis of this input administration.

3.6 Reading declarers.

We will now show how to recognize declarers not unlike those of ALGOL 68. During recognition of a declarer, some (head, tail) pairs may be added to the worktable (3.2), and the mode of the declarer is represented by the index of one such pair.

Let μ stand for a declarer, $\bar{\mu}$ for its index, τ for a tag, $\bar{\tau}$ for its key.

int gives an index to (0, int)
real gives (0, real)
long μ gives (long, $\bar{\mu}$)
ref μ gives (ref, $\bar{\mu}$) and
struct ($\mu_1\tau_1, \mu_2\tau_2, \dots, \mu_n\tau_n$) gives
 (struct, $\bar{\mu}_1, \bar{\mu}_2, \dots, \bar{\mu}_n, 0$)
 $(\bar{\mu}_1, \bar{\tau}_1) (\bar{\mu}_2, \bar{\tau}_2) \dots (\bar{\mu}_n, \bar{\tau}_n)$

In adding a pair to the worktable, we take care that no two copies of a pair are ever inserted: in adding a pair that was already there, the index of the old copy is obtained. This implies that equivalent modes automatically get equal indices.

'action' add to decl.

'pointer' pdecl.

declarer + mode:

primitive declarator + mode;

long declarator + mode;

ref declarator + mode;

struct declarator + mode.

'restore'

primitive declarator + mode:

rq + int, add to decl + 0 + int + mode;

rq + real, add to decl + 0 + real + mode.

```

long declarator + mode:
    rq + long, declarer + mode, add to decl + long + mode + mode.
ref declarator + mode:
    rq + ref, declarer + mode, add to decl + ref + mode + mode.
struct declarator + mode:
    rq + struct, rq + open, fields + mode, add to decl + struct +
                                                mode + mode.

fields + mode - m1 - m2:
    rq + close, make + mode + 0;
    field + m1, fields + m2, add to decl + m1 + m2 + mode.
field + mode - tag:
    declarer + mode, identifier + tag, add to decl + mode + tag + mode.

'unrestore'
add to decl + hd + tl + mode - oldp:
    already there + hd + tl + pdecl + oldp, make + mode + oldp;
    make + oldp + pwork, insert 2 + hd + tl + pdecl, make + mode +
                                                oldp.

```

Note that for the declarators we want restoring parsers. Also note that affixes that have no defining occurrence, e.g., `int`, `real`, `long`, etc., are terminals, so that a representation for them has to be given as input to the resulting compiler.

For `already there` and `insert 2`, see 3.2.

3.7 Collecting defining occurrences.

We will now show how to collect defining occurrences of identifiers. We assume the existence of a global pointer *blocknumber* that is automatically updated by some other part of the compiler description.

A defining occurrence of an identifier is an occurrence in a declaration. For each defining occurrence we want to store its *blocknumber* and *mode* in the worktable. We will not allow an identifier to be defined twice in one block.

```

'global' 'pointer' blocknumber.
'action' define identifier
declaration - mode - idf:
    declarer + mode, identifier + idf,
    rep: (define identifier + idf + mode,
          (rq + comma,
            (identifier + idf, :rep;
             declarer + mode, identifier + idf, :rep;
             errormessage + incorrect declaration);
            rq + semicolon)).
define identifier + idf + mode - def - dummy:
    get defn + idf + def,
    (search 2 + def + blocknumber + dummy, errormessage +
                                     defined twice;
     insert 2 + def + blocknumber + mode, put defn + idf + def).

```

A more elaborate version of this last example can be found in [2].

COMPILER CCMPILER TOWARDS ALGOL60

```

1
2
3
4 [4 COMPILER COMPILER DESCRIBED IN CDL *****]
5
6 'SHORT'
7
8 [4.1 GENERAL ENVIRONMENT]
9
10 [4.1.1 INTERFACE WITH MACHINE]
11 'EXTERNAL' 'ACTION'
12 NEW PAGE, EXIT, PRSYM, CSYM.
13
14 'EXTERNAL' 'POINTER' RESYM.
15
16 'MACRO' 'FLAG'
17 WAS LETTER='1'>9~'1'<36,
18 WAS LETGIT='1'<36,
19 WAS DIGIT='1'<10,
20 WAS SPECCH='1'>63.
21
22 'MACRO' 'POINTER'
23 NIX=63, MINUSCODE=65, SPACECODE=93, TABCODE=118, NLRCODE=119, ACCC=120, QUOTE=121.
24
25 [4.1.2 STACKS]
26 'MACRO' 'POINTER'
27 MIN TAG =100001, MAX TAG =110000,
28 MIN BOLD=200001, MAX BOLD=200600,
29 MIN SPEC=300001, MAX SPEC=300100,
30 MIN CONS=400001, MAX CONS=400300,
31 MIN LOC =500001, MAX LOC =500200,
32 MIN GLOB=600001, MAX GLOB=601000,
33 MIN MACR=700001, MAX MACR=701100,
34 MIN TEXT=800001, MAX TEXT=804000.
35
36 'LIST'
37 TTAG [MIN TAG :MAX TAG ],
38 TROLL [MIN BOLD:MAX ROLD],
39 TSPEC [MIN SPEC:MAX SPEC],
40 TCONS [MIN CONS:MAX CONS],
41 LLOC [MIN LOC :MAX LOC ],
42 LLOB [MIN GLOB:MAX GLOB],
43 LTEXT [MIN TEXT:MAX TEXT],
44 LMACR [MIN MACR:MAX MACR].
45
46 'MACRO' 'FLAG'
47 WAS TAG =MIN TAG 'LE'~'1'~'1'<PTAG,
48 WAS BOLD=MIN BOLD 'LE'~'1'~'1'<PBOLD,
49 WAS SPEC=MIN SPEC 'LE'~'1'~'1'<PSPEC,
50 WAS CONS=MIN CONS 'LE'~'1'~'1'<PCONS.

```

```

46 [4.1.3 MACROS]
47 'MACRO' 'ACTION'
48 GET '3':='1'['2'],
49 PUT '1'['2']:= '3',
50 MAKE='1':='2',
51 SET = '1':='2',
52 MARK='1':='1',
53 ADD='3':='1'+ '2',
54 SUBTR='3':='1'- '2',
55 ARDMULT='4':='1'* '2'+ '3',
56 D VREM='3':='1'/"'2'+ '4':='1'- '2'* '3',
57 PACK3='4':='( '1'*128+'2)*128+'3',
58 UNPACK3='2':='1'/"'16384; '4':='1'-16384+'2'; '3':='4'/"'128; '4':='4'-128+'3',
59 PACK2='3':='8192*'1'+ '2',
60 UNPACK2='2':='1'/"'8192; '3':='1'-8192*'2',
61 GET TAIL='2':='1'- '1'/"'8192*8192,
62 INCR = '1':='1'+1,
63 DECR = '1':='1'-1.
64
65 'MACRO' 'FLAG'
66 MARKED='1'<0,
67 LEFT=ABS('1')<ABS('2'),
68 LESS='1'<'2',
69 LSEQ='1'<='2',
70 EQUAL='1'='2',
71
71 'MACRO' 'POINTER'
72 FALSE='FALSE',
73 TRUE='TRUE',
74
74 [DATA STRUCTURES AND THEIR ACCESS]
75 'MACRO' 'ACTION'
76 GET NPARS='2':='TTAG['1'-8], PUT NPARS=TTAG['1'-8]:'2',
77 GET MTEXT='2':='TTAG['1'-7], PUT MTEXT=TTAG['1'-7]:'2',
78 GET PLACE='2':='TTAG['1'-6], PUT PLACE=TTAG['1'-6]:'2',
79 GET STATE='2':='TTAG['1'-5], PUT STATE=TTAG['1'-5]:'2',
80 GET SORT = '2':='TTAG['1'-4], PUT SORT =TTAG['1'-4]:'2',
81 GET TYPE = '2':='TTAG['1'-3], PUT TYPE =TTAG['1'-3]:'2',
82 GET LEFT = '2':='TTAG['1'-2], PUT LEFT =TTAG['1'-2]:'2',
83 GET RIGHT='2':='TTAG['1'-1], PUT RIGHT=TTAG['1'-1]:'2',
84
84 [4.1.4 POINTERS AND FLAGS]
85 'POINTER'
86 PTAG,PBOLD,PSPEC,PCONS,
87 PLOC,PLOB,PTEXT,PMACR,XLOC.
88
89 'FLAG' GIVE TEXT,GIVE TRACE,LEGIBLE.
90 'POINTER' FIRST TAG.
91

```

```
92
93
94
95
96 [4.2 OUTPUT]
97
98 [4.2.1 PRINTER SECTION]
99 'ACTION'OUT,OUTINT,PRINT,PRINT1,NLCR,TAB,SPACE,OUTINT1,PRCHAR,SPACES,TABS,
100 POSITION,SHIFT2LINES.
101 'POINTER'POS.
102 OUT+X-EL:
103 WAS TAG +X,PRINT+TTAG +X;
104 WAS BOLD+X,PRINT+TBOLD+X;
105 WAS SPEC+X,PRINT+TSPEC+X;
106 WAS CONS + X, GET+CONS+X+EL,OUTINT+EL;
107 OUTINT+X.
108 OUTINT+X-QUOT-REM:
109 SPACE,
110 LESS+X+0,MAKE+REM+X,MARK+REM,PRCHAR+MINUSCODE,OUTINT+REM;
111 EQUAL+X+0,PRCHAR+0,SPACE;
112 DIVREM+X+10+QUOT+REM,OUTINT+QUOT,PRCHAR+REM,SPACE.
113 OUTINT+X-QUOT-REM:
114 EQUAL+X+0;
115 DIVREM+X+10+QUOT+REM,OUTINT+QUOT,PRCHAR+REM.
116 PRINT+LIST+Y-X-LX;
117 MAKE+X+Y,SPACE,
118 RST: GET+LIST+X+LX,
119 (MARKED+LX,MARK+LX,PRINT1+LX,PRINT1+LX,INCR+X,RST).
120 PRINT1+X-X1-X2-X3;
121 UNPACK3+X+X1+X2+X3.
122 PRCHAR+X1,PRCHAR+X2,PRCHAR+X3.
123 POSITION+X-TBS-SPCES;
124 SUBTR+X+POS+SPCES,LEQ+0+SPCES,SPACES+SPCES;
125 NLCR,DIVREM+X+8+TBS+SPCES,TABS+TBS,
126 (EQUAL+SPCES+0;SPACES+SPCES).
127 SPACES+X-N;
128 MAKE+N+0,
129 SPC: LESS+N+X,SPACE,INCR+N,:SPC;
130 TABS+X-N;
131 MAKE+N+0,
132 TBS: LESS+N+X,TAB,INCR+N,:TBS;
```

133 NLCR:

23

134 PRSYM +NLRCODE,MAKE+POS+0.

24

135 SHIFT 2 LINES-OLD POS:

24

136 MAKE+OLDPOS+POS,NLCR,NLCR,SPACES+OLD POS.

25

25

137 TAB:PRSYM+TABCODE,ADD+POS+8+POS.

26

138 SPACE:

27

139 PRCHAR+SPACECODE.

140 PRCHAR+X:

27

141 EQUAL+X+NLRCODE,NLCR;

28

142 EQUAL+X+TABCODE,TAB;

28

143 EQUAL+X+NIX;

28

144 PRSYM+X,INCR+POS.

28

28

145

28

146 [4.2.2 TRACE ADMINISTRATION]

29

147 'ACTION',TRACE,SIGNAL,INFORM,ERROR.

29

148 TRACE+X:

30

149 GIVE TRACE,POSITION+32,INFORM+X;

30

150 SIGNAL+X:

31

151 POSITION+16,INFORM+X.

152 INFORM+X,STATE-SORT-TYPE:

31

153 GET STATE+X+STATE,GET SORT+X+SORT,GET TYPE+X+TYPE,

32

154 OUT+STATE,OUT+SORT,OUT+TYPE,OUT+X.

32

155 ERROR+TEXT+INFO-OLD POS:

32

156 MAKE+OLD POS+POS,NLCR,OUT+TEXT,

33

157 (WAS TAG+INFO,INFORM+INFO,POSITION+OLD POS;

33

158 EQUAL+INFO+0,POSITION+OLD POS;

33

159 OUT+INFO,POSITION+OLD POS).

33

160

33

161 [4.2.3 CARD PUNCH SECTION]

33

162 'ACTION',WRITE,WRITEINT,WRITEINT1,PUNCH,PUNCH1,NEW CARD,TAB CARD,SPACE CARD,

34

163 PUCHAR.

34

164 'MACRO','POINTER';

34

165 MAX CARD=72,DOLLAR CODE=133.

34

166 'POINTER' CARD,CPOS.

34

167 WRITE +X - EL;
168 WAS TAG + X, PUNCH + TTAG + X;
169 WAS BOLD + X, PUNCH + TBOLD + X;
170 WAS SPEC + X, PUNCH + TSPEC + X;
171 WAS CONS + X, GET+TCONS+X+EL, WRITE INT+EL;
172 WRITE INT+X.

173 WRITE INT+X-QUOT-REM;
174 LESS+X+0, MAKE+REM+X, MARK+REM, PUCHAR+MINUSCODE, WRITE INT+REM;
175 EQUAL+X+0, PUCHAR+0;
176 DIVREM+X+10+QUOT+REM, WRITE INT+QUOT, PUCHAR+REM.

177 WRITE INT+X-QUOT-REM;
178 EQUAL+X+0;
179 DIVREM+X+10+QUOT+REM, WRITE INT+QUOT, PUCHAR+REM.

180 PUNCH+LIST+Y-X-LX;
181 MAKE+X+Y,
182 RST: GET+LIST+X+LX,
183 (MARKED+LX, MARK+LX, PUNCH1+LX;
184 PUNCH1+LX, INCR+X, :RST).

185 PUNCH1+X-X1-X2-X3;
186 UNPACK3+X+X1+X2+X3.
187 PUCHAR+X1, PUCHAR+X2, PUCHAR+X3.

188 NEW CARD:
189 SPC: LSEQ +CPOS + MAX CARD, CSYM+SPACECODE, INCR+CPOS, :SPC;
190 MAKE+CPOS+1, WRITE INT+CARD, CSYM+NLCR CODE, INCR+CARD, MAKE+CPOS+1.

191 TAB CARD-SPCES-TBS:
192 DIVREM+CPOS+8+TBS+SPCES,
193 SPC: LESS+SPCES +8, SPACE CARD, INCR+SPCES , :SPC;

194 SPACE CARD:
195 PUCHAR+SPACECODE.

196 PUCHAR+X:
197 EQUAL + X + NLCRCODE, NEW CARD;
198 EQUAL + X + TABCODE, TAB CARD;
199 EQUAL + X + NIX;
200 LESS+CPOS+MAX CARD, CSYM+X, INCR+CPOS;
201 CSYM+DOLLARCODE, CSYM+NLCRCODE, MAKE+CPOS+1, PUCHAR+X.

202
203 [4,2,4 RESULT ADMINISTRATION]
204 ACTION'G, LINE UP, PUTABS, U, L, NEW LINE, BLANK LINE, TAB LINE.

205 FLAG' LINED UP.

46 206 'POINTER' INDENTATION.
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207 G + X:
208 LINE UP, WRITE + X.
209 LINE UP:
210 LINED UP;
211 LEGIBLE, PUTABS + INDENTATION, SET + LINED UP + TRUE.
212 PUTABS + N - N1:
213 LSEQ + N + 0;
214 MAKE + N1 + N, DECR + N1, TAB CARD, PUTABS + N1.
215 U* INCR + INDENTATION.
216 L: DECR + INDENTATION.
217 NEW LINE:
218 LEGIBLE, LINED UP, BLANK LINE.
219 BLANK LINE:
220 NEW CARD, SET + LINED UP + FALSE.
221 TAB LINE: LEGIBLE, TAB CARD.
222
223
224
225
226
227 [4.3 INPUT]
228 'ACTION' POST MORTEM.
229
230 'POINTER'
231 LINE, CHAR, INPT.
232
233 [4.3.1 READING CHARACTERS]
234 'ACTION' NEXT CHAR, DISPLAY CHARACTER, NEXT NON SPACE CHAR.
235 LETTER + X:
236 WAS LETTER + CHAR, MAKE + X + CHAR, NEXT NON SPACE CHAR.
237 LEFT + X:

238 WAS LETGIT+CHAR,MAKE+X+CHAR,NEXT NON SPACE CHAR.

239 DIGIT+X:

240 WAS DIGIT +CHAR,MAKE+X+CHAR,NEXT NON SPACE CHAR.

241 BOLDCHAR+X:

242 EQUAL+CHAR+QUOTE,MAKE+X+ACCC,NEXTCHAR;

243 MAKE+X+CHAR,NEXTCHAR.

244 ACCENT:EQUAL+CHAR+ACCC,NEXTCHAR.

245 SPECCHAR + X:

246 WAS SPECCH+CHAR,MAKE+X+CHAR,NEXT NON SPACE CHAR.

247 NEXTCHAR:

248 DISPLAY CHARACTER,MAKE + CHAR + RESYM.

249 DISPLAY CHARACTER:

250 EQUAL+CHAR+NLCRCODE,INCR+LINE,

251 (GIVE TEXT,POSITION+133,OUTINT1+CARD,

252 POSITION+48,OUTNT1+LINE,POSITION+52;);

253 GIVE TEXT,PRCHAR+CHAR)

254 NEXT NON SPACE CHAR:

255 CHR: NEXTCHAR.

256 (EQUAL+CHAR+SPACECODE,:CHR:).

257 LAYOUT SYMBOL:

258 EQUAL + CHAR + SPACECODE;

259 EQUAL + CHAR + NLCRCODE;

260 EQUAL + CHAR + TABCODE .

261

262 [4.3.2 TAG SYMBOLS]

263 'ACTION'STACK TAG,ENTER TAG,RESERVE ADMIN SPACE.

264

READ TAG + X - T - T1 - T2 - T3 - T4:

LETTER+T1,MAKE+X+PTAG,

NXT: (LETGIT + T2,

(LETGIT + T3,

(LETGIT + T4, PACK3 + T1 + T2 + T3 + T,

STACK TAG+T,MAKE+T1+T4,:NXT;

PACK3 + T1 + T2 + T3 + T, :LST);

PACK3 + T1 + T2 + NIX + T, :LST);

PACK3 + T1 + NIX + T,

LST: MARK + T, STACK TAG + T).

274

STACK TAG + X:

LSEQ + PTAG + MAX TAG, PUT + TTAG + PTAG + X, INCR + PTAG;

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276 ERROR + TAG FULL+0,POST MORTEM,EXIT.

277 ENTER TAG + X - Y - X1 - Y1 - WX - WY;

278 MAKE + Y + FIRST TAG,

279 NXY: MAKE + X1 + X, MAKE + Y1 + Y,

280 NXW: GET+TTAG+X1+WX,GET+TTAG+Y1+WY,

281 (EQUAL + WX + WY,

282 (MARKED + WX,

283 (EQUAL + Y + X,RESERVE ADMIN SPACE;

284 MAKE+PTAG+X,MAKE+X+Y);

285 INCR+X1,INCR+Y1,;NXW);

286 LEFT+WX+WY,GET LEFT+Y+WY,

287 (EQUAL+WY+0,PUT LEFT+Y+X,RESERVE ADMIN SPACE;

288 MAKE+Y+WY,;NXY);

289 GET RIGHT+Y+WY,

290 (EQUAL+WY+0,PUT RIGHT+Y+X,RESERVE ADMIN SPACE;

291 MAKE+Y+WY,;NXY)).

292 RESERVE ADMIN SPACE:

293 ADD+PTAG+8+PTAG,LESS+PTAG+MAX TAG,PUT NPARS+PTAG+999,

294 PUT MTEXT+PTAG+0,PUT PLACE+PTAG+0,PUT STATE+PTAG+0,PUT SORT+PTAG+0,

295 PUT TYPE+PTAG+0,PUT LEFT+PTAG+0,PUT RIGHT+PTAG+0;

296 ERROR+TAG FULL+0,POST MORTEM,EXIT.

297

298 [4.3.3 BOLD SYMBOLS]

299 'ACTION'STACK BOLD,ENTER BOLD.

300 READ BOLD + X - B1-B2-B3-D:

301 ACCENT, MAKE+X+PBOLD,

302 (ACCENT,PACK3+NIX+NIX+D,MARK+D,STACK BOLD+D;

303 BOLDCHAR+B1,

304 RST: (ACCENT,PACK3+B1+NIX+NIX+D,MARK+D,STACK BOLD+D;

305 BOLDCHAR+B2,

306 (ACCENT,PACK3+B1+B2+NIX+D,MARK+D,STACK BOLD+D;

307 BOLDCHAR+B3,PACK3+B1+B2+B3+D,

308 (ACCENT,MARK+D,STACK BOLD+D;

309 BOLDCHAR+B1,STACK BOLD+D,;RST))).

310 STACK BOLD + X:

311 LSEQ + PBOLD + MAX BOLD, PUT + TBOLD + PBOLD + X, INCR+PBOLD;

312 ERROR + BOLD FULL+0,POST MORTEM,EXIT.

313

314 ENTER BOLD + X - Y - X1 - Y1 - WX - WY:

315 MAKE+Y+MIN BOLD,

316 NXY: MAKE + X1 + X, MAKE + Y1 + Y,

317 NXW: GET+TBOLD+X1+WX,GET+TBOLD+Y1+WY,

318 (EQUAL + WX + WY,

319 (MARKED + WX,

320 (EQUAL+Y+X;

321 MAKE+PBOLD+X,MAKE+X+Y);

322 INCR+X1,INCR+Y1,;NXW);

323 SKP: MARKED + WY, MAKE + Y + Y1, INCR + Y, ;NXY;

INCR + Y1, GET + TBOLD + Y1 + WY, ;SKP).

324
325 [4.3.4 SPECIAL SYMBOLS]
326 'ACTION' STACK SPEC, ENTER SPEC.

327 READ SPEC + X - S - S1;
328 SPECCHAR+S1, MAKE+X+PSPEC, PACK3+S1+NIX+NIX+S, MARK+S, STACK SPEC+S.

329 STACK SPEC + X;
330 LSEQ + PSPEC + MAX SPEC, PUT + TSPEC + PSPEC + X, INCR+PSPEC;
331 ERROR + SPEC FULL+0, POST MORTEM, EXIT.

332 ENTER SPEC + X - Y - WX - WY;
333 MAKE+Y+MIN SPEC, GET+TSPEC+X+WX,
334 NXY: GET+TSPEC+Y+WY,
335 (EQUAL + WX + WY,
336 (EQUAL+Y+X;
337 MAKE+PSPEC+X, MAKE+X+Y);
338 INCR + Y, :NXY).

339
340 [4.3.5 CONSTANTS]
341 'ACTION' STACK CONS, ENTER CONS.

342 READ CONS + X - D - S;
343 DIGIT + D, MAKE + S + D,
344 RST: (DIGIT + D, ADDMULT + S + 10 + D + S, :RST;
345 MAKE + X + PCONS, STACK CONS + S).

346 STACK CONS + X;
347 LSEQ + PCONS + MAX CONS, PUT + TCONS + PCONS + X, INCR + PCONS;
348 ERROR + CONS FULL+0, POST MORTEM, EXIT.

349 ENTER CONS + X - Y - S - T;
350 GET + TCONS + X + S, MAKE + Y + MIN CONS,
351 NXT: GET + TCONS + Y + T,
352 (EQUAL + S + T,
353 (EQUAL + X + Y;
354 MAKE + PCONS + X, MAKE + X + Y);
355 INCR + Y, :NXT).

356
357 [4.3.6 INPUT ADMINISTRATION]
358 'ACTION' INITIALIZE FOR READING, NEXT SYMBOL, READ.

359 INITIALIZE FOR READING;
360 MAKE+GIVE TEXT+FALSE, MAKE+PCONS+MIN CONS,
361 MAKE+PTAG+MIN TAG, MAKE+PBOLD+MIN BOLD, MAKE +PSPEC+MIN SPEC, MAKE+LINE+0,
362 MAKE+CHAR+RESYM, RESERVE ADMIN SPACE, MAKE+FIRST TAG+PTAG.

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363 IS TAG + X: WAS TAG + INPT, MAKE + X + INPT, NEXT SYMBOL.

364 IS BOLD+X: WAS BOLD+INPT, MAKE+X+INPT, NEXT SYMBOL.

365 IS CONS + X: WAS CONS + INPT, MAKE + X + INPT, NEXT SYMBOL.

366 R + X: AHEAD + X, NEXT SYMBOL.

367 AHEAD + X: EQUAL + INPT + X.

368 NEXT SYMBOL:

369 READ+INPT.

370 READ+X:

371 SKP: LAYOUT SYMBOL, NEXTCHAR, SKP;

372 READ TAG + X, ENTER TAG+X;

373 READ BOLD+X, ENTER BOLD+X;

374 READ SPEC+X, ENTER SPEC+X;

375 READ CONS+X, ENTER CONS+X;

376 ERROR+WRONG INPUT CODE+CHAR, NEXTCHAR, READ+X.

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382 [4.4 COMPILER COMPILER TO ALGOL60]

383 'FLAG' GIVE RSTO, ADVANCED.

384 'POINTER' HANDLE, BOUNDX, SPECX, TO NEXT, TO LOST, TO END, CUR LAB, NEW LAB.

385

386 [4.4.1 AUXILIARY ACTIONS]

387 'ACTION'

388 FORGET LOCALS, WARN, ENTER LOCAL, ENTER GLOBAL, ENTER MACRO, ENTER TEXT,

389 ADD PLACE, REPLACE.

390 ADD PLACE+TAG-H-K-Q:

391 GET PLACE+TAG+H,

392 (EQUAL + H + 0, REPLACE + TAG + H;

393 GET+LTEXT+H+Q, GET TAIL+Q+K, LESS+K+LINE, REPLACE+TAG+H;).

394 REPLACE + TAG + H:

395 PUT PLACE+TAG+PTEXT,

396 (EQUAL+H+0, ENTER TEXT+LINE;

397 SUBTR+H+MIN TEXT+H, INCR+H, PACK2+H+LINE+H, ENTER TEXT+H).

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398 FORGET LOCALS-LOC:
399 RST: EQUAL+PLOC-MIN LOC:
400 : DECR+PLOC,GET+LLOC+PLOC+LOC,WARN+LOC,
401 : PUT STATE+LOC+0,PUT SORT+LOC+0,PUT TYPE+LOC+0,IRST.
101
102
103
402 WARN+X=STATE:
403 GET STATE+X=STATE,
404 (EQUAL+STATE+BLANK,TRACE+X;
405 EQUAL+STATE+0,TRACE+X;
406 SIGNAL+X).
101
102
103
407 ENTER LOCAL + HEAD:
408 LSEQ + PLOC + MAX LOC, PUT + LLOC + PLOC + HEAD, INCR + PLOC,
409 (LSEQ + PLOC + XLOC; MAKE + XLOC + PLOC);
410 ERROR + LOC FULL+0,POST MORTEM,EXIT.
102
103
104
411 ENTER GLOBAL + HEAD:
412 LSEQ + PGLOB + MAX GLOB, PUT + LGLOB + PGLOB + HEAD, INCR + PGLOB;
413 ERROR + GLOB FULL+0,POST MORTEM,EXIT.
103
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105
414 ENTER MACRO + HEAD:
415 LSEQ + PMACR + MAX MACR, PUT + LMACR + PMACR + HEAD, INCR + PMACR;
416 ERROR + MACR FULL+0,POST MORTEM,EXIT.
104
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106
417 ENTER TEXT + X:
418 LSEQ + PTEXT + MAX TEXT, PUT + LTEXT + PTEXT + X, INCR + PTEXT;
419 ERROR + TEXT FULL+0,POST MORTEM,EXIT.
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420
421 [4.4.2 TREATMENT OF TYPES]
422 'ACTION'
423 DEFINE ACTION,DEFINE LABEL,DEFINE PREDICATE,DEFINE AFFIX,APPLY ACTION,
424 APPLY LABEL,REDEFINE,NEW PLACE,NEW NPARS,APPLY,APPLY PREDICATE.
106
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108
425 DEFINE ACTION+HEAD:
426 REDEFINE+HEAD+DEFINED+GLOBAL+ACTION.
107
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427 DEFINE PREDICATE+HEAD:
428 REDEFINE+HEAD+DEFINED+GLOBAL+PREDICATE.
108
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110
429 DEFINE AFFIX+HEAD:
430 REDEFINE+HEAD+DEFINED+LOCAL+POINTER.
109
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111
431 DEFINE LABEL+LAB:
432 WAS TAG +LAB,REDEFINE+LAB+DEFINED+LOCAL+LABEL;
433 WAS CONS+LAB.
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112
434 APPLY ACTION+HEAD:
435 REDEFINE+HEAD+APPLIED+GLOBAL+ACTION.
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112 436 APPLY PREDICATE+HEAD:
113 437 REDEFINE+HEAD+APPLIED+GLOBAL+PREDICATE.

113 438 APPLY LABEL+LAB:
114 439 WAS TAG +LAB,REDEFINE+LAB+APPLIED+LOCAL+LABEL;
114 440 WAS CONS+LAB.

114 441 APPLY+X:
115 442 NEW PLACE+X,TRY STATE+X+APPLIED.

115 443 REDEFINE+X+NSTATE+NSORT+NTYPE+OLD POS:
116 444 REDEFINE1+X+NSTATE+NSORT+NTYPE,NEW PLACE+X;
116 445 MAKE+OLD POS +POS,NLCR,OUT+X IMPOSSIBLE,INFORM+X,
116 446 OUT+NSTATE,OUT+NSORT,OUT+NTYPE,POSITION+OLD POS.

116 447 REDEFINE1+X+NSTATE+NSORT+NTYPE:
117 448 TRY TYPE+X+NTYPE,TRY SORT+X+NSORT, TRY STATE+X+NSTATE.

BACKTRACK?
BACKTRACK?

117 449 TRY STATE+X+S-Q:
118 450 GET STATE+X+Q,EQUAL+Q+0,PUT STATE+X+S;
118 451 EQUAL+Q+DEFINED,EQUAL+S+APPLIED,PUT STATE+X+BLANK;
118 452 EQUAL+Q+APPLIED,
118 453 (EQUAL+S+APPLIED;
118 454 EQUAL+S+DEFINED,PUT STATE+X+BLANK);
118 455 EQUAL+Q+BLANK ,EQUAL+S+APPLIED.

118 456 TRY SORT+X+P-Q:
119 457 GET SORT+X+Q,EQUAL+Q+0,PUT SORT+X+P,
119 458 (EQUAL+P+LOCAL,ENTER LOCAL+X;
119 459 EQUAL+P+GLOBAL,ENTER GLOBAL+X;
119 460 EQUAL+P+MACRO,ENTER GLOBAL+X);
119 461 EQUAL+P+Q.

119 462 TRY TYPE+X+P-Q:
120 463 GET TYPE+X+Q,EQUAL+Q+0,PUT TYPE+X+P;
120 464 EQUAL+P+Q.

120 465 NEW PLACE+X-SORT:
121 466 GET SORT+X+SORT,EQUAL+SORT+LOCAL; ADD PLACE+X.

121 467 NEW NPARS+X+P-Q:
122 468 GET NPARS+X+Q,EQUAL+Q+P;
122 469 EQUAL+Q+999,PUT NPARS+X+P;
122 470 ERROR+WRONG NUMBER OF PARAMETERS+X.

122 471 WAS ACTION+HEAD-T:
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472 GET TYPE+HEAD+T,EQUAL+T+ACTION.

473 WAS AFFIX+HEAD-T:GET TYPE+HEAD+T, 123
474 EQUAL+T+POINTER; 124
475 EQUAL+T+FLAG; 124
476 EQUAL+T+LIST. 124

477 WAS MACRO+HEAD-S: 124
478 GET SORT+HEAD+S,EQUAL+S+MACRO. 125

479 WAS PARAMLESS MACRO+HEAD: 125
480 WAS TAG+HEAD,WAS MACRO+HEAD,WAS AFFIX+HEAD. 126

BACKTRACK?

481 126
482 [4.4.3 SPECIFICATIONS] 126
483 'ACTION'TREAT SPEC LIST,READ MACRO. 127

484 SPECIFICATION: 127
485 EXTERNAL SPECIFICATION; 128
486 INTERNAL SPECIFICATION; 128
487 MACRO SPECIFICATION. 128

488 EXTERNAL SPECIFICATION-TYPE: 128
489 R+EXTERNAL,IS TYPE+TYPE, 129
TREAT SPEC LIST+DEFINED+GLOBAL+TYPE. 129

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490 IS TYPE+X:MAKE+X+INPT, 129
491 R+ACTION; 130
492 R+PREDICATE; 130
493 R+POINTER; 130
494 R+FLAG. 130

495 TREAT SPEC LIST-STATE+SORT+TYPE-HEAD: 130
496 NXT: IS TAG+HEAD,REDEFINE+HEAD+STATE+SORT+TYPE, 131
497 (R+COMMA:NXT;SHIFT2LINES,R+POINT). 131

498 INTERNAL SPECIFICATION-TYPE: 131
499 IS TYPE OF LOCAL+TYPE,TREAT SPEC LIST+APPLIED+GLOBAL+TYPE. 132

500 IS TYPE OF LOCAL+X:MAKE+X+INPT, 132
501 R+ACTION; 133
502 R+PREDICATE. 133

503 MACRO SPECIFICATION-TYPE-HEAD: 133
504 R+MACRO,IS TYPE+TYPE, 134

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505 NXT: IS TAG+HEAD, 134
REDEFINE+HEAD+DEFINED+MACRO+TYPE,READ MACRO+HEAD. 134

BACKTRACK?

506 (R+COMMA,:NXT;SHIFT2LINES,R+POINT),

507 READ MACRO+HEAD-X:

508 PUT MTEXT+HEAD+PMACR,

509 NXT: (AHEAD + POINT, ENTER MACRO + POINT;

510 AHEAD + COMMA, ENTER MACRO + POINT;

511 IS TAG+X,TRY STATE+X+APPLIED,

ENTER MACRO+X,:NXT;

512 ENTER MACRO + INPT, NEXT SYMBOL, :NXT).

513

514 [4.4.4 DECLARATIONS]

515 'ACTION'TREAT DECL LIST,CONSTANT TEXT.

516 'ACTION'PUT PARAMLESS MACRO.

517 DECLARATION:

518 POINTER DECLARATION;

519 FLAG DECLARATION;

520 LIST DECLARATION.

521 POINTER DECLARATION:

522 R+POINTER,TREAT DECL LIST+Q INTEGER+POINTER.

523 TREAT DECL LIST+ALGOLTYPE+TYPE-HEAD:

524 BLANK LINE,G+ALGOLTYPE,

525 NXT: IS TAG+HEAD,REDEFINE+HEAD+DEFINED+GLOBAL+TYPE,G+HEAD,

(R+COMMA,G+COMMA,:NXT;

527 SHIFT2LINES,R+POINT,G+SEMICOLON).

528 FLAG DECLARATION:

529 R+FLAG,TREAT DECL LIST+Q BOOLEAN+FLAG.

530 LIST DECLARATION-HEAD:

531 R+LIST,BLANK LINE,G+Q INTEGER,G+Q ARRAY,U,

532 NXT: IS TAG+HEAD,

G+HEAD,REDEFINE+HEAD+DEFINED+GLOBAL+LIST,

533 R+SUB,

G+SUB,CONSTANT TEXT,

534 R+COLON,

G+COLON,CONSTANT TEXT,

535 R+BUS,

G+BUS,

(R+COMMA,G+COMMA,NEW LINE,:NXT;

537 SHIFT2LINES,R+POINT,G+SEMICOLON,L).

538 CONSTANT TEXT - SYMB:

539. NXT: IS TAG+ SYMB,APPLY+SYMB,

540 (WAS PARAMLESS MACRO+SYMB,PUT PARAMLESS MACRO+SYMB,:NXT;

541 G + SYMB, :NXT);

542 IS CONS + SYMB, G + SYMB, :NXT;

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BACKTRACK?

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543 R + PLUS, G + PLUS, :NXT; 142
544 R + MINUS, G + MINUS, :NXT; 142
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145 [4.4.5 AFFIX EXPRESSIONS]
546 ACTION'PUTAFFIX EXPRESSION,PUT DIRECT,AFFIX PACK OPTION,TREAT AFFIX,PUT SINGLE, 143
547 PUT PARAMLESS MACRO,PUT MACRO,GET PARAMETERS,GET PAR,PUT PAR. 143
548
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145
549 PUT AFFIX EXPRESSION+HEAD: 143
550 TRY SORT+HEAD+GLOBAL,PUT DIRECT+HEAD; 144
551 TRY SORT+HEAD+MACRO, 144
552 (WAS ACTION+HEAD,PUT MACRO+HEAD; 144
553 G+OPEN,PUT MACRO+HEAD,G+CLOSE). 144
144
145
554 PUT DIRECT+HEAD-NPARS: 144
555 MAKE+NPARS+0,G+HEAD,AFFIX PACK OPTION+NPARS,NEW NPARS+HEAD+NPARS. 145
145
556 AFFIX PACK OPTION+NPARS: 145
557 R + PLUS, G + OPEN, TREAT AFFIX, 146
558 RST: (INCR+NPARS,R+PLUS,G+COMMA,TREAT AFFIX,:RST; 146
559 G + CLOSE); 146
146
560 TREAT AFFIX - AFFIX: 146
561 IS TAG + AFFIX, PUT SINGLE + AFFIX,APPLY+AFFIX; 147
562 IS BOLD+AFFIX,G+STRINGQUOTE,G+AFFIX,G+STRINGQUOTE; 147
563 IS CONS + AFFIX, G + AFFIX; 147
564 ERROR+WRONG AFFIX+AFFIX. 147
147
565 PUT SINGLE+X: 147
566 WAS AFFIX+X, 148
567 (WAS MACRO+X,PUT PARAMLESS MACRO+X; G+X); 148
568 REDEFINE+X+APPLIED+GLOBAL+POINTER,G+X. 148
148
569 PUT PARAMLESS MACRO+Y-X-M: 148
570 GET MTEXT+Y+X,NEW NPARS+Y+0, 149
571 NXT: INCR + X, GET + LMACR + X + M, 149
572 (EQUAL + M + POINT; G + M, :NXT). 149
149
573 PUT MACRO+X+P1-P2-P3-P4-P5-Q-M-NPARS: 149
574 MAKE+NPARS+0,GET MTEXT+X+G, 150
575 GET PARAMETERS+NPARS+P1+P2+P3+P4+P5, NEW NPARS+X+NPARS, 150
576 NXT: INCR + Q, GET + LMACR + Q + M, 150
577 (EQUAL + M + POINT; 150
578 EQUAL+M+ONE ,PUT PAR+P1+1,:NXT; 150
579 EQUAL+M+TWO ,PUT PAR+P2+2,:NXT; 150
580 EQUAL+M+THREE,PUT PAR+P3+3,:NXT; 150
581 EQUAL+M+FOUR ,PUT PAR+P4+4,:NXT; 150
582 EQUAL+M+FIVE ,PUT PAR+P5+5,:NXT; 150
583 WAS PARAMLESS MACRO+M,PUT PARAMLESS MACRO+M,:NXT; 150
584 G + M, :NXT). 150
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585 GET PARAMETERS+P+P1+P2+P3+P4+P5: 151
 586 GET PAR+P+P1,GET PAR+P+P2,GET PAR+P+P3,GET PAR+P+P4,GET PAR+P+P5. 152

587 GET PAR+P+X: 152

588 R+PLUS, INCR+P, 152
 589 (IS TAG+X,APPLY+X; IS CONS+X); 152
 590 MAKE + X + X PARAMETER ERROR.

591 PUT PAR + X + Y: 152

592 WAS CONS+X,G+X; 153
 593 EQUAL+X+X PARAMETER ERROR,ERROR+X+Y; 153
 594 WAS BOLD+X,G+STRINGQUOTE,G+X,G+STRINGQUOTE; 153
 595 PUT SINGLE + X.

596 153

597 [4.4.6 BUILDING STONES OF A RULE] 153

598 'ACTION' PUT CONNECT AND,PUT CONNECT OR,END JUMP,END LABEL,LOST LABEL,PUT INIT, 154
 599 PUT RESTORE,FRESH LABEL,TRUE RESULT,FALSE RESULT,PUT JUMP,PUT LABEL,SEMICOL, 154
 600 BEGIN,ENDER,BECOMES.

601 PUT CONNECT AND: 154

602 GIVE RSTO,PUT JUMP+CUR LAB,INCR+TO NEXT; 155
 603 ADVANCED,ERROR+POSSIBLY BACKTRACK NECESSARY+0,PUT JUMP+0,INCR+TO LOST; 155
 604 PUT JUMP+CUR LAB,INCR+TO NEXT.

605 PUT CONNECT OR: 155

606 AHEAD+SEMICOLON,PUT RESTORE,FRESH LABEL, 156
 607 (EQUAL+TO NEXT+0,ERROR+ALTERNATIVE NEVER REACHED+0;) 156
 608 AHEAD+CLOSE, 156
 609 (EQUAL+TO NEXT+0; 156
 610 PUT LABEL+CUR LAB,FRESH LABEL, 156
 611 (GIVE RSTO;PUT JUMP+0,INCR+TO LOST)); 156
 612 PUT RESTORE,LOST LABEL.

613 LOST LABEL: 156

614 EQUAL+TO LOST+0, 157
 615 (EQUAL+TO NEXT+0;FALSE RESULT); 157
 616 PUT LABEL+0,FALSE RESULT.

617 END LABEL: 157

618 EQUAL+TO END+0;PUT LABEL+999. 158

619 END JUMP: 158

620 EQUAL+INPT+POINT,EQUAL+TO LOST+0,EQUAL+TO NEXT+0; 159
 621 PUT JUMP+999,INCR+TO END.

622 PUT INIT: 159

623 GIVE RSTO,G+INIT RESTORE; 160

624 PUT RESTORE: 160

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625 EQUAL+TO NEXT+0;
626 PUT LABEL+CUR LAB,
627 (GIVE RSTO,G+DO RESTORE; ).
161
162
628 FRESH LABEL:
629 MAKE+CUR LAB+NEW LAB, INCR+NEW LAB.
161
162
630 TRUE RESULT:
631 WAS ACTION+HANDLE;
632 G+HANDLE,BECOMES,G+Q TRUE,SEMICOL.
163
164
633 FALSE RESULT:
634 WAS ACTION+HANDLE;
635 G+HANDLE,BECOMES,G+Q FALSE,SEMICOL.
163
164
636 PUT JUMP + LAB:
637 G+Q GOTO, G + LAB, G + SEMICOLON, NEW LINE, APPLY LABEL+LAB.
164
165
638 PUT LABEL + LAB:
639 NEW LINE,L,G+LAB,G+COLON,U,DEFINE LABEL+LAB,TAB LINE.
165
166
640 SEMICOL: G+SEMICOLON.
166
641 BECOMES:
642 G+COLON,G+EQUALS.
167
168
643 BEGIN: NEW LINE, G + Q BEGIN,U,TAB LINE.
168
644 ENDB: NEW LINE, L, G + Q END.
169
645
646 [4.4.7 GENERAL FORM OF A RULE]
647 'ACTION'OPTIONAL BOUND AFFIXES,REST BOUND AFFIXES,OPTIONAL FREE AFFIXES,RESTLHS.
170
171
648 RULE:
649 LEFT HAND SIDE,MIDDLE, RIGHT HAND SIDE,
172
650 FORGET LOCALS,SHIFT2LINES,R+POINT.
172
651 LEFT HAND SIDE:
652 IS TAG+HANDLE,BLANK LINE,MAKE+BOUNDX+0,MAKE+SPECX+0,
173
653 (WAS ACTION+HANDLE,DEFINE ACTION+HANDLE,REST LHS;
173
654 G+Q BOOLEAN,DEFINE PREDICATE+HANDLE,REST LHS).
173
655 REST LHS:
173
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BACKTRACK?
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BACKTRACK?

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656 G+G. PROCEDURE,G+HANDLE,OPTIONAL BOUND AFFIXES,NEW NPARS+HANDLE+BOUNDX, 174
657 SEMICOL. 175
175 175
658 OPTIONAL BOUND AFFIXES-AFFX: 175
659 BOUND AFFIX1+AFFX,G+OPEN,G+AFFX,REST BOUND AFFIXES; 176
660 BOUND AFFIX2+AFFX,G+OPEN,G+AFFX,REST BOUND AFFIXES,G+AFFX; 176
176 176
661 REST BOUND AFFIXES-AFFX: 175
662 BOUND AFFIX1+AFFX,G+COMMA,G+AFFX,REST BOUND AFFIXES; 176
663 BOUND AFFIX2+AFFX,G+COMMA,G+AFFX,REST BOUND AFFIXES,G+AFFX,DECR+SPECX, 176
664 (EQUAL+SPECX+0;G+COMMA); 176
665 G+CLOSE, 176
666 (EQUAL+SPECX+0;SEMICOL,G+Q INTEGER). 176
176 177
667 BOUND AFFIX1+AFFX: 176
668 R+TIMES,IS TAG+AFFX, DEFINE AFFIX+AFFX,INCR+BOUNDX. 177
177 178
BACKTRACK?
669 BOUND AFFIX2+AFFX: 177
670 R+PLUS,IS TAG+AFFX, DEFINE AFFIX+AFFX,INCR+BOUNDX,INCR+SPECX. 178
178 179
671 MIDDLE: 178
672 BEGIN, OPTIONAL FREE AFFIXES, R + COLON,PUT INIT. 179
179 180
673 OPTIONAL FREE AFFIXES - AFFX: 179
674 R + MINUS, G + Q INTEGER, 180
675 FFX: IS TAG + AFFX, G + AFFX, DEFINE AFFIX + AFFX, 180
676 (R + MINUS, G + COMMA, :FFX; 180
677 G + SEMICOLON); 180
180 181
678 RIGHT HAND SIDE: 180
679 MAKE+CUR LAB+1,MAKE+NEW LAB+2,MAKE+TO LOST+0,MAKE+TO END+0, 181
680 NXT: ALTERNATIVE,PUT CONNECT OR, 181
681 (R+SEMICOLON,:NXT; 181
682 END LABEL,END BR,SEMICOL). 181
181 182
683 ALTERNATIVE-LAB-OLD NEXT: 181
684 MAKE+TO NEXT+0,SET+ADVANCED+FALSE, 182
685 NXT: R+COMMA,:NXT; 182
686 R+COLON,IS TAG+LAB, PUT JUMP+LAB; 182
687 R+OPEN,MAKE+OLD NEXT+TO NEXT,CHOICE, R+CLOSE, 182
182 182
688 (GIVE RSTO,MAKE+TO NEXT+1;MAKE+TO NEXT+OLD NEXT); 182
689 MEMBER,:NXT; 182
690 TRUE RESULT,END JUMP. 182
182 183
691 CHOICE-LAB: 182
183 183

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692 MAKE+LAB+CUR.LAB,BEGIN,PUT INIT,FRESH LABEL,
 693 NXT: ALTERNATIVE,PUT CONNECT OR,
 694 (R+SEMICOLON;NXT;ENDBR,SEMICOL,MAKE+CUR LAB+LAB).

695 MEMBER-HEAD:

696 IS TAG+HEAD,APPLY+HEAD,
 697 (R+COLON,PUT LABEL+HEAD;
 698 TRY TYPE+HEAD+PREDICATE,
 699 G+GIF,G+NOT,PUT AFFIX EXPRESSION+HEAD,G+Q THEN,
 700 PUT CONNECT AND,SET+ADVANCED+TRUE;
 701 TRY TYPE+HEAD+FLAG,
 702 G+GIF,G+NOT,PUT AFFIX EXPRESSION+HEAD,G+Q THEN,
 703 PUT CONNECT AND;
 704 TRY TYPE+HEAD+ACTION,
 705 PUT AFFIX EXPRESSION+HEAD,SEMICOL).

706
 707 [4.4.8 OTHER BUILDING STONES OF A GRAMMAR]
 708 ACTION SKIP UNTIL POINT,TERMINALS.

709 COMMAND:

710 R+RSTO ON,SET+GIVE RSTO+TRUE;
 711 R+RSTO OFF,SET+GIVE RSTO+FALSE;
 712 R+SHORT,MAKE+LEGIBLE+FALSE;
 713 R+LONG,MAKE+LEGIBLE+TRUE;
 714 R + TRACE ON, SET + GIVE TRACE + TRUE;
 715 R + TRACE OFF, SET + GIVE TRACE + FALSE.

716 COMMENT:

717 R+SUB,BLANK LINE,G+Q COMMENT,
 718 RST: (R + BUS, G + SEMICOLON;
 719 G+INPT,READ+INPT,:RST).

720
 721 STARTING SYMBOL-HEAD:
 722 R+RESULT,IS TAG+HEAD, NEW PAGE,

BACKTRACK?

723 (TRY TYPE+HEAD+ACTION,APPLY ACTION+HEAD,TERMINALS,G+HEAD;
 724 APPLY PREDICATE+HEAD,TERMINALS,G+HEAD).

725 SKIP UNTIL POINT:

726 NXT: ERROR + SKIPPED + INPT,
 727 (R + POINT; NEXT SYMBOL,:NXT).

728

729 [4.4.9 TERMINALS]

730 ACTION: TERMINALS,FORGET GLOBALS,MAY BE TERMINAL,PUT DECL,FORGET TERMINALS,
 731 PUT CALL.

732 TERMINALS:

733 BLANK LINE,FORGET GLOBALS,NEW LINE,EQUAL+PLOC+MIN LOC;
 734 APPLY ACTION+XREAD,APPLY ACTION+XOUT,APPLY ACTION+INIT READ.

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735 ..... BLANK LINE, G + INIT READ, G + SEMICOLON, NEW LINE,
736 ..... BLANK LINE, FORGET TERMINALS, NEW LINE, BLANK LINE.
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737 FORGET GLOBALS-P-TERM:
738 MAKE+P+MIN GLOB, MAKE+PLOC+MIN LOC,
739 NXT: EQUAL + P + PGLOB;
740 GET+LGLOB+P+TERM, MAY BE TERMINAL+TERM, INCR+P, :NXT.
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741 MAY BE TERMINAL+X:
742 WAS TERMINAL+X, PUT DECL+X, ENTER LOCAL+X;
743 WARN+X.
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744 WAS TERMINAL+X-S:
745 GET STATE+X+S, EQUAL+S+APPLIED,
746 GET SORT +X+S, EQUAL+S+GLOBAL,
747 GET TYPE+X+S, EQUAL+S+POINTER.
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748 PUT DECL + X:
749 NEW LINE, G+Q INTEGER, G+X, G+SEMICOLON, POSITION+64, INFORM+X.
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750 FORGET TERMINALS-P-TERM:
751 MAKE+P+MIN LOC,
752 NXT: EQUAL+P+PLOC;
753 GET+LLOC+P+TERM, PUT CALL+X READ+TERM, INCR+P, :NXT.
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754 PUT CALL +PROC + X:
755 NEW LINE, G + PROC, G + OPEN, G + X, G + CLOSE, G + SEMICOLON.
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756
757 [4.4.10 POST MORTEM]
758 'ACTION', POST MORTEM, DICTIONARY, LIST TAGS, ENTRIES.
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197
198

759 POST MORTEM:
760 NLCR, OUTINT+MIN TAG, OUTINT+PTAG, OUTINT+MAX TAG,
761 NLCR, OUTINT+MIN BOLD, OUTINT+PBOLD, OUTINT+MAX BOLD,
762 NLCR, OUTINT+MIN SPEC, OUTINT+PSPEC, OUTINT+MAX SPEC,
763 NLCR, OUTINT+MIN CONS, OUTINT+PCONS, OUTINT+MAX CONS,
764 NLCR, OUTINT+MIN LOC, OUTINT+XLOC, OUTINT+MAX LOC,
765 NLCR, OUTINT+MIN GLOB, OUTINT+PGLOB, OUTINT+MAX GLOB,
766 NLCR, OUTINT+MIN MACR, OUTINT+PMACR, OUTINT+MAX MACR,
767 NLCR, OUTINT+MIN TEXT, OUTINT+PTEXT, OUTINT+MAX TEXT,
768 NLCR, OUTINT+LINE, OUTINT+CARD, NLCR, NLCR, DICTIONARY.
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769 DICTIONARY:
770 LIST TAGS + FIRST TAG.
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771 LIST TAGS + X = P-STATE-SORT-TYPE:
772 EQUAL + X + 0;
773 GET LEFT+X+P, LIST TAGS+P,
774 GET STATE+X+STATE, GET SORT+X+SORT, GET TYPE+X+TYPE.
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775 EQUAL+SORT+0,:SKP;
776 NLCR,OUT+X,POSITION+32,OUT+STATE,OUT+SORT,OUT+TYPE,
777 POSITION+59,ENTRIES+X,
778 SKP: GET RIGHT+X+R,LIST TAG$+P.
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285 [4.5. HEART OF COMPILER COMPILER]
286 'ACTION'SENTENCE,START SYSTEM,COMPILER DESCRIPTION.
287 'POINTER' INIT RESTORE,DO RESTORE,XREAD,XOUT,INIT READ,TITLE.
288 SENTENCE:
289 START SYSTEM, BLANK LINE,
290 BEGIN,COMPILER DESCRIPTION,ENDBR,
291 BLANK LINE,NEW PAGE,POST MORTEM.
292 START SYSTEM:
293 READ+INIT RESTORE,READ+DO RESTORE,READ+X READ,READ+X OUT,
294 READ+INIT READ,READ+TITLE,
295 SET+GIVERS+O+FALSE,SET+GIVE TRACE+FALSE,SET+GIVE TEXT+TRUE,
296 MAKE+PLOC+MIN LOC,MAKE+PLOB+MIN GLOB,MAKE+PTEXT+MIN TEXT,
297 MAKE+PMACR+MIN MACR,MAKE+CARD+0.MAKE+CPOS+1,MAKE+XLOC+0,
298 SET+LINED UP+FALSE,SET+LEGIBLE+TRUE,
299 MAKE+INDENTATION+0,MAKE+LINE+0,MAKE+POS+0,OUT+TITLE,READ+INPT,
300 COMPILER DESCRIPTION:
301 NXT: SPECIFICATION, :NXT;
302 DECLARATION, :NXT;
303 COMMAND, :NXT;
304 COMMENT, :NXT;
305 STARTING SYMBOL;
306 RULE, :NXT;
307 SKIP UNTIL POINT, :NXT.
308
309 'RESULT' SENTENCE.
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APPLIED GLOBAL POINTER TAGFULL
APPLIED GLOBAL POINTER BOLDFULL
APPLIED GLOBAL POINTER SPECFULL
APPLIED GLOBAL POINTER CONSFULL
APPLIED GLOBAL POINTER WRONGINPUTCODE
APPLIED GLOBAL POINTER BLANK
APPLIED GLOBAL POINTER LOGFULL
APPLIED GLOBAL POINTER GLOBFULL
APPLIED GLOBAL POINTER MACRFULL
APPLIED GLOBAL POINTER TEXTFULL
APPLIED GLOBAL POINTER DEFINED
APPLIED GLOBAL POINTER GLOBAL
APPLIED GLOBAL POINTER ACTION
APPLIED GLOBAL POINTER PREDICATE
APPLIED GLOBAL POINTER LOCAL
APPLIED GLOBAL POINTER POINTER
APPLIED GLOBAL POINTER LABEL
APPLIED GLOBAL POINTER APPLIED
APPLIED GLOBAL POINTER XIMPOSSIBLE
APPLIED GLOBAL POINTER MACRO
APPLIED GLOBAL POINTER WRONGNUMBEROFPARAMETERS
APPLIED GLOBAL POINTER FLAG
APPLIED GLOBAL POINTER LIST
APPLIED GLOBAL POINTER EXTERNAL
APPLIED GLOBAL POINTER COMPA
APPLIED GLOBAL POINTER POINT
APPLIED GLOBAL POINTER INTEGER
APPLIED GLOBAL POINTER SEMICOLON
APPLIED GLOBAL POINTER QBOCLEAN
APPLIED GLOBAL POINTER GARPAY
APPLIED GLOBAL POINTER SUB
APPLIED GLOBAL POINTER COLON
APPLIED GLOBAL POINTER BUS
APPLIED GLOBAL POINTER PLUS
APPLIED GLOBAL POINTER MINUS
APPLIED GLOBAL POINTER OPEN
APPLIED GLOBAL POINTER CLOSE
APPLIED GLOBAL POINTER STRINGQUOTE
APPLIED GLOBAL POINTER WRONGAFFIX
APPLIED GLOBAL POINTER ONE
APPLIED GLOBAL POINTER TWO
APPLIED GLOBAL POINTER THREE
APPLIED GLOBAL POINTER FOUR
APPLIED GLOBAL POINTER FIVE
APPLIED GLOBAL POINTER XPARAMETERERROR
APPLIED GLOBAL POINTER POSSIBLYBACKTRACKNECESSARY
APPLIED GLOBAL POINTER ALTERNATIVEEVERREACHED
APPLIED GLOBAL POINTER QTRUE
APPLIED GLOBAL POINTER QFALSE
APPLIED GLOBAL POINTER QGOTO
APPLIED GLOBAL POINTER EQUALS
APPLIED GLOBAL POINTER QBEGIN
APPLIED GLOBAL POINTER QEND
APPLIED GLOBAL POINTER QPROCEDURE
APPLIED GLOBAL POINTER TIMES
APPLIED GLOBAL POINTER QIF
APPLIED GLOBAL POINTER NOT
APPLIED GLOBAL POINTER QTHEN
APPLIED GLOBAL POINTER RSTOON

APPLIED GLOBAL POINTER RSTOUFF
APPLIED GLOBAL POINTER SHORT
APPLIED GLOBAL POINTER LONG
APPLIED GLOBAL POINTER TRACEON
APPLIED GLOBAL POINTER TRACEOFF
APPLIED GLOBAL POINTER GCOMMENT
APPLIED GLOBAL POINTER RESULT
APPLIED GLOBAL POINTER SKIPPED

100001 104390 110000
200001 204173 200600
300001 304118 300100
400001 400348 400300
500001 500068 500200
600001 600353 601000
700001 700534 701100
800001 801910 804000
809 211

ACCC	MACRO POINTER	244	242	20
ACCENT	GLOBAL PREDICATE	308	306	304
ACTION	APPLIED GLOBAL POINTER	723	704	501
ADD	MACRO ACTION	783	293	137
ADDRULT	MACRO ACTION	344	56	
ADDRPLAC	GLOBAL ACTION	466	390	389
ADVANCED	GLOBAL FLAG	700	684	603
AFFIXPACKOPTION	GLOBAL ACTION	556	555	547
AHEAD	GLOBAL PREDICATE	608	606	510
ALTERNATIVE	APPLIED GLOBAL POINTER	607		
ALTERNATIVE	GLOBAL PREDICATE	693	683	680
APPLIED	APPLIED GLOBAL POINTER	745	568	511
APPLYACTION	GLOBAL ACTION	734	723	434
APPLYLABEL	GLOBAL ACTION	637	438	424
APPLYPREDICATE	GLOBAL ACTION	724	436	424
APPLY	GLOBAL ACTION	696	589	561
BECOMES	GLOBAL ACTION	641	635	632
BEGIN	GLOBAL ACTION	790	692	672
BLANKLINE	GLOBAL ACTION	791	789	736
BLANK	APPLIED GLOBAL POINTER	455	454	451
BOLDCHAR	GLOBAL PREDICATE	309	307	305
BOLDFULL	APPLIED GLOBAL POINTER	312		
BOUNDARYX1	GLOBAL PREDICATE	667	662	659
BOUNDARYX2	GLOBAL PREDICATE	669	663	660

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[illegible]

ISCONS	GLOBAL PREDICATE	589	563	542	365				
ISTAG	GLOBAL PREDICATE	722	696	686	675	670	668	652	589 561 539 532 525 511 505 496 363
ISTYPE	GLOBAL PREDICATE	504	490	489					
ISTYPEOFLCAL	GLOBAL PREDICATE	500	499						
LABEL	APPLIED GLOBAL POINTER	439	432						
LAYOUTSYMBOL	GLOBAL PREDICATE	371	257						
LEFTANDSIDE	GLOBAL PREDICATE	651	649						
LEPT	MACRO FLAG	286	67						
LEGIBLE	GLOBAL FLAG	798	713	712	221	218	211	89	
LESS	MACRO FLAG	393	293	200	193	174	132	129	110 68
LETGIT	GLOBAL PREDICATE	268	267	266	237				
LETTER	GLOBAL PREDICATE	265	235						
LGLOB	GLOBAL LIST	740	412	38					
LINEDUP	GLOBAL FLAG	798	220	218	211	210	205		
LINEUP	GLOBAL ACTION	209	208	204					
LINE	GLOBAL POINTER	799	768	397	396	393	361	252	250 231
LISTDECLARATION	GLOBAL PREDICATE	530	520						
LISTTAGS	GLOBAL ACTION	778	773	771	770	758			
LIST	APPLIED GLOBAL POINTER	532	531	476					
LLCC	GLOBAL LIST	753	408	400	37				
LMACR	GLOBAL LIST	576	571	415	40				
LOCAL	APPLIED GLOBAL POINTER	466	458	439	432	430			
LOCFULL	APPLIED GLOBAL POINTER	410							
LONG	APPLIED GLOBAL POINTER	713							
LOSTLABEL	GLOBAL ACTION	613	612	598					
LSEQ	MACRO FLAG	418	415	412	409	408	347	330	311 275 213 189 124 69
LTEXT	GLOBAL LIST	781	418	393	39				
L	GLOBAL ACTION	644	639	537	216	204			
MACRFULL	APPLIED GLOBAL POINTER	416							
MACRSPECIFICATION	GLOBAL PREDICATE	503	487						

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POS	GLOBAL POINTER	799	445	156	144	137	136	134	124	101
POSSIBLYBACKTRACKNECESSARY	APPLIED GLOBAL POINTER	603								
POSTMORTEM	GLOBAL ACTION	791	759	758	419	416	413	410	348	331 312 296 276 229
PRCHAR	GLOBAL ACTION	253	140	139	122	115	112	111	110	99
PREDICATE	APPLIED GLOBAL POINTER	698	502	492	437	428				
PRINT	GLOBAL ACTION	120	119	99						
PRINT	GLOBAL ACTION	116	105	104	103	99				
PRSYM	GLOBAL ACTION	144	137	134	12					
PSPEC	GLOBAL POINTER	762	361	337	330	328	87			
PTAG	GLOBAL POINTER	760	362	361	295	294	293	284	275	265 87
PTEXT	GLOBAL POINTER	796	767	418	395	88				
PUCHAR	GLOBAL ACTION	201	196	195	187	179	176	175	174	163
PUNCH1	GLOBAL ACTION	185	184	183	162					
PUNCH	GLOBAL ACTION	180	170	169	168	162				
PUT	MACRO ACTION	418	415	412	408	347	330	311	275	50
PUTABS	GLOBAL ACTION	214	212	211	204					
PUTAFFIXEXPRESSION	GLOBAL ACTION	705	702	699	549	547				
PUTCALL	GLOBAL ACTION	754	753	731						
PUTCONNECTAND	GLOBAL ACTION	703	700	601	598					
PUTCONNECTOR	GLOBAL ACTION	693	680	605	598					
PUTDECL	GLOBAL ACTION	748	742	730						
PUTDIRECT	GLOBAL ACTION	554	550	547						
PUTINIT	GLOBAL ACTION	692	672	622	598					
PUTJUMP	GLOBAL ACTION	686	636	621	611	604	603	602	599	
PUTLABEL	GLOBAL ACTION	697	638	626	618	616	610	599		
PUTLEFT	MACRO ACTION	295	287	82						
PUTMACRO	GLOBAL ACTION	573	553	552	548					
PUTMTEXT	MACRO ACTION	508	294	77						
PUTNPARS	MACRO ACTION	469	293	76						
PUTPARAMLESSMACRO	GLOBAL ACTION	583	569	567	548	540	516			

PUTPAR	GLOBAL ACTION	591	582	581	580	579	578	548
PUTPLACE	MACRO ACTION	395	294	78				
PUTRESTORE	GLOBAL ACTION	624	612	606	599			
PUTRIGHT	MACRO ACTION	295	290	83				
PUTSINGLE	GLOBAL ACTION	595	565	561	547			
PUTSORT	MACRO ACTION	457	401	294	80			
PUTSTATE	MACRO ACTION	454	451	450	401	294	79	
FUTTYPE	MACRO ACTION	463	401	295	81			
GARRAY	APPLIED GLOBAL POINTER	531						
GBEGIN	APPLIED GLOBAL POINTER	643						
GBCOLEAN	APPLIED GLOBAL POINTER	654	529					
GCOMMENT	APPLIED GLOBAL POINTER	717						
GEND	APPLIED GLOBAL POINTER	644						
GFALSE	APPLIED GLOBAL POINTER	635						
QGCTO	APPLIED GLOBAL POINTER	637						
GIF	APPLIED GLOBAL POINTER	702	699					
GINTEGER	APPLIED GLOBAL POINTER	749	674	666	531	522		
GPROCEDURE	APPLIED GLOBAL POINTER	656						
GTPEN	APPLIED GLOBAL POINTER	702	699					
GTRUE	APPLIED GLOBAL POINTER	632						
QUOTE	MACRO POINTER	242	20					
READBOLD	GLOBAL PREDICATE	373	300					
READCONS	GLOBAL PREDICATE	375	342					
READMACRO	GLOBAL ACTION	507	505	483				
READSPEC	GLOBAL PREDICATE	374	327					
READTAG	GLOBAL PREDICATE	372	264					
READ	GLOBAL ACTION	809	799	794	793	719	376	370 369 358
REDEFINE1	GLOBAL PREDICATE	447	444					
REDEFINE	GLOBAL ACTION	568	532	525	505	496	443	439 437 435 432 430 428 426 424
REPLACE	GLOBAL ACTION	394	393	392	389			

RESERVEADDRESS	GLOBAL ACTION	362	292	290	287	283	263
RESTEQUINADIFFIXES	GLOBAL ACTION	663	662	661	660	659	647

RESTBOUND AFFIXES	663	662	661	660	659	647
GLOBAL ACTION						

RESULTS **GLOBAL ACTION** **655** **654** **653** **647**

RESULT **APPLIED GLOBAL POINTER** **722**

RESYN **GLOBAL POINTER** **362** **248** **13**

RIGHT HAND SIDE	GLOBAL PREDICATE	678	649
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RSTOCFF	APPLIED GLOBAL POINTER	711
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RSTOCK APPLIED GLOBAL POINTER 710

RULE **GLOBAL PREDICATE** **806 648**

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489 306

SERVICOLON **APPLIED GLOBAL POINTER** 755 749 735 718 694 681 677 640 637 606 537 527

SEMICOL **GLOBAL ACTION** 705 694 682 666 657 640 635 632 599

[illegible][illegible]

SHIFT LINES	GLOBAL ACTION	537	527	506	497	135	100
		650					

SHORT..... **APPLIED GLOBAL POINTER**..... **712**

SIGNAL	GLOBAL ACTION	406	150	147
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SKIPPED APPLIED GLOBAL POINTER 726

SKI PLANTILPCINT.....	GLOBAL ACTION.....	807	725	708
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SPACECARD GLOBAL ACTION 194 193 162

SPACECODE	MACRO POINTER	258	256	195	189	139	20
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SPACES	136	127	126	124	99
GLOBAL ACTION	136	127	126	124	99

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SPECCHAR	GLOBAL PREDICATE	328	245
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SPECFULL	APPLIED GLOBAL POINTER	331
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SPECIFICATION.....	GLOBAL PREDICATE.....	801	484
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SPECX	GLOBAL POINTER	670	666	664	663	652	384
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STACKBOLD	GLOBAL ACTION	310	309	308	306	304	302	299
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STACKCONS	GLOBAL ACTION	346	345	341
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STACKSPEC	GLOBAL ACTION	329	328	326
STACKTAG	GLOBAL ACTION	274	273	263
STARTINGSYNBCL	GLOBAL PREDICATE	805	721	
STARTSYSTEM	GLOBAL ACTION	792	789	786
STRINGOJCTE	APPLIED GLOBAL POINTER	594	562	
SUPTR	MACRO ACTION	397	124	55
SUB	APPLIED GLOBAL POINTER	717	533	
TABCARD	GLOBAL ACTION	221	214	198 191 162
TABCODE	MACRO POINTER	260	198	142 137 20
TAB	GLOBAL ACTION	142	137	132 99
TABLINE	GLOBAL ACTION	643	639	221 204
TABS	GLOBAL ACTION	130	125	99
TAGFULL	APPLIED GLOBAL POINTER	296	276	
TBOLD	GLOBAL LIST	323	316	311 169 104 34
TCONS	GLOBAL LIST	351	350	347 171 106 36
TERMINALS	GLOBAL ACTION	732	730	724 723 708
TEXTFULL	APPLIED GLOBAL POINTER	419		
THREE	APPLIED GLOBAL POINTER	580		
TIMES	APPLIED GLOBAL POINTER	668		
TITLE	GLOBAL POINTER	799	794	787
TOEND	GLOBAL POINTER	679	621	618 384
TOLOST	GLOBAL POINTER	679	620	614 611 603 384
TOEXT	GLOBAL POINTER	688	687	684 625 620 615 609 607 604 602 384
TRACEOFF	APPLIED GLOBAL POINTER	715		
TRACEON	APPLIED GLOBAL POINTER	714		
TRACE	GLOBAL ACTION	405	404	148 147
TREATAFFIX	GLOBAL ACTION	560	558	557 547
TREATDECLIST	GLOBAL ACTION	529	523	522 515
TREATSPECLIST	GLOBAL ACTION	499	495	489 483
TRIRESULT	GLOBAL ACTION	690	630	599

TRUE	MACRO POINTER	798	795	714	713	710	700	211	73
TRYSCT	GLOBAL PREDICATE	551	550	456	448				
TRYSTATE	GLOBAL PREDICATE	511	449	448	442				
TRYTYPE	GLOBAL PREDICATE	723	704	701	698	462	448		
TSPEC	GLOBAL LIST	334	333	330	170	105	35		
TTAG	GLOBAL LIST	280	275	168	103	33			
TWO	APPLIED GLOBAL POINTER	579							
UNPACK2	MACRO ACTION	781	61						
UNPACK3	MACRO ACTION	186	121	59					
U	GLOBAL ACTION	643	639	531	215	204			
WARN	GLOBAL ACTION	743	402	470	388				
WASACTION	GLOBAL PREDICATE	653	634	631	552	471			
WASAFFIX	GLOBAL PREDICATE	566	480	473					
WASBOLD	MACRO FLAG	594	364	169	104	43			
WASCCNS	MACRO FLAG	592	440	433	365	171	106	45	
WASDIGIT	MACRO FLAG	240	17						
WASLETGIT	MACRO FLAG	238	16						
WASLETTER	MACRO FLAG	236	15						
WASMACRO	GLOBAL PREDICATE	567	480	477					
WASPARAMLESSMACRO	GLOBAL PREDICATE	583	540	479					
WASSPECCH	MACRO FLAG	246	18						
WASSPEC	MACRO FLAG	170	105	44					
WASTAG	MACRO FLAG	480	439	432	363	168	157	103	42
WASTERMINAL	GLOBAL PREDICATE	744	742						
WRITEINT1	GLOBAL ACTION	179	177	176	162				
WRITEINT	GLOBAL ACTION	190	174	173	172	171	162		
WRITE	GLOBAL ACTION	208	167	162					
WRNGAFFIX	APPLIED GLOBAL POINTER	564							
WRNGINPUTCODE	APPLIED GLOBAL POINTER	376							
WRNGNUMBEROFPARAMETERS	APPLIED GLOBAL POINTER	470							

XIMPCSSIBLE

APPLIED GLOBAL POINTER 445

XLOC

GLOBAL POINTER 797 764 409 88

XOUT

GLOBAL POINTER 793 787 734

XPARAMETERERROR

APPLIED GLOBAL POINTER 593 590

XREAD

GLOBAL POINTER 793 787 753 734

A1262J.19. KCSTER

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1  'BEGIN'
2  'COMMENT' 4,1,GENERALCOMPILERDESCRIBED INCOL*****
3  'COMMENT' 4,1,GENERALENVIRONMENT;
4  'COMMENT' 4,1,1,INTERFACETW:MACHIN;
5  'COMMENT' 4,1,2,STACKS;
6  'INTEGER' 8,1,ARRAY,TTAG(100001:100001),TBOLD(200001:200600),TSPEC(300001:300100),TCONS(400001:400300),LLOC(500001:500200),LGL08(600001:60100),LTEXT(80001:804000),LMACR(700001:701100);
7  'COMMENT' 4,1,3,MACROS;
8  'COMMENT' 4,1,4,1,3,MACROS;
9  'COMMENT' 4,1,4,1,4,1,3,MACROS;
10 'INTEGER' PTAG,PBOLD,PSPEC,PCONS,PLOC,PGL08,PTEXT,PMACR,XLOC;
11 'EOCLEAN' GIVETEXT,GIVETRACE,LEGIBLE;
12 'INTEGER' FIRSTTAG;
13 'COMMENT' 4,2,OUTPUT;
14 'COMMENT' 4,2,1,PRINTINTERSECTION;
15 'INTEGER' POS;
16 'PROCEDURE' OUT(X); 'INTEGER' X; 'BEGIN' 'INTEGER' EL; 'IF' -(100001,LE'X'<PTAG) 'THEN' 'GOTO' 1; 'PRINT' (TTAG,X); 'GOTO' 999; 1: 'IF' -(200001,LE'X'<PBOLD) 'THEN' 'GOTO' 2; 'PRINT' (TBOLD,X); 'GOTO' 999; 2: 'IF' -(300001,LE'X'<PSPEC) 'THEN' 'GOTO' 3; 'PRINT' (TSPEC,X); 'GOTO' 999; 3: 'IF' -(400001,LE'X'<PCONS) 'THEN' 'GOTO' 4; 'EL' = TCONS(X); 'OUTINT' (EL); 'GOTO' 999; 4: 'OUTINT' (X); 999: 'END';
17 'PROCEDURE' OUTINT(X); 'INTEGER' X; 'BEGIN' 'INTEGER' QUOT,REM; 'SPACE'; 'IF' -(X<0) 'THEN' 'GOTO' 1; 'REM' = X; 'REM' = -REM; 'PRCHAR' (65); 'OUTINT' (REM); 'GOTO' 999; 1: 'IF' -(X=0) 'THEN' 'GOTO' 2; 'PRCHAR' (0); 'SPACE'; 'GOTO' 999; 2: 'QUOT' = X/'10'; 'REM' = X-10*QUOT; 'OUTINT1' (QUOT); 'END';
18 'PROCEDURE' OUTINT1(X); 'INTEGER' X; 'BEGIN' 'INTEGER' QUOT,REM; 'IF' -(X=0) 'THEN' 'GOTO' 1; 'GOTO' 999; 1: 'QUOT' = X/'10'; 'REM' = X-10*QUOT; 'OUTINT1' (QUOT); 'PRCHAR' (REM); 999: 'END';
19 'PROCEDURE' PRINT(LIST,Y); 'INTEGER' Y; 'BEGIN' 'INTEGER' X,LX; X=Y; 'SPACE'; 'RST': LX=LIST[X]; 'BEGIN' 'IF' -(LX<0) 'THEN' 'GOTO' 2; 'LX' = -LX; 'PRINT' (LX); 'GOTO' 999; 2: 'PRINT' (LX); X=X+1; 'GOTO' 'RST'; 'END'; 999: 'END';
20 'PROCEDURE' PRINT1(X); 'INTEGER' X; 'BEGIN' 'INTEGER' X1,X2,X3; X1=X/'16384'; X3=X-16384*X1; X2=X3/'128'; X3=X3-128*X2; 'PRCHAR' (X1); 'PRCHAR' (X2); 'PRCHAR' (X3); 'END';
21 'PROCEDURE' POSITION(X); 'INTEGER' X; 'BEGIN' 'INTEGER' TBS,SPCES,SPCES; 'IF' -(X<0) 'THEN' 'GOTO' 1; 'SPCES' (SPCES); 'GOTO' 999; 1: 'NLCR'; TBS=X/'8'; 'SPCES' = X-8*TBS; 'TBS' (TBS); 'BEGIN' 'IF' -(SPCES=0) 'THEN' 'GOTO' 3; 'SPCES' (SPCES); 'GOTO' 999; 'END'; 999: 'END';
22 'PROCEDURE' SPACES(X); 'INTEGER' X; 'BEGIN' 'INTEGER' N; N=0; 'SPC': 'IF' -(N<X) 'THEN' 'GOTO' 1; 'SPACE' N; N=N+1; 'GOTO' 'SPC'; 1: 'END';
23 'PROCEDURE' TABS(X); 'INTEGER' X; 'BEGIN' 'INTEGER' N; N=0; 'TBS': 'IF' -(N<X) 'THEN' 'GOTO' 1; 'TAB' N; N=N+1; 'GOTO' 'TBS'; 1: 'END';
24 'PROCEDURE' NLCR; 'BEGIN' 'PRSYM' (119); 'POS' = 0; 'END';
25 'PROCEDURE' SHIFT2(LINES); 'BEGIN' 'INTEGER' OLDPOS; OLDPOS = POS; 'NLCR'; 'NLCR'; 'SPACES' (OLDPOS); 'END';
26 'PROCEDURE' TAB; 'BEGIN' 'PRSYM' (118); 'POS' = POS+8; 'END';
27 'PROCEDURE' SPACE; 'BEGIN' 'PRCHAR' (93); 'END';
28 'PROCEDURE' PRCHAR(X); 'INTEGER' X; 'BEGIN' 'IF' -(X=119) 'THEN' 'GOTO' 1; 'NLCR'; 'GOTO' 999; 1: 'IF' -(X=118) 'THEN' 'GOTO' 2; 'TAB'; 'GOTO' 999; 2: 'IF' -(X=63) 'THEN' 'GOTO' 3; 'GOTO' 999; 3: 'PRSYM' (X); 'POS' = POS+1; 999: 'END';
29 'COMMENT' 4,2,2,TRACEADMINISTRATION;
30 'PROCEDURE' TRACE(X); 'INTEGER' X; 'BEGIN' 'IF' -GIVETRACE 'THEN' 'GOTO' 1; 'POSITION' (32); 'INFORM' (X); 'GOTO' 999; 1: 999: 'END';
31 'PROCEDURE' SIGNAL(X); 'INTEGER' X; 'BEGIN' 'POSITION' (16); 'INFORM' (X); 'END';
32 'PROCEDURE' INFORM(X); 'INTEGER' X; 'BEGIN' 'INTEGER' STATE, SORT, TYPE; STATE = TTAG(X-5); SORT = TTAG(X-4); TYPE = TTAG(X-3); 'OUT' (STATE); 'OUT' (SORT); 'OUT' (TYPE); 'OUT' (X); 'END';
33 'PROCEDURE' ERROR(TEXT,INFO); 'INTEGER' INFO,TEXT; 'BEGIN' 'INTEGER' OLDPOS; OLDPOS = POS; 'NLCR'; 'OUT' (TEXT); 'BEGIN' 'IF' -(100001,LE'INFO'<PTAG) 'THEN' 'GOTO' 2; 'INFORM' (INFO); 'POSITION' (OLDPOS); 'GOTO' 999; 2: 'IF' -(INFO=0) 'THEN' 'GOTO' 3; 'POSITION' (OLDPOS); 'GOTO' 999; 3: 'OUT' (INFO); 'POSITION' (OLDPOS); 'GOTO' 999; 'END'; 999: 'END';
34 'COMMENT' 4,2,3,CARDPUNCHSECTION;
35 'INTEGER' CARD,CPOS;
36 'PROCEDURE' WRITE(X); 'INTEGER' X; 'BEGIN' 'INTEGER' EL; 'IF' -(100001,LE'X'<PTAG) 'THEN' 'GOTO' 1; 'PUNCH' (TTAG,X); 'GOTO' 999; 1: 'IF' -(200001,LE'X'<PBOLD) 'THEN' 'GOTO' 2; 'PUNCH' (TBOLD,X); 'GOTO' 999; 2: 'IF' -(300001,LE'X'<PSPEC) 'THEN' 'GOTO' 3; 'PUNCH' (TSPEC,X); 'GOTO' 999; 3: 'IF' -(400001,LE'X'<PCONS) 'THEN' 'GOTO' 4; 'EL' = TCONS(X); 'WRITE' (EL); 'GOTO' 999; 4: 'WRITE' (X); 999: 'END';
37 'PROCEDURE' WRITEINT(X); 'INTEGER' X; 'BEGIN' 'INTEGER' QUOT,REM; 'IF' -(X<0) 'THEN' 'GOTO' 1; 'REM' = X; 'REM' = -REM; 'PCHAR' (65); 'WRITE' (REM); 'GOTO' 999; 1: 'IF' -(X=0) 'THEN' 'GOTO' 2; 'PCHAR' (0); 'GOTO' 999; 2: 'QUOT' = X/'10'; 'REM' = X-10*QUOT; 'WRITEINT1' (QUOT); 'END';
38 'PROCEDURE' WRITEINT1(X); 'INTEGER' X; 'BEGIN' 'INTEGER' QUOT,REM; 'IF' -(X=0) 'THEN' 'GOTO' 1; 'GOTO' 999; 1: 'QUOT' = X/'10'; 'REM' = X-10*QUOT; 'WRITEINT1' (QUOT); 'PCHAR' (REM); 999: 'END';
39 'PROCEDURE' PUNCH(LIST,Y); 'INTEGER' Y; 'LIST': X=LX; X=Y; 'RST': LX=LIST[X]; 'BEGIN' 'IF' -(LX<0) 'THEN' 'GOTO' 2; 'LX' = -LX; 'PUNCH' (LX); 'GOTO' 999; 2: 'PUNCH' (LX); X=X+1; 'GOTO' 'RST'; 'END'; 999: 'END';

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40 'PROCEDURE PUNCH1(X); 'INTEGER' X; 'BEGIN' 'INTEGER' X1,X2,X3;X1:=X/'16384;X3:=X-16384*X1;X2:=X3/'128;X3:=X3-128*X2;PUNCHAR(X1);PUNCHAR(X2);P
UCHAR(X3);'END';
41 'PROCEDURE NEWCARD; 'BEGIN' SPC; 'IF' ~(CPOS<72) 'THEN' 'GOTO' 1; CSYM(93); CPOS:=CPOS+1; WRITE INT(CARD) ICSYM(119); CARD:=CARD+
1; CPOS:= 'END';
42 'PROCEDURE TABCARD; 'BEGIN' 'INTEGER' SPACES, TBS; TBS:=CPOS/'18; SPACES:=CPOS-8*TBS; SPC; 'IF' ~(SPACES<8) 'THEN' 'GOTO' 1; SPACECARD; SPACES:=SPACES+1; 'G
OTO' SPC; 'END';
43 'PROCEDURE SPACECARD; 'BEGIN' PUCHAR(93); 'END';
44 'PROCEDURE PUCHAR(X); 'INTEGER' X; 'BEGIN' 'IF' ~(X=119) 'THEN' 'GOTO' 1; NEWCARD; 'GOTO' 999; 1: 'IF' ~(X=118) 'THEN' 'GOTO' 2; TABCARD; 'GOTO' 999; 2: 'IF' ~
(X=63) 'THEN' 'GOTO' 3; 'GOTO' 999; 3: 'IF' ~(CPOS<72) 'THEN' 'GOTO' 4; CSYM(X); CPOS:=CPOS+1; 'GOTO' 999; 4: CSYM(133); CSYM(119); CPOS:=1; PUCHAR(X); 999: 'END';
45 'COMMENT' 4.2.4 RESULT ADMINISTRATION;
46 'POCLEAN' LINEDUP;
47 'INTEGER' INDENTATION;
48 'PROCEDURE G(X); 'INTEGER' X; 'BEGIN' LINEDUP; WRITE(X); 'END';
49 'PROCEDURE LINEDUP; 'BEGIN' 'IF' ~LINEDUP 'THEN' 'GOTO' 1; 'GOTO' 999; 1: 'IF' ~LEGIBLE 'THEN' 'GOTO' 2; PUTABS(INDENTATION); LINEDUP:= 'TRUE'; 'GOTO' 999; 2:
999: 'END';
50 'PROCEDURE PUTABS(N); 'INTEGER' N; 'BEGIN' 'INTEGER' N1; 'IF' ~(N<0) 'THEN' 'GOTO' 1; 'GOTO' 999; 1: N1:=N; N1:=N1-1; TABCARD; PUTABS(N1); 999: 'END';
51 'PROCEDURE U; 'BEGIN' INDENTATION:=INDENTATION+1; 'END';
52 'PROCEDURE L; 'BEGIN' INDENTATION:=INDENTATION-1; 'END';
53 'PROCEDURE NEWLINE; 'BEGIN' 'IF' ~LEGIBLE 'THEN' 'GOTO' 1; 'IF' ~LINEDUP 'THEN' 'GOTO' 1; BLANKLINE; 'GOTO' 999; 1: 999: 'END';
54 'PROCEDURE BLANKLINE; 'BEGIN' NEWCARD; LINEDUP:= 'FALSE'; 'END';
55 'PROCEDURE TABLINE; 'BEGIN' 'IF' ~LEGIBLE 'THEN' 'GOTO' 1; TABCARD; 'GOTO' 999; 1: 999: 'END';
56 'COMMENT' 4.3 INPUT;
57 'INTEGER' LINE, CHAR, INPT;
58 'COMMENT' 4.3.1 READING CHARACTERS;
59 'POCLEAN' 'PROCEDURE LETGIT(X); 'INTEGER' X; 'BEGIN' 'IF' ~(CHAR>9*CHAR36) 'THEN' 'GOTO' 1; X:=CHAR; NEXTNONSPACECHAR; LETTER:= 'TRUE'; 'GOTO' 999; 1: L
ETTER:= 'FALSE'; 999: 'END';
60 'POCLEAN' 'PROCEDURE LETGIT(X); 'INTEGER' X; 'BEGIN' 'IF' ~(CHAR<36) 'THEN' 'GOTO' 1; X:=CHAR; NEXTNONSPACECHAR; LETGIT:= 'TRUE'; 'GOTO' 999; 1: LETGIT:=
'FALSE'; 999: 'END';
61 'POCLEAN' 'PROCEDURE DIGIT(X); 'INTEGER' X; 'BEGIN' 'IF' ~(CHAR<10) 'THEN' 'GOTO' 1; X:=CHAR; NEXTNONSPACECHAR; DIGIT:= 'TRUE'; 'GOTO' 999; 1: DIGIT:= 'F
ALSE'; 999: 'END';
62 'POCLEAN' 'PROCEDURE BOLDCHAR(X); 'INTEGER' X; 'BEGIN' 'IF' ~(CHAR=121) 'THEN' 'GOTO' 1; X:=120; NEXTCHAR; BOLDCHAR:= 'TRUE'; 'GOTO' 999; 1: X:=CHAR; NEXT
CHAR; BOLDCHAR:= 'TRUE'; 999: 'END';
63 'POCLEAN' 'PROCEDURE ACCENT; 'BEGIN' 'IF' ~(CHAR=120) 'THEN' 'GOTO' 1; NEXTCHAR; ACCENT:= 'TRUE'; 'GOTO' 999; 1: ACCENT:= 'FALSE'; 999: 'END';
64 'POCLEAN' 'PROCEDURE SPECCHAR(X); 'INTEGER' X; 'BEGIN' 'IF' ~(CHAR>63) 'THEN' 'GOTO' 1; X:=CHAR; NEXTNONSPACECHAR; SPECCHAR:= 'TRUE'; 'GOTO' 999; 1: SPEC
CHAR:= 'FALSE'; 999: 'END';
65 'PROCEDURE NEXTCHAR; 'BEGIN' DISPLAY CHARACTER; CHAR:=RESYM; 'END';
66 'PROCEDURE DISPLAY CHARACTER; 'BEGIN' 'IF' ~(CHAR=119) 'THEN' 'GOTO' 1; LINE:=LINE+1; 'BEGIN' 'IF' ~GIVETEXT 'THEN' 'GOTO' 2; POSITION(133); OUT INT1(CAR
D); POSIT: CH(48); OUT INT1 LINE; POSITION(52); 'GOTO' 999; 2: 'GOTO' 999; 'END'; 1: 'IF' ~GIVETEXT 'THEN' 'GOTO' 4; PRCHAR(CHAR); 'GOTO' 999; 4: 999: 'END';
67 'PROCEDURE NEXTNONSPACECHAR; 'BEGIN' CHR: NEXTCHAR; 'BEGIN' 'IF' ~(CHAR=93) 'THEN' 'GOTO' 2; 'GOTO' CHR; 2: 'GOTO' 999; 'END'; 999: 'END';
68 'POCLEAN' 'PROCEDURE LAYOUTSYMBOL; 'BEGIN' 'IF' ~(CHAR=93) 'THEN' 'GOTO' 1; LAYOUTSYMBOL:= 'TRUE'; 'GOTO' 999; 1: 'IF' ~(CHAR=119) 'THEN' 'GOTO' 2; LAYOUT
SYMBOL:= 'TRUE'; 'GOTO' 999; 2: 'IF' ~(CHAR=118) 'THEN' 'GOTO' 3; LAYOUTSYMBOL:= 'TRUE'; 'GOTO' 999; 3: LAYOUTSYMBOL:= 'FALSE'; 999: 'END';
69 'COMMENT' 4.3.2 TAGSYMBOLS;
70 'POCLEAN' 'PROCEDURE READTAG(X); 'INTEGER' X; 'BEGIN' 'INTEGER' T1,T2,T3,T4; 'IF' ~LETTER(T1) 'THEN' 'GOTO' 1; X:=T3; T4; 'IF' ~LETTER(T1) 'THEN' 'GOTO' 1; X:=T3; T4; 'IF' ~LETTER(T1) 'THEN' 'GOTO' 2; 'BEGIN' 'IF' ~LETGIT(T3) 'THEN' 'GOTO' 3; 'BEGIN' 'IF' ~LETGIT(T4) 'THEN' 'GOTO' 4; T:= (T1*128+T2)*128+T3; STACKTAG(T); T1:=T4; 'GOTO' NEXT; 4: T:= (T1
*128+T2)*128+T3; 'GOTO' LST; 'END'; 3: T:= (T1*128+T2)*128+63; 'GOTO' LST; 'END'; 2: T:= (T1*128+63)*128+63; LST: T:=T; STACKTAG(T); READTAG:= 'TRUE'; 'GOTO' 999;
'END'; 1: READTAG:= 'FALSE'; 999: 'END';
71 'PROCEDURE STACKTAG(X); 'INTEGER' X; 'BEGIN' 'IF' ~(PTAG<110000) 'THEN' 'GOTO' 1; TTAG(PTAG)=X; PTAG:=PTAG+1; 'GOTO' 999; 1: ERRORC(TAGFULL,0); POSTWOR
K: EXIT; 999: 'END';
72 'PROCEDURE ENTERTAG(X); 'INTEGER' X; 'BEGIN' 'INTEGER' Y,X1,Y1,WX,WY; Y1:=FIRSTTAG; NXY: X1:=X; Y1:=Y; NXW: WX:=TTAG[X1]; WY:=TTAG[Y1]; 'BEGIN' 'IF' ~(W
X=WX) 'THEN' 'GOTO' 2; 'BEGIN' 'IF' ~(WX<0) 'THEN' 'GOTO' 3; 'BEGIN' 'IF' ~(Y=X) 'THEN' 'GOTO' 4; RESERVEADMINSYMBOL; 'GOTO' 999; 4: PTAG:=X; X:=Y; 'GOTO' 999; 'END'; 3: X:=Y1; Y1:=Y; 'GOTO' 999; 'END'; 2: 'IF' ~(ABS(WX)<ABS(WY)) 'THEN' 'GOTO' 7; WY:=TTAG[Y-2]; 'BEGIN' 'IF' ~(WY=0) 'THEN' 'GOTO' 8; TTAG[Y-2]:=X; RESERVEADMINSY
MBOL; 'GOTO' 999; 8: Y:=WY; 'GOTO' NXY; 'END'; 7: WY:=TTAG[Y-1]; 'BEGIN' 'IF' ~(WY=0) 'THEN' 'GOTO' 11; TTAG[Y-1]:=X; RESERVEADMINSYMBOL; 'GOTO' 999; 11: Y:=WY; 'GOTO' N
XY; 'END'; 'END'; 999: 'END';
73 'PROCEDURE RESERVEADMINSYMBOL; 'BEGIN' PTAG:=PTAG+8; 'IF' ~(PTAG<110000) 'THEN' 'GOTO' 1; TTAG(PTAG-8):=999; TTAG(PTAG-7):=0; TTAG(PTAG-6):=0; TTAG(PTAG-5):=0; TTAG(PTAG-4):=0; TTAG(PTAG-3):=0; TTAG(PTAG-2):=0; TTAG(PTAG-1):=0; 'GOTO' 999; 1: ERRORC(TAGFULL,0); POSTMORTEM: EXIT; 999: 'END';
74 'COMMENT' 4.3.3 BOLD SYMBOLS;
75 'POCLEAN' 'PROCEDURE READBOLD(X); 'INTEGER' X; 'BEGIN' 'INTEGER' B1,B2,B3,D; 'IF' ~ACCENT 'THEN' 'GOTO' 1; X:=PBOLD; 'BEGIN' 'IF' ~ACCENT 'THEN' 'GOTO' 2;
D:= (B3*128+B2)*128+B3; D:=D; STACKBOLD(D); READBOLD:= 'TRUE'; 'GOTO' 999; 2: 'IF' ~BOLDCHAR(B1) 'THEN' 'GOTO' 3; RST: 'BEGIN' 'IF' ~ACCENT 'THEN' 'GOTO' 4; D:= (B1
*128+B3)*128+B3; D:=D; STACKBOLD(D); READBOLD:= 'TRUE'; 'GOTO' 999; 4: 'IF' ~BOLDCHAR(B2) 'THEN' 'GOTO' 5; 'BEGIN' 'IF' ~ACCENT 'THEN' 'GOTO' 6; D:= (B1*128+B2)*128
+B3; D:=D; STACKBOLD(D); READBOLD:= 'TRUE'; 'GOTO' 999; 6: 'IF' ~BOLDCHAR(B3) 'THEN' 'GOTO' 7; D:= (B1*128+B2)*128+B3; 'BEGIN' 'IF' ~ACCENT 'THEN' 'GOTO' 8; D:=D; D:=D+
1; 'GOTO' 999; 8: 999: 'END';

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TACKBOLD);READBOLD:=TRUE;'GOTO'99;8:'IF'-(BOLDCHAR(B1))THEN'GOTO'9;STACKBOLD(D);'GOTO'RST;9:'GOTO'0;'END';7:'GOTO'0;'END';5:'GOTO'0;'END';3
:'GOTO'0;'END';1:0:READBOLD:=FALSE;'999:'END';
76 'PRCCEDURE'STACKBOLD(X);'INTEGER'X;'BEGIN'IF'-(PBOLD<200600)THEN'GOTO'1;TBOLD(PBOLD):=X;PBOLD:=PBOLD+1;'GOTO'999;1:ERROR(BOLDFULL,0);
PCSTMORTEW:EXIT;999:'END';
77 'PRCCEDURE'ENTERBOLD(X);'INTEGER'X;'BEGIN'INTEGER'Y,X1,Y1,WX,WY;Y:=200001;NXY:X1:=X;Y1:=Y;NXY:WX:=TBOLD(X1);WY:=TBOLD(Y1);'BEGIN'IF'-(
WX=WY)THEN'GOTO'2;'BEGIN'IF'-(WX<0)THEN'GOTO'3;'BEGIN'IF'-(Y<0)THEN'GOTO'4;'GOTO'999;4:PBOLD:=X;X:=Y;'GOTO'999;'END';3:X1:=X1+1;Y1:=Y1+1
;'GOTO'NXY;'END';2:SKP:'IF'-(WY<0)THEN'GOTO'7;Y:=Y1;Y1:=Y+1;'GOTO'NXY;7:Y1:=Y1+1;WY:=TBOLD(Y1);'GOTO'SKP;'END';999:'END';
78 'COMMENT'4.3.4SPECIALSYMBOLS;
79 'BOCLEAN'PROCEDURE'READSPEC(X);'INTEGER'X;'BEGIN'INTEGER'S,S1;'IF'-(SPECCHAR(S1))THEN'GOTO'1;X:=PSPEC;S1:=S1*128+63;S1:=S1+128+63;S:=S;STAC
KSPEC(S);READSPEC:=TRUE;'GOTO'999;1:READSPEC:=FALSE;'999:'END';
80 'PRCCEDURE'STACKSPEC(X);'INTEGER'X;'BEGIN'IF'-(PSPEC<300100)THEN'GOTO'1;TSPEC(PSPEC):=X;PSPEC:=PSPEC+1;'GOTO'999;1:ERROR(SPECFULL,0);
PCSTMORTEW:EXIT;999:'END';
81 'PRCCEDURE'ENTERSPEC(X);'INTEGER'X;'BEGIN'INTEGER'Y,WX,WY;Y:=300001;WX:=TSPEC(X);NXY:WY:=TSPEC(Y);'BEGIN'IF'-(WX=WY)THEN'GOTO'2;'BEG
IN'IF'-(Y<0)THEN'GOTO'3;'GOTO'999;3:PSPEC:=X;X:=Y;'GOTO'999;'END';2:Y:=Y+1;'GOTO'NXY;'END';999:'END';
82 'COMMENT'4.3.5CONSTANTS;
83 'BOCLEAN'PROCEDURE'READCONS(X);'INTEGER'X;'BEGIN'INTEGER'D,S;'IF'-(DIGIT(D))THEN'GOTO'1;S:=D;RST:'BEGIN'IF'-(DIGIT(D))THEN'GOTO'2;S:=
S*10+D;'GOTO'RST;2:X:=PCONS;STACKCONS(S);READCONS:=TRUE;'GOTO'999;'END';1:READCONS:=FALSE;'999:'END';
84 'PRCCEDURE'STACKCONS(X);'INTEGER'X;'BEGIN'IF'-(PCONS<400300)THEN'GOTO'1;TCONS(PCONS):=X;PCONS:=PCONS+1;'GOTO'999;1:ERROR(CONSFULL,0);
PCSTMORTEW:EXIT;999:'END';
85 'PRCCEDURE'ENTERCONS(X);'INTEGER'X;'BEGIN'INTEGER'Y,S,T;S:=TCONS(X);Y:=400001;NXT:T:=TCONS(Y);'BEGIN'IF'-(S=T)THEN'GOTO'2;'BEGIN'IF
'-(X=Y)THEN'GOTO'3;'GOTO'999;3:PCONS:=X;X:=Y;'GOTO'999;'END';2:Y:=Y+1;'GOTO'NXT;'END';999:'END';
86 'COMMENT'4.3.6INPUTADMINISTRATION;
87 'PRCCEDURE'INITIALIZEFORREADING;'BEGIN'GIVETEXT:=FALSE;'PCONS:=400001;PTAG:=100001;PBOLD:=200001;PSPEC:=300001;LINE:=0;CHAR:=RESYM;RESE
RVEADMINPAGE;FIRSTTAG:=PTAG;'END';
88 'BOCLEAN'PROCEDURE'STAG(X);'INTEGER'X;'BEGIN'IF'-(100001<LE'INPT<INPT<PTAG)THEN'GOTO'1;X:=INPT;NEXTSYMBOL;ISTAG:=TRUE;'GOTO'999;1
:ISTAG:=FALSE;'999:'END';
89 'BOCLEAN'PROCEDURE'ISBOLD(X);'INTEGER'X;'BEGIN'IF'-(200001<LE'INPT<INPT<PBOLD)THEN'GOTO'1;X:=INPT;NEXTSYMBOL;ISBOLD:=TRUE;'GOTO'99
9;1:ISBOLD:=FALSE;'999:'END';
90 'BOCLEAN'PROCEDURE'ISCONS(X);'INTEGER'X;'BEGIN'IF'-(400001<LE'INPT<INPT<PCONS)THEN'GOTO'1;X:=INPT;NEXTSYMBOL;ISCONS:=TRUE;'GOTO'99
9;1:ISCONS:=FALSE;'999:'END';
91 'BOCLEAN'PROCEDURE'RAHEAD(X);'INTEGER'X;'BEGIN'IF'-(AHEAD(X))THEN'GOTO'1;NEXTSYMBOL;R:=TRUE;'GOTO'999;1:R:=FALSE;'999:'END';
92 'BOCLEAN'PROCEDURE'RAHEAD(X);'INTEGER'X;'BEGIN'IF'-(INPT=X)THEN'GOTO'1;AHEAD:=TRUE;'GOTO'999;1:AHEAD:=FALSE;'999:'END';
93 'PRCCEDURE'NEXTSYMBOL;'BEGIN'READ(INPT);'END';
94 'PRCCEDURE'READ(X);'INTEGER'X;'BEGIN'SKP:'IF'-(LAYOUTSYMBOL)THEN'GOTO'1;NEXTCHAR;'GOTO'SKP;1:'IF'-(READTAG(X))THEN'GOTO'2;ENTERTAG(X);'G
OTO'999;2:'IF'-(READBOLD(X))THEN'GOTO'3;ENTERBOLD(X);'GOTO'999;3:'IF'-(READSPEC(X))THEN'GOTO'4;ENTERSPEC(X);'GOTO'999;4:'IF'-(READCONS(X))THEN'G
OTO'5;ENTERCONS(X);'GOTO'999;5:ERROR(WRONGINPUTCODE,CHAR);NEXTCHAR;READ(X);999:'END';
95 'COMMENT'4.4COMPILERCOMPILETCALGOL60;
96 'BOCLEAN'GIVERSTO,ADVANCED;
97 'INTEGER'HANDLE,BOUNDX,SPECX,TONEXT,TOEND,CURLAB,NEWLAB;
98 'COMMENT'4.4.1AUXILIARYACTIONS;
99 'PRCCEDURE'ADDPAGE(TAG);'INTEGER'TAG;'BEGIN'INTEGER'H,K,Q;H:=TTAG(TAG-6);'BEGIN'IF'-(H=0)THEN'GOTO'2;REPLACE(TAG,H);'GOTO'999;2:3:=
LTEXT(H);K:=Q//8192*8192;'IF'-(K<LINE)THEN'GOTO'3;REPLACE(TAG,H);'GOTO'999;3:'GOTO'999;'END';999:'END';
100 'PRCCEDURE'REPLACE(TAG,H);'INTEGER'H,TAG;'BEGIN'TTAG(TAG-6):=PTEXT;'BEGIN'IF'-(H=0)THEN'GOTO'2;ENTERTEXT(LINE);'GOTO'999;2:H:=H+80000
1;H:=H+1;P:=8192*H+LINE;ENTERTEXT(H);'GOTO'999;'END';999:'END';
101 'PRCCEDURE'FORGETLOCALS;'BEGIN'INTEGER'LCC;RST:'IF'-(PLOC=500001)THEN'GOTO'1;'GOTO'999;1:PLOC:=PLOC-1;LOC:=LLOC(PLOC);WARN(LCC);TTAG[
LCC-5]:=0;TTAG[LOC-4]:=0;TTAG[LOC-3]:=0;'GOTO'RST;999:'END';
102 'PRCCEDURE'WARN(X);'INTEGER'X;'BEGIN'INTEGER'STATE;STATE:=TTAG[X-5];'BEGIN'IF'-(STATE=BLANK)THEN'GOTO'2;TRACE(X);'GOTO'999;2:'IF'-(S
TATE=0)THEN'GOTO'3;TRACE(X);'GOTO'999;3:SIGNAL(X);'GOTO'999;'END';999:'END';
103 'PRCCEDURE'ENTERLOCAL(HEAD);'INTEGER'HEAD;'BEGIN'IF'-(PLOC<500200)THEN'GOTO'1;LLOC(PLOC):=HEAD;PLOC:=PLOC+1;'BEGIN'IF'-(PLOC<XLOC)'T
HEN'GOTO'2;'GOTO'999;2:XLOC:=PLOC;'GOTO'999;'END';1:ERROR(LOCFULL,0);POSTMORTEW:EXIT;999:'END';
104 'PRCCEDURE'ENTERGLOBAL(HEAD);'INTEGER'HEAD;'BEGIN'IF'-(PGLOB<601000)THEN'GOTO'1;LGLOB(PGLOB):=HEAD;PGLOB:=PGLOB+1;'GOTO'999;1:ERROR(G
LOBFULL,0);POSTMORTEW:EXIT;999:'END';
105 'PRCCEDURE'ENTERMACRO(HEAD);'INTEGER'HEAD;'BEGIN'IF'-(PMACR<701100)THEN'GOTO'1;LMACR(PMACR):=HEAD;PMACR:=PMACR+1;'GOTO'999;1:ERROR(MA
CRFULL,0);POSTMORTEW:EXIT;999:'END';
106 'PRCCEDURE'ENTERTEXT(X);'INTEGER'X;'BEGIN'IF'-(PTEXT<804000)THEN'GOTO'1;LTEXT(PTEXT):=X;PTEXT:=PTEXT+1;'GOTO'999;1:ERROR(TEXTFULL,0);
PCSTMORTEW:EXIT;999:'END';
107 'COMMENT'4.4.2TREATMENTOFTYPES;
108 'PRCCEDURE'DEFINACTION(HEAD);'INTEGER'HEAD;'BEGIN'REDEFINE(HEAD,DEFINED,GLOBAL,ACTION);'END';
109 'PRCCEDURE'DEFINPREDICATE(HEAD);'INTEGER'HEAD;'BEGIN'REDEFINE(HEAD,DEFINED,GLOBAL,PREDICATE);'END';
110 'PRCCEDURE'DEFINAFFIX(HEAD);'INTEGER'HEAD;'BEGIN'REDEFINE(HEAD,DEFINED,LOCAL,POINTER);'END';

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111 'PROCEDURE' DEFINE LABEL(LAB); 'INTEGER' LAB; 'BEGIN' 'IF' (100001 'LE' LAB < LAB < PTAG) 'THEN' 'GOTO' 1; 'REDEFINE' (LAB, DEFINED, LOCAL, LABEL); 'GOTO' 999; 1
112 'IF' (400001 'LE' LAB < LAB < PCONS) 'THEN' 'GOTO' 2; 'GOTO' 999; 2; 999: 'END';
113 'PROCEDURE' APPLY ACTION(HEAD); 'INTEGER' HEAD; 'BEGIN' 'REDEFINE' (HEAD, APPLIED, GLOBAL, ACTION); 'END';
114 'PROCEDURE' APPLY PREDICATE(HEAD); 'INTEGER' HEAD; 'BEGIN' 'REDEFINE' (HEAD, APPLIED, GLOBAL, PREDICATE); 'END';
115 'PROCEDURE' APPLY LABEL(LAB); 'INTEGER' LAB; 'BEGIN' 'IF' (100001 'LE' LAB < LAB < PTAG) 'THEN' 'GOTO' 1; 'REDEFINE' (LAB, APPLIED, LOCAL, LABEL); 'GOTO' 999; 1;
116 'IF' (400001 'LE' LAB < LAB < PCONS) 'THEN' 'GOTO' 2; 'GOTO' 999; 2; 999: 'END';
117 'PROCEDURE' APPLY(X); 'INTEGER' X; 'BEGIN' 'NEWPLACE' (X); 'IF' (TRYSTATE(X, APPLIED)) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
118 'PROCEDURE' REDEFINE(X, NSTATE, NSORT, NTYPE); 'INTEGER' X; 'BEGIN' 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
119 'GOTO' 1; 'NEWPLACE' (X); 'GOTO' 999; 1; CLDPOS: POS; NLCR: OUT(X); 'INFORM' (X); 'OUT' (NSTATE); 'OUT' (NSTATE); 'OUT' (NSTATE); 'POSITION' (CLDPOS); 999: 'END';
120 'BOCLEAN' 'PROCEDURE' REDEFINE(X, NSTATE, NSORT, NTYPE); 'INTEGER' X; 'BEGIN' 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
121 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
122 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
123 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
124 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
125 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
126 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
127 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
128 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
129 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
130 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
131 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
132 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
133 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
134 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
135 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
136 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
137 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
138 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';
139 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'INTEG' (X); 'IF' (Q=0) 'THEN' 'GOTO' 1; 'TRYSTATE' (X, NSTATE, NSORT, NTYPE); 'END';

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GOTO'1;REDEFINE(HEAD,DEFINED,GLOBAL,TYPE),G(HEAD),TYPE,'G(HEAD)',BEGIN'IF'-(R(COMMA)),THEN'GOTO'2;G(COMMA)),THEN'GOTO'1;IF'-(R(POINT)),THEN'GOTO'3;
G(SEMICOLON)),GOTO'999;3;'GOTO'0;'END';10;999;'END';
140 'BOCLEAN','PROCEDURE'FLAGDECLARATION;'BEGIN'IF'-(R(FLAG)),THEN'GOTO'1;TREATDECLLIST(BOCLEAN,FLAG);FLAGDECLARATION:=TRUE;'GOTO'999;1;FL
AGDECLARATION:=FALSE;'999;'END';
141 'BOCLEAN','PROCEDURE'LISTDECLARATION;'BEGIN'INTEGER'HEAD;'IF'-(R(LIST)),THEN'GOTO'1;BLANKLINE;G(QUANTITY);G(ARRAY);UNEXT;'IF'-(ISTAG(HEAD
)),'THEN'GOTO'0;G(HEAD);REDEFINE(HEAD,DEFINED,GLOBAL,LIST);'IF'-(R(SUB)),THEN'GOTO'0;G(SUB);CONSTANTTEXT;'IF'-(R(COLON)),THEN'GOTO'0;G(COLON);CONST
ANTTEXT;'IF'-(R(BUS)),THEN'GOTO'0;G(BUS);'BEGIN'IF'-(R(COMMA)),THEN'GOTO'2;G(COMMA);NEWLINE;'GOTO'NEXT;2;SHIFT2LINES;'IF'-(R(PCINT)),THEN'GOTO'3;G(
SEMICOLON);LISTDECLARATION:=TRUE;'GOTO'0;'END';10;LISTDECLARATION:=FALSE;'999;'END';
142 'PROCEDURE'CONSTANTTEXT;'BEGIN'INTEGER'SYMB;NEXT;'IF'-(ISTAG(SYMB)),THEN'GOTO'1;APPLY(SYMB);'BEGIN'IF'-(WASPARAMLESSMACRO(SYMB)),THEN'GOTO'1;GOTO
0;2;PUTPARAMLESSMACRO(SYMB);'GOTO'NEXT;2;G(SYMB);'GOTO'NEXT;'END';11;'IF'-(ISCONS(SYMB)),THEN'GOTO'4;G(SYMB);'GOTO'NEXT;4;'IF'-(R(PLUS)),THEN'GOTO'5;G
(PLUS);'GOTO'NEXT;5;'IF'-(R(MINUS)),THEN'GOTO'6;G(MINUS);'GOTO'NEXT;6;'END';
143 'COMMENT'4,4,5AFFIXEXPRESSIONS;
144 'PROCEDURE'PUTAFFIXEXPRESSION(HEAD);'BEGIN'IF'-(TRYSORT(HEAD,GLOBAL)),THEN'GOTO'1;PUTDIRECT(HEAD);'GOTO'999;1;'IF'-(TRYSORT
(HEAD,MACRO)),THEN'GOTO'2;'BEGIN'IF'-(WASACTION(HEAD)),THEN'GOTO'3;PUTMACRO(HEAD);G(CLOSE));'GOTO'999;'END';12;
999;'END';
145 'PROCEDURE'PUTDIRECT(HEAD);'BEGIN'INTEGER'HEAD;'BEGIN'INTEGER'NPARS;NPARS:=0;G(HEAD);AFFIXPACKOPTION(NPARS);NEWNPARS(HEAD,NPARS);'END';
146 'PROCEDURE'AFFIXPACKOPTION(NPARS);'BEGIN'IF'-(R(PLUS)),THEN'GOTO'1;G(OPEN);TREATAFFIX;RST;'BEGIN'NPARS:=NPARS+1;'IF'-(R(PL
US)),THEN'GOTO'2;G(COFMA);TREATAFFIX;'GOTO'RST;2;G(CLOSE));'GOTO'999;'END';11;999;'END';
147 'PROCEDURE'TREATAFFIX;'BEGIN'INTEGER'AFFIX;'IF'-(ISTAG(AFFIX)),THEN'GOTO'1;PUTSINGLE(AFFIX);APPLY(AFFIX);'GOTO'999;1;'IF'-(ISBOLD(AFFIX)),THEN'
'GOTO'2;G(STRINGQUOTE);G(AFFIX);G(STRINGQUOTE);'GOTO'999;2;'IF'-(ISCONS(AFFIX)),THEN'GOTO'3;G(AFFIX);'GOTO'999;3;ERROR(WRONGAFFIX,APFX);'999;'END';
148 'PROCEDURE'PUTSINGLE(X);'BEGIN'IF'-(WASAFFIX(X)),THEN'GOTO'1;'BEGIN'IF'-(WASMACRO(X)),THEN'GOTO'2;PUTPARAMLESSMACRO(X);'GOTO'
999;2;G(X);'GOTO'999;'END';11;REDEFINE(X,APPLIED,GLOBAL,POINTER);G(X);'999;'END';
149 'PROCEDURE'PUTPARAMLESSMACRO(Y);'BEGIN'INTEGER'Y;'BEGIN'INTEGER'X,M;X:=TAG(Y-7);NEWNPARS(Y,0);NEXTX:=X+1;M:=LMACR[X];'BEGIN'IF'-(M=POINT),'T
HEN'GOTO'2;'GOTO'999;2;G(M);'GOTO'NEXT;'END';999;'END';
150 'PROCEDURE'PUTMACRO(X);'BEGIN'INTEGER'X;'BEGIN'INTEGER'P1,P2,P3,P4,P5,Q,M,NPARS;NPARS:=0;Q:=TAG(X-7);GETPARAMETERS(NPARS,P1,P2,P3,P4,P5);NEW
NPARS(X,NPARS);NEXTQ:=Q+1;M:=LMACR[Q];'BEGIN'IF'-(M=POINT)),THEN'GOTO'2;'GOTO'999;2;'IF'-(M=ONE)),THEN'GOTO'3;PUTPAR(P1,1);'GOTO'NEXT;3;'IF'-(M=
TWO)),THEN'GOTO'4;PUTPAR(P2,2);'GOTO'NEXT;4;'IF'-(M=THREE)),THEN'GOTO'5;PUTPAR(P3,3);'GOTO'NEXT;5;'IF'-(M=FOUR)),THEN'GOTO'6;PUTPAR(P4,4);'GOTO'NEXT
;6;'IF'-(M=FIVE)),THEN'GOTO'7;PUTPAR(P5,5);'GOTO'NEXT;7;'IF'-(WASPARAMLESSMACRO(M)),THEN'GOTO'8;PUTPARAMLESSMACRO(M);'GOTO'NEXT;8;G(M);'GOTO'NEXT';
'END';999;'END';
151 'PROCEDURE'GETPARAMETERS(P,P1,P2,P3,P4,P5);'BEGIN'P5,P4,P3,P2,P1,P;'BEGIN'GETPAR(P,P2);GETPAR(P,P3);GETPAR(P,P4);GETPAR(P
,P5);'END';
152 'PROCEDURE'GETPAR(P,X);'BEGIN'IF'-(R(PLUS)),THEN'GOTO'1;P:=P+1;'BEGIN'IF'-(ISTAG(X)),THEN'GOTO'2;APPLY(X);'GOTO'999;2;'IF'-(
ISCONS(X)),THEN'GOTO'3;'GOTO'999;3;'GOTO'0;'END';11;X:=XPARAMETERERROR;'GOTO'999;0;'999;'END';
153 'PROCEDURE'PUTPAR(X,Y);'BEGIN'IF'-(X=XPARAMETERERROR),THEN'GOTO'1;G(X);'GOTO'999;1;'IF'-(X=XPARAMETERERROR),THEN'GOTO'
2;ERROR(X,Y);'GOTO'999;2;'IF'-(X=200001,'LE'X<X<PBOLE)),THEN'GOTO'3;G(STRINGQUOTE);G(X);'GOTO'999;3;PUTSINGLE(X);'999;'END';
154 'COMMENT'4,4,6BUILDINGSTONESOFARULE;
155 'PROCEDURE'PUTCONNECTAND;'BEGIN'IF'-(GIVERSTO),THEN'GOTO'1;PUTJUMP(CURLAB);TONEXT:=TONEXT+1;'GOTO'999;1;'IF'-(ADVANCED),THEN'GOTO'2;ERROR
(POSSIBLYPACKTRACKNECESSARY,0);PUTJUMP(0);TOLOST:=TOLOST+1;'GOTO'999;2;PUTJUMP(CURLAB);TONEXT:=TONEXT+1;'999;'END';
156 'PROCEDURE'PUTCONNECTOR;'BEGIN'IF'-(AHEAD(SEMICOLON)),THEN'GOTO'1;PUTRESTORE;FRESHLABEL;'BEGIN'IF'-(TONEXT=0)),THEN'GOTO'2;ERRRCR(ALTERN
ATIVE,NEVERREACHED,0);'GOTO'999;2;'GOTO'999;'END';11;'IF'-(AHEAD(CLOSE)),THEN'GOTO'4;'BEGIN'IF'-(TONEXT=0)),THEN'GOTO'5;'GOTO'999;5;PUTLABEL(CURLA
B);FRESHLABEL;'BEGIN'IF'-(GIVERSTO),THEN'GOTO'7;'GOTO'999;7;PUTJUMP(0);TOLOST:=TOLOST+1;'GOTO'999;'END';'END';4;PUTRESTORE;LOSTLABEL;'999;'END';
157 'PROCEDURE'LOSTLABEL;'BEGIN'IF'-(TOLOST=0)),THEN'GOTO'1;'BEGIN'IF'-(TONEXT=0)),THEN'GOTO'2;'GOTO'999;2;FALSERESULT;'GOTO'999;'END';11;P
UTLABEL(0);FALSERESULT;'999;'END';
158 'PROCEDURE'ENDLABEL;'BEGIN'IF'-(TOEND=0)),THEN'GOTO'1;'GOTO'999;1;PUTLABEL(999);'999;'END';
159 'PROCEDURE'ENDJUMP;'BEGIN'IF'-(INPT=POINT)),THEN'GOTO'1;'IF'-(TOLOST=0)),THEN'GOTO'1;'IF'-(TONEXT=0)),THEN'GOTO'1;'GOTO'999;1;PUTJUMP(9
99);TOEND:=TOEND+1;'999;'END';
160 'PROCEDURE'PUTINIT;'BEGIN'IF'-(GIVERSTO),THEN'GOTO'1;G(INITRESTORE));'GOTO'999;1;'999;'END';
161 'PROCEDURE'PUTRESTORE;'BEGIN'IF'-(TONEXT=0)),THEN'GOTO'1;'GOTO'999;1;PUTLABEL(CURLAB);'BEGIN'IF'-(GIVERSTO),THEN'GOTO'3;G(DOORESTORE));'G
OTO'999;3;'GOTO'999;'END';999;'END';
162 'PROCEDURE'FRESHLABEL;'BEGIN'CURLAB:=NEWLAB;NEWLAB:=NEWLAB+1;'END';
163 'PROCEDURE'TRUESERESULT;'BEGIN'IF'-(WASACTION(HANDLE)),THEN'GOTO'1;'GOTO'999;1;G(HANDLE);BECOMES;G(OTRUE);SEMICOL;'999;'END';
164 'PROCEDURE'FALSERESULT;'BEGIN'IF'-(WASACTION(HANDLE)),THEN'GOTO'1;'GOTO'999;1;G(HANDLE);BECOMES;G(OFALSE);SEMICOL;'999;'END';
165 'PROCEDURE'PUTJUMP(LAB);'BEGIN'G(QGOTO);G(LAB);G(SEMICOLON);NEWLINE;APPLYLABEL(LAB);'END';
166 'PROCEDURE'PUTLABEL(LAB);'BEGIN'LAB;'BEGIN'NEWLINE;G(LAB);G(LAB);G(COLON);U:DEFINELABEL(LAB);TABLINE;'END';
167 'PROCEDURE'SEMICOL;'BEGIN'G(SEMICOLON);'END';
168 'PROCEDURE'BECOMES;'BEGIN'G(COLON);G(EQUALS);'END';
169 'PROCEDURE'BEGIN;'BEGIN'NEWLINE;G(GBEGIN);U:TABLINE;'END';
170 'PROCEDURE'ENDBR;'BEGIN'NEWLINE;G(QEND);'END';
171 'COMMENT'4,4,7GENERALFORMOFARULE;
172 'BOCLEAN','PROCEDURE'RULE;'BEGIN'IF'-(LEFTHANDSIDE),THEN'GOTO'1;'IF'-(MIDDLE),THEN'GOTO'0;'IF'-(RIGHTHANDSIDE),THEN'GOTO'0;FORGETLCCALS;SHI

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FT2LINES: IF '-R(POINT)' THEN 'GOTO' 0; RULE: = 'TRUE'; 'GOTO' 999; 1: 0; RULE: = 'FALSE'; 999: 'END';
173 'EOCLEAN' 'PROCEDURE' 'LEFTHANDSIDE'; 'BEGIN' 'IF' '-ISTAG(HANDLE)' THEN 'GOTO' 1; 3: BLANKLINE; BOUNDX: = 0; SPECX: = 0; 'BEGIN' 'IF' '-WASACTION(HANDLE)' THEN
'GOTO' 2; 'DEFINEACTION(HANDLE); RESTLHSLEFTHANDSIDE: = 'TRUE'; 'GOTO' 999; 2: G(QBOOLEAN); 'DEFINEPREDICATE(HANDLE); RESTLHSLEFTHANDSIDE: = 'TRUE'; 'GOTO' 999
; 'END'; 1: LEFTHANDSIDE: = 'FALSE'; 999: 'END';
174 'PROCEDURE' 'RESTLHS'; 'BEGIN' G(OPRCEDURE); G(HANDLE); 'OPTIONALBOUNDAFFIXES'; NEWPARS(HANDLE, BOUNDX); SEMICOL; 'END';
175 'PROCEDURE' 'OPTIONALBOUNDAFFIXES'; 'BEGIN' 'INTEGER' AFFX; 'IF' '-BOUNDAFFIX1(AFFX)' THEN 'GOTO' 1; G(OPR); G(AFFX); RESTBOUNDAFFIXES; 'GOTO' 999; 1: 'IF'
'-BOUNDAFFIX2(AFFX)' THEN 'GOTO' 2; G(OPR); G(AFFX); RESTBOUNDAFFIXES; G(AFFX); 'GOTO' 999; 2: 999: 'END';
176 'PROCEDURE' 'RESTBOUNDAFFIXES'; 'BEGIN' 'INTEGER' AFFX; 'IF' '-BOUNDAFFIX1(AFFX)' THEN 'GOTO' 1; G(COMMA); G(AFFX); RESTBOUNDAFFIXES; 'GOTO' 999; 1: 'IF'
'-BOUNDAFFIX2(AFFX)' THEN 'GOTO' 2; G(COMMA); G(AFFX); RESTBOUNDAFFIXES; G(AFFX); 'GOTO' 999; 2: 999: 'END';
177 'EOCLEAN' 'PROCEDURE' 'BOUNDAFFIX1(AFFX)'; 'BEGIN' 'IF' '-ISTAG(AFFX)' THEN 'GOTO' 1; 'IF' '-ISTAG(AFFX)' THEN 'GOTO' 0; 'DEFINEAFFIX(AFFX)'; BOUNDX: = BOUNDX+1; BOUNDAFFIX1: = 'TRUE'; 'GOTO' 999; 1: 0; BOUNDAFFIX1: = 'FALSE'; 999: 'END';
178 'EOCLEAN' 'PROCEDURE' 'BOUNDAFFIX2(AFFX)'; 'BEGIN' 'IF' '-R(PLUS)' THEN 'GOTO' 1; 'IF' '-ISTAG(AFFX)' THEN 'GOTO' 0; 'DEFINEAFFIX(AFFX)'; BOUNDX: = BOUNDX+1; SPECX: = SPECX+1; BOUNDAFFIX2: = 'TRUE'; 'GOTO' 999; 1: 0; BOUNDAFFIX2: = 'FALSE'; 999: 'END';
179 'EOCLEAN' 'PROCEDURE' 'MIDDLE'; 'BEGIN' 'BEGIN' 'OPTIONALPREAFFIXES'; 'IF' '-R(COLON)' THEN 'GOTO' 1; PUTINIT; MIDDLE: = 'TRUE'; 'GOTO' 999; 1: MIDDLE: = 'FALSE'
; 999: 'END';
180 'PROCEDURE' 'OPTIONALPREAFFIXES'; 'BEGIN' 'INTEGER' AFFX; 'IF' '-R(MINUS)' THEN 'GOTO' 1; G(QUINTEGER); FFX: 'IF' '-ISTAG(AFFX)' THEN 'GOTO' 0; G(AFFX); 'DEF
INEAFFIX(AFFX)'; 'BEGIN' 'IF' '-R(MINUS)' THEN 'GOTO' 2; G(COMMA); 'GOTO' FFX; 2: G(SEMICOLON); 'GOTO' 999; 'END'; 1: 'GOTO' 999; 0: 999: 'END';
181 'EOCLEAN' 'PROCEDURE' 'RIGHTHANDSIDE'; 'BEGIN' 'CURLAB: = 1; NEWLAB: = 2; TOEND: = 0; NEXT: 'IF' '-ALTERNATIVE' THEN 'GOTO' 1; PUTCONNECTOR; 'BEGIN' 'IF'
'-R(SEMICOLON)' THEN 'GOTO' 2; 'GOTO' NEXT; 2: ENDLABEL; ENDRSEMICOL; RIGHTHANDSIDE: = 'TRUE'; 'GOTO' 999; 'END'; 1: RIGHTHANDSIDE: = 'FALSE'; 999: 'END';
182 'EOCLEAN' 'PROCEDURE' 'ALTERNATIVE'; 'BEGIN' 'INTEGER' LAB, OLDNEXT; TONEXT: = 0; ADVANCED: = 'FALSE'; NEXT: 'IF' '-R(COMMA)' THEN 'GOTO' 1; 'GOTO' NEXT; 1: 'IF'
'-R(COLON)' THEN 'GOTO' 2; 'IF' '-ISTAG(LAB)' THEN 'GOTO' 0; PUTJUMP(LAB); ALTERNATIVE: = 'TRUE'; 'GOTO' 999; 2: 'IF' '-R(OPEN)' THEN 'GOTO' 3; OLDNEXT: = TONEXT; 'IF'
'-R(CLOSE)' THEN 'GOTO' 0; 'IF' '-R(CLOSE)' THEN 'GOTO' 0; 'BEGIN' 'IF' '-GIVERSTO' THEN 'GOTO' 4; TONEXT: = 1; ALTERNATIVE: = 'TRUE'; 'GOTO' 999; 4: TONEXT: = OLDNEXT; ALTERN
ATIVE: = 'TRUE'; 'GOTO' 999; 'END'; 3: 'IF' '-MEMBER' THEN 'GOTO' 6; 'GOTO' NEXT; 6: TRUERESULT; ENDJUMP; ALTERNATIVE: = 'FALSE'; 999: 'END';
183 'EOCLEAN' 'PROCEDURE' 'CHOICE'; 'BEGIN' 'INTEGER' LAB1, LAB2; CURLAB: = LAB; BEGIN; PUTINIT; FRESHLABEL; NEXT: 'IF' '-ALTERNATIVE' THEN 'GOTO' 1; PUTCONNECTOR; 'BE
GIN' 'IF' '-R(SEMICOLON)' THEN 'GOTO' 2; 'GOTO' NEXT; 2: ENDR; SEMICOL; CURLAB: = LAB; CHOICE: = 'TRUE'; 'GOTO' 999; 'END'; 1: CHOICE: = 'FALSE'; 999: 'END';
184 'EOCLEAN' 'PROCEDURE' 'MEMBER'; 'BEGIN' 'INTEGER' HEAD; 'IF' '-ISTAG(HEAD)' THEN 'GOTO' 1; APPLY(HEAD); 'BEGIN' 'IF' '-R(COLON)' THEN 'GOTO' 2; PUTLABEL(HEAD)
; MEMBER: = 'TRUE'; 'GOTO' 999; 2: 'IF' '-TRYTYPE(HEAD, PREDICATE)' THEN 'GOTO' 3; G(QUF); G(NOT); PUTAFFIXEXPRESSION(HEAD); G(QUF); PUTCONNECTANDADVANCED: = 'TRUE'; MEMBER: = 'TRUE'; 'GOTO' 999; 3: 'IF' '-TRYTYPE(HEAD, FLAG)' THEN 'GOTO' 4; G(QUF); G(NOT); PUTAFFIXEXPRESSION(HEAD); G(QUF); PUTCONNECTAND; MEMBER: = 'TRUE'; 'GOTO' 999; 4: 'IF' '-TRYTYPE(HEAD, ACTION)' THEN 'GOTO' 5; PUTAFFIXEXPRESSION(HEAD); SEMICOL; MEMBER: = 'TRUE'; 'GOTO' 999; 5: 'GOTO' 0; 'END'; 1: 0; MEMBER: = 'FALSE'; 999: 'END';
185 'COMMENT' '4.4. ROTHERBUILDINGSTONES OF A GRAMMAR';
186 'EOCLEAN' 'PROCEDURE' 'COMMAND'; 'BEGIN' 'IF' '-R(RSTOON)' THEN 'GOTO' 1; GIVERSTO: = 'TRUE'; COMMAND: = 'TRUE'; 'GOTO' 999; 1: 'IF' '-R(RSTOOF)' THEN 'GOTO' 2
; GIVERSTO: = 'FALSE'; COMMAND: = 'TRUE'; 'GOTO' 999; 2: 'IF' '-R(SHORT)' THEN 'GOTO' 3; LEGIBLE: = 'FALSE'; COMMAND: = 'TRUE'; 'GOTO' 999; 3: 'IF' '-R(LONG)' THEN 'GOTO' 4
; LEGIBLE: = 'TRUE'; COMMAND: = 'TRUE'; 'GOTO' 999; 4: 'IF' '-R(TRACEON)' THEN 'GOTO' 5; GIVE-RACE: = 'TRUE'; COMMAND: = 'TRUE'; 'GOTO' 999; 5: 'IF' '-R(TRACEOFF)' THEN 'GOTO' 6; GIVE-RACE: = 'FALSE'; COMMAND: = 'TRUE'; 'GOTO' 999; 6: COMMAND: = 'FALSE'; 999: 'END';
187 'EOCLEAN' 'PROCEDURE' 'COMMENT'; 'BEGIN' 'IF' '-R(SUB)' THEN 'GOTO' 1; BLANKLINE; G(QUCOMMENT); RST: 'BEGIN' 'IF' '-R(BUS)' THEN 'GOTO' 2; G(SEMICOLON); COMME
NT: = 'TRUE'; 'GOTO' 999; 2: G(INPT); READ(INPT); 'GOTO' RST; 'END'; 1: COMMENT: = 'FALSE'; 999: 'END';
188 'EOCLEAN' 'PROCEDURE' 'STARTINGSYMBOL'; 'BEGIN' 'INTEGER' HEAD; 'IF' '-R(RESULT)' THEN 'GOTO' 1; 'IF' '-ISTAG(HEAD)' THEN 'GOTO' 0; NEWPAGE; 'BEGIN' 'IF' '-R
TYPE(HEAD, ACTION)' THEN 'GOTO' 2; APPLYACTION(HEAD); TERMINALS; G(HEAD); STARTINGSYMBOL: = 'TRUE'; 'GOTO' 999; 2: APPLYPREDICATE(HEAD); TERMINALS; G(HEAD); ST
ARTINGSYMBOL: = 'TRUE'; 'GOTO' 999; 'END'; 1: 0; STARTINGSYMBOL: = 'FALSE'; 999: 'END';
189 'PROCEDURE' 'SKIPUNTILPOINT'; 'BEGIN' NEXT; ERROR(SKIPPED, INPT); 'BEGIN' 'IF' '-R(POINT)' THEN 'GOTO' 2; 'GOTO' 999; 2: NEXTSYMBOL; 'GOTO' NEXT; 'END'; 999: 'E
ND';
190 'COMMENT' '4.4.9 TERMINALS';
191 'PROCEDURE' 'TERMINALS'; 'BEGIN' 'BLANKLINE; FORGETGLOBALS; NEWLINE; 'IF' '-(PLOC=500001)' THEN 'GOTO' 1; 'GOTO' 999; 1: APPLYACTION(XREAD); APPLYACTION(X
OUT); APPLYACTION(INITREAD); BLANKLINE; G(INITREAD); G(SEMICOLON); NEWLINE; BLANKLINE; 999: 'END';
192 'PROCEDURE' 'FORGETGLOBALS'; 'BEGIN' 'INTEGER' P, TERM; P: = 600001; PLOC: = 500001; NEXT: 'IF' '-(P=PGLOB)' THEN 'GOTO' 1; 'GOTO' 999; 1: TERM: = LGLOB(P); MAYSET
TERMINAL(TERM); P: = P+1; 'GOTO' NEXT; 999: 'END';
193 'PROCEDURE' 'MAYBETERMINAL(X)'; 'BEGIN' 'IF' '-WASTERMINAL(X)' THEN 'GOTO' 1; PUTDECL(X); ENTERLOCAL(X); 'GOTO' 999; 1: WARN(X); 999: 'END';
194 'EOCLEAN' 'PROCEDURE' 'WASTERMINAL(X)'; 'BEGIN' 'INTEGER' X; 'BEGIN' 'INTEGER' S; S: = TAG(X-5); 'IF' '-(S=APPLIED)' THEN 'GOTO' 1; S: = TAG(X-4); 'IF' '-(S=GLOBAL)'
THEN 'GOTO' 1; S: = TAG(X-3); 'IF' '-(S=POINT)' THEN 'GOTO' 1; WASTERMINAL: = 'TRUE'; 'GOTO' 999; 1: WASTERMINAL: = 'FALSE'; 999: 'END';
195 'PROCEDURE' 'PUTDECL(X)'; 'BEGIN' 'NEWLINE; G(QUINTEGER); G(X); G(SEMICOLON); POSITION(64); INFORM(X); 'END';
196 'PROCEDURE' 'FORGETTERMINALS'; 'BEGIN' 'INTEGER' P, TERM; P: = 500001; NEXT: 'IF' '-(P=PLOC)' THEN 'GOTO' 1; 'GOTO' 999; 1: TERM: = LLOC(P); PUTCALL(XREAD, TERM)
; P: = P+1; 'GOTO' NEXT; 999: 'END';
197 'PROCEDURE' 'PUTCALL(PROC, X)'; 'BEGIN' 'NEWLINE; G(OPEN); G(X); G(CLOSE); G(SEMICOLON); 'END';
198 'COMMENT' '4.4.10 POSTMORTEM';
199 'PROCEDURE' 'POSTMORTEM'; 'BEGIN' 'NLCR; OUTINT(100001); OUTINT(PTAG); OUTINT(110000); NLCR; OUTINT(200001); OUTINT(PBOLD); OUTINT(200600); NLCR; OUTIN
T(300001); OUTINT(PSPEC); OUTINT(300100); NLCR; OUTINT(400001); OUTINT(PCONS); OUTINT(400300); NLCR; OUTINT(500001); OUTINT(XLOC); OUTINT(500200); NLCR; OUT
INT(600001); OUTINT(PGLOB); OUTINT(601000); NLCR; OUTINT(700001); OUTINT(PMACR); OUTINT(701000); NLCR; OUTINT(800001); OUTINT(800100); NLCR; OUTINT(900001); OUTINT(900200); NLCR; OUTINT(900300); OUTINT(900400); NLCR; OUTINT(900500); OUTINT(900600); NLCR; OUTINT(900700); OUTINT(900800); NLCR; OUTINT(900900); OUTINT(901000); NLCR; OUTINT(901100); OUTINT(901200); NLCR; OUTINT(901300); OUTINT(901400); NLCR; OUTINT(901500); OUTINT(901600); NLCR; OUTINT(901700); OUTINT(901800); NLCR; OUTINT(901900); OUTINT(902000); NLCR; OUTINT(902100); OUTINT(902200); NLCR; OUTINT(902300); OUTINT(902400); NLCR; OUTINT(902500); OUTINT(902600); NLCR; OUTINT(902700); OUTINT(902800); NLCR; OUTINT(902900); OUTINT(903000); NLCR; OUTINT(903100); OUTINT(903200); NLCR; OUTINT(903300); OUTINT(903400); NLCR; OUTINT(903500); OUTINT(903600); NLCR; OUTINT(903700); OUTINT(903800); NLCR; OUTINT(903900); OUTINT(904000); NLCR; OUTINT(904100); OUTINT(904200); NLCR; OUTINT(904300); OUTINT(904400); NLCR; OUTINT(90450
```

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001 OUTINT(LINE);CUTINT(CARD);NLCR;NLCR;NLCR;DICTIONARY;'END';
002 'PROCEDURE'DICTIONARY;'BEGIN'LISTTAGS(FIRSTTAG);'END';
003 'PROCEDURE'LISTTAGS(X);'BEGIN'INTEGER'X';
004 TTAG(X-5);SORT:=TTAG(X-4);TYPE:=TTAG(X-3);'IF'-(SORT=0)'THEN'GOTO'2;'GOTO'SKP;2:NLCR;OUT(X);POSITION(32);OUT(STATE);OUT(TYPE);POSITION(59);ENTRIES(X);SKP:P:=TTAG(X-1);LISTTAGS(P);999:'END';
005 'PROCEDURE'ENTRIES(P);'BEGIN'INTEGER'P';
006 'PROCEDURE'ENTRIES(P);'BEGIN'INTEGER'P';
007 'K:=H-81;C=CUTINT(K);'BEGIN'IF'-(Q=0)'THEN'GOTO'3:NLCR;'GOTO'999;3:Q:=Q+800001;Q:=Q-1;'GOTO'LNE;'END';
008 'COMMENT'4.5:'HEARTOFCOMPILERCOMPILER';
009 'INTEGER'INTRESTORE,DORSTORE,XREAD,XOUT,INTREAD,TITLE;
010 'PROCEDURE'SENTENCE;'BEGIN'STARTSYSTEM;BLANKLINE;BEGIN;
011 'PROCEDURE'STARTSYSTEM;'BEGIN'READ(INTRESTORE);READ(XREAD);READ(OUT);READ(TITLE);GIVERSTO:=FALSE;'GIVE
012 TRACE:=FALSE;'GIVETEXT:=TRUE;'PLOC:=500001;PLOC:=600001;PTEXT:=800001;PMACR:=700001;CARD:=0;CPOS:=1;XLOC:=0;LINEDUP:=FALSE;'LEGIBLE:=TRUE';
013 IDENTAT,CN:=0;LINE:=0;POS:=0;OUT(TITLE);READ(INT);'END';
014 'PROCEDURE'COMPILERDESCRIPTION;'BEGIN'NEXT:'IF'-SPECIFICATION'THEN'GOTO'1;'GOTO'NEXT;1:'GOTO'NEXT;1:'IF'-DECLARATION'THEN'GOTO'2;'GOTO'NEXT;2:'IF'-COM
015 MAND'THEN'GOTO'3;'GOTO'NEXT;3:'IF'-COMMENT'THEN'GOTO'4;'GOTO'NEXT;4:'IF'-STARTINGSYMBOL'THEN'GOTO'5;'GOTO'999;5:'IF'-RULE'THEN'GOTO'6;'GOTO'NEXT;6:SKIPJUNTILPOINT;'GOTO'NEXT;999:'END';
016 'INTEGER'TAGFULL;'INTEGER'BCDFULL;'INTEGER'SPECFULL;'INTEGER'GLOBAL;'INTEGER'CONFULL;'INTEGER'WRONGINPUTCODE;'INTEGER'LOCAL;'INTEGER'POINTER
017 GLOFFULL;'INTEGER'MACRFULL;'INTEGER'TEXTFULL;'INTEGER'DEFINED;'INTEGER'GLOBAL;'INTEGER'PREDICATE;'INTEGER'LOCAL;'INTEGER'POINTER
018 'INTEGER'LABEL;'INTEGER'APPLIED;'INTEGER'XIMPOSSIBLE;'INTEGER'MACRO;'INTEGER'WRONGNUMBEROFPARAMETERS;'INTEGER'FLAG;'INTEGER'LIST;'INTEGER'EXTERN
019 AL;'INTEGER'COMMA;'INTEGER'POINT;'INTEGER'QUINTEGER;'INTEGER'SEMICOLON;'INTEGER'QBOOLEAN;'INTEGER'QARRAY;'INTEGER'SUB;'INTEGER'COLON;'INTEGER'BU
020 S;'INTEGER'PLUS;'INTEGER'MINUS;'INTEGER'OPEN;'INTEGER'CLOSE;'INTEGER'STRINGQUOTE;'INTEGER'ONE;'INTEGER'TWO;'INTEGER'THREE;'IN
021 TEGER'FOUR;'INTEGER'FIVE;'INTEGER'POSSIBLYBACKTRACKNECESSARY;'INTEGER'ALTERNATIVENEVERREACHED;'INTEGER'QTRUE;'INTEGER'
022 QFALSE;'INTEGER'QGOTO;'INTEGER'QEQUALS;'INTEGER'QEND;'INTEGER'QBEGIN;'INTEGER'QEND;'INTEGER'QBEGIN;'INTEGER'QEND;'INTEGER'QEND;'INTEGE
023 R'RTSTON;'INTEGER'RTSTOFF;'INTEGER'SHORT;'INTEGER'LONG;'INTEGER'TRACEON;'INTEGER'TRACEOFF;'INTEGER'QCOMMENT;'INTEGER'RESULT;'INTEGER'SKIP
024 PED;
025 INITIALIZEFORREADING;
026 READ(TAGFULL);READ(BOLDFULL);READ(SPECFULL);READ(CONSFULL);READ(WRONGINPUTCODE);READ(BLANK);READ(LOCFULL);READ(GLOFFULL);READ(MACRFULL);
027 READ(TEXTFULL);READ(DEFINED);READ(GLOBAL);READ(ACTION);READ(PREDICATE);READ(LOCAL);READ(POINT);READ(APPLIED);READ(XIMPOSSIBLE);REA
028 D(MACRO);READ(WRONGNUMBEROFPARAMETERS);READ(FLAG);READ(LIST);READ(QUINTEGER);READ(SEMICOLON);READ(QBOOLEAN);READ(QARRAY);READ(SUB);READ(COLON);READ(ONE);READ(TWO);READ(THREE);READ(FOUR);READ(FIVE);READ(XPARAMETERERROR);READ(POSSIBLYBACKTRACKNECESSARY);READ(ALTERNATIVENEVERREACHED);READ(QTRUE);READ(QFALSE);READ(QGOTO);READ(QEQUALS);READ(QEND);READ(QBEGIN);READ(QEND);READ(QCOMMENT);READ(RESULT);READ(TRACEON);READ(TRACEOFF);READ(DISPORT);READ(LONG);READ(EQUALS);READ(QBEGIN);READ(QEND);READ(QCOMMENT);READ(RESULT);READ(SKIPPED);
029 SENTENCE;'END';
030

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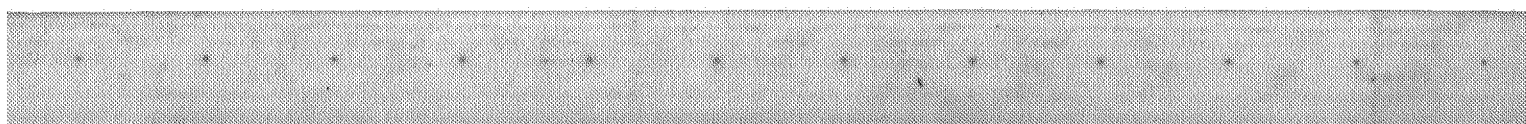

References.

- [1] J. Feldman, D. Gries, "Translator Writing Systems", Communications of the ACM, Volume 11 number 2 (1968).
- [2] C.H.A. Koster, "Affix Grammars", in: J.E.L. Peck (Editor), "ALGOL 68 implementation", North-Holland Publishing Company (1971).
- [3] A. van Wijngaarden (Editor), "Report on the Algorithmic Language ALGOL 68", Numerische Mathematik, 14 (1969).
- [4] P. Naur (Editor), "Revised Report on the Algorithmic Language ALGOL 60", Regnecentralen, Copenhagen (1962).
- [5] D.E. Knuth, "Semantics of Context Free Languages", Math. Systems Theory 2, (1968).
- [6] D.E. Knuth, "Top-down syntax analysis", International Summer School on Computer Programming, Copenhagen (1967).
- [7] J. Earley, H. Sturgis, "A Formalism for Translator Interactions", Communications of the ACM, Volume 13 number 10 (1970).

The differences between versions 17 and 19 of the CDL-compiler; first update to MR 127/71.

1. p. 23, bottom line, ignored, add:
; moreover any sequence of symbols, not containing a
\$-character, enclosed between \$-characters is also ignored
2. p. 24, l. 16, equals =, add:
not ¬
3. p. 24, l. 25, list ..., add:
comment symbol 'comment'
4. p. 25, l. 17, flag symbol, add:
; list symbol
5. p. 26, l. 2 from bottom, list (global, add:
or macro
6. p. 27, l. 13, 'external' ... flag, add:
'external' 'list' defined global list
7. p. 27, l. 19, 'macro'.... pointer, add:
'macro' 'list' defined macro list
8. p. 30, l. 8, 2.6. Comments. *input*; , replace by
2.6. Comments.
comment: skipped comment; translated comment.
skipped comment: sub, rest skipped comment.
rest skipped comment: bus; nonbus, rest skipped comment.
translated comment: comment symbol, rest translated comment.
rest translated comment: point; nonpointnonsemicolon,
rest translated comment.

(Here, nonbus stands for any symbol except] and nonpoint
nonsemicolon for any symbol except . or ; .)



A skipped comment is skipped and a translated comment is translated into an ALGOL 60 <comment>; e.g. [input] is skipped while 'comment' input. is translated into '*comment*' *input*;. The text of a comment may only consist of allowed CDL-symbols (see 2.2), as opposed to the sequence of characters between two \$-characters (see 2.2).

9. p. 30, bottom line, replace by:

(semicolon, right hand side;) .

10. p. 31, lines 1-8, replace by:

alternative:

label, colon, alternative;
affix expression, (comma, alternative;);
group; jump; .

affix expression:

handle, optional affixes;
not, tag.

11. p. 32, line 10, affix expression, add:

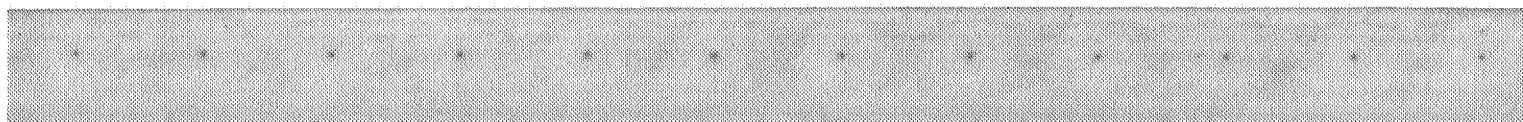
that starts with a handle

12. p. 33, after line 20, α ; , add:

The translation of an affix expression that starts with a not is:

if tag then goto λ ;

where λ is determined by means of the algorithm on the previous page.



13. p. 33, bottom, , add:

The above translation scheme makes use of integer labels, which, for some applications, are undesirable. Therefore a symbol is prefixed to every label in the translation; this symbol is the terminal "label prefix" of the CDL-compiler, which is read in during the initialization phase of the compiler. The prefix should be chosen so that the resulting names will not clash with existing names. The present value of "label prefix" is "l"; in the examples in this report its value is "" (i.e. the empty symbol).

14. p. 36, l. 4, delete: Various themselves.

15. p. 36, l. 25, occurrence. , add:

Line numbers of specification are preceded by S:,
line numbers of definition are preceded by D:.

16. after page 36 add the following pages:

2.10. Error detection and recovery.

All errors fall in one of the following four classes.

1) Warnings:

backtrack?	An alternative may fail although one of its predicates may have succeeded.
alternative never reached	One of the previous alternatives in this right hand side will always succeed.
nonfalse	This predicate will always succeed, and is in effect an action.
may be false	Possibly none of the alternatives of this action will succeed.

2) Unexpected symbols:

wrong input code	Some unreadable item is met on the input medium.
symbol missing, assumed: symbol	1. The grammar requires at this point the presence of the indicated symbol, which, however, was not found in the input. The symbol is assumed to be present, so that e.g. "'list' a[1:10." will yield a correct translation. 2. The grammar requires at this point the presence of a tag, which, however, was not found in the input. The tag "dummy" is inserted. 3. The grammar requires at this point the presence of one of the symbols: 'action' 'predicate'

'pointer'

'flag'

'list'

None of these, however, was found. The symbol 'predicate' was inserted.

symbol not allowed inside comment symbol

The indicated symbol violates the restrictions on the symbols allowed inside a comment. In this version of the compiler the symbol can only be a semicolon.

affix missing

The grammar requires at this point the presence of an affix, which, however, was not found. The tag "dummy" was inserted.

skipped --- state, sort, type, symbol

The indicated symbol was skipped because it was unacceptable as the first symbol of a compiler description.

3) Conflict of attributes

impossible redefinition state, sort, type, tag, new state, new sort,
new type

The indicated tag with the indicated state, sort and type is specified, defined or applied with the conflicting attributes new state, new sort and new type.

wrong number of parameters state, sort, type, tag

The number of affixed following the indicated tag is incompatible with that number in the definition of the tag.

wrong sort state, sort, type, tag

The indicated local tag is applied as a handle.

